Programmable In-Network Security for Context-aware BYOD Policies

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BYOD: Bring Your Own Device Risks



- BYOD devices: Less well managed and easier to be compromised
- Need to access control for BYOD clients

"Context-aware" policies for BYOD



- Making precise security decisions by dynamically adapting to security contexts
- How to enforce these policies?

State of the art: SDN-based defense



- Client modules collect client-side information
- BYOD policies are managed and enforced in an SDN "app"

Limitations of the SDN-based defense



- Low defense agility: Context updates need to traverse the software controller
- Vulnerable to control plane DoS attacks [AvantGuard CCS'13]
- Root cause: Lower processing speed of the SDN controller software

Research question

Can we address the limitations of SDN-based BYOD defense?



Opportunity: Programmable data planes



- Switch features:
 - Programmable parser: Customized protocols
 - Stateful processing
 - Arithmetic operations
 - General-purpose control plane
- High performance : <1us delay for Tbps traffic
- Can we transform these hardware features to security benefits for BYOD? 7

P4: Language for data plane programming

Customized headers

```
header myTunnel_t {
   bit<16> proto_id;
   bit<16> dst id;
```

```
struct headers {
    ethernet t
                 ethernet;
    myTunnel_t
                 myTunnel;
    ipv4_t
                 ipv4;
```

```
Match/action processing
table ipv4 lpm {
   key = {
       hdr.ipv4.dstAddr: lpm;
    }
   actions = {
       ipv4 forward;
       drop;
       NoAction;
   size = 1024;
   default action = NoAction();
```

Stateful registers

// count the number of bytes seen since the last probe register<bit<32>>(MAX PORTS) byte cnt reg; // remember the time of the last probe register<time_t>(MAX_PORTS) last_time_reg;

Reconfigures switch pipeline for header manipulation

}

- Has the potential to enforce BYOD policies at linespeed
- Downside: P4 is low-level, non-trivial to develop and maintain

Poise: Programmable In-network Security



- Language: An expressive language for defining BYOD policies
- **Compiler**: Generates device configurations and switch programs
- P4 data plane design: A dynamic and efficient security primitive

Outline

Motivation

- Poise Design
 - The Poise language
 - Compiling Poise policies
 - Data plane design
- Evaluation
- Conclusion

The Poise language



- An expressive language for writing context-aware policies
 - Predicates on customized client contexts
 - Support pre-defined primitive actions

Compiling Poise policies

<pre>if match (sslver <= 6.5)</pre>	table decision_tab
then drop	۱ key = { ctx.sslver: exact}
<pre>header ctx_t { sslver: 16 }</pre>	<pre>entries = { <= 6.5.0: dec = DROP > 6.5.0: dec = ALLOW } }</pre>

- Contexts (sslver) are compiled to customized header fields
- Security actions (if-else) are compiled to match/action table entries
- Advanced features: Policy composition, resource optimizations, etc

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An efficient in-network primitive



- **Problem**: How to spread context information from client to switch?
 - Strawman solution: Tag every packet with context -- high overhead!
- Idea: Periodic context packets per flow: Headers + context, no data
 - **Dynamic**: Decisions are based on the latest context
 - Efficient: Data packets unmodified (no embedded contexts)
 - Adjustable accuracy: Tunable context packet period

Key data structure: A per-flow table



Per-flow table

- A match/action table to maintain the latest per-flow decision
- Technical challenges:
 - New flow insertion delay (~1ms)
 - Controlling the size of the table
 - Handling DoS attacks (e.g., many new flows) See more details in our paper!

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Evaluation setup

- Prototype implementation
 - Compiler: Bison + Flex
 - Android client module: a kernel module on Linux 3.18.31
 - ~6000 LoC
- Evaluation setup
 - Tofino Wedge 100BF switch 32 X 100 Gbps = 3.2 Tbps

What we have evaluated

- **Correctness**: Can Poise enforce BYOD policies correctly?
- **Overhead**: How much delay or throughput degradation can Poise incur?
- Scalability: How complex/large policies can Poise support?
- Poise vs. SDN: Is Poise resilient to control plane saturation attacks?

- SDN-based solution: PBS NDSS'16
 - Floodlight v1.2 + Open vSwitch v2.9.2
- Methodology:
 - DoS attacker: Launch frequent context changes
 - Measure how normal user traffic are affected

Poise vs. SDN: First packet delay



- SDN: Takes ~1 second for the first packet to arrive under heavy attacks
- Poise: Remains at a constant level

Poise vs. SDN: New flow installation



- SDN: Fails to install new flows under heavy attacks
- Poise: Almost always installs 100% new flows
- Poise is highly resilient to DoS attacks from malicious clients

Conclusion

- Motivation
 - SDN-based BYOD defense has limitations
- Poise: Programmable In-Network Security
 - An expressive policy language
 - Compiler for generating P4 programs
 - An efficient in-network security primitive
- Poise transforms the hardware features to security benefits

Thank you for listening!

Contact: <u>qiaokang@rice.edu</u> – Looking for 2021 summer internship <u>https://github.com/qiaokang92/poise</u>