LODIC : Logical Distributed Counting for Scalable File Access

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Background

- The number of cores is rapidly increasing
- Main memory is getting larger and larger
- Manycore scalability becomes a serious issue in the modern OS design





Reference counter

• Number of accesses for a given object





Reference counter in the kernel object

File descriptor table
struct files_struct



File object
struct file



Directory entry struct dentry



read-write semaphore
struct rw_semaphore



Physical page frame struct page





Manycore scalability in kernel object access





Distributed reference counter

- Allocate local counter for each core
- Update operation : update the local counter
- Counter query : scan all local counters





Issues in distributed reference counter

- Memory pressure
 - Memory overhead increase in proportion to number of CPUs and objects
- Query latency
 - For reclaim the object, checking all local counter increase query latency
 - Overhead of obtaining the global state of the counter







Existing works for per-core reference counter



Design of Logical Distributed Counter



Characteristics of kernel objects



Per-core vs. Per-process view



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DOK UNIVER

(1939) 1939

KAIST

Per-process





Cause for counter contention



LODIC : Logical Distributed Counter

• LODIC

- Counter contention is caused by the contention among the processes
- Distributed counter with local counters are defined in per-process basis

- Used characteristics
 - Popularity : Define the counter with respect to the degree of sharing
 - Access brevity : Not consider the reference split





Objective



The number of counters are proportional to the degree of sharing



Key techniques

• File mapping

Map the file block to the process address space

• Reverse Mapping

- Do not use existing rmap mechanism that is not scalable
- Reverse Mapping based upon the process address space, file's address space
- Counter Embedding
 - Use the un-used bits in the page table entry



Allocate the counter with per-processes basis

How to allocate a counter of physical page on process address space?







Key technique 1: File Mapping





Find VMA of mapped page

How to quickly find the VMA of page mapped shared file block?





VMA

Key Technique 2: Reverse Mapping



Process-Space based reverse mapping





Key technique 3: Counter Embedding

- Embed the local counter at PTE
- For local counter, use un-used bits in PTE.











Evaluation



Throughput on shared file block read

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- 120-cores (15 cores/CPU, 8 socket, Intel Xeon E7-8870), 780 GB DRAM, Linux 4. 11. 6
- DRBH Workload on Fxmark

DOK UNIV

DKi

<u>만양</u> 1939

ΚΔΙΣΤ



B: Baseline , F: File-based reverse mapping , P: Process-based reverse mapping

Web server throughput

- 50 client processes, 50 server processes
- NGINX : Reverse proxy server that handles client request
- wrk benchmark : Make the client process to read request for the same file





Counter query latency

- fadvise() : System call to reclaim the page
- File size : 1GB

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- LODIC(10%): 10% of file blocks are mapped
- LODIC(20%) : 20% of file blocks are mapped
- LODIC(100%): 100% of file blocks are mapped

GINIA

FECH



Conclusion

• We take process-centric view in designing the distributed counting scheme

" Counter contention is caused by the contention among the processes, not by the contention on the processors "

- Number of local counters : With respect to the actual degree of sharing
- Memory pressure : Almost none
- Throughput on the shared block read increases by 65x
- Web server performance increases by 2.5x
- Memory pressure decreases by 13x against per-core distributed reference counters





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