# Determinism Is Not Enough: What Really Makes Multithreaded Programs Hard to Get Right and What Can Be Done about It?

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# One-slide summary

- Multithreaded programs: critical but hard to get right
- Many blamed *nondeterminism*
- Deterministic multithreading (DMT): one input → one schedule
- But, determinism is neither sufficient nor necessary for reliability!
- True culprit is rather quantitative: too many schedules
- Stable Multithreading (StableMT): all inputs → a small set of schedules

#### Background and motivation

# Multithreaded programs: hard to get right

- Plagued with concurrency bugs
  - Data races, atomicity violations, order violations, deadlocks, etc
- Concurrency bugs: bad
  - Have taken lives in the Therac 25 incidents and caused the 2003 Northeast blackout
  - May be exploited by attackers to violate confidentiality, integrity, and availability of critical systems [Hotpar 12]

#### Concurrency bug examples

Thread 0 mutex_lock(M) *obj = mutex_unlock(M)	Thread 1 mutex lock(M)	Thread 0 mutex lock(M)	Thread 1 mutex_lock(M) free(obj) mutex_unlock(M)
	free( <mark>obj</mark> )	*obj =	
	mutex_unlock(M)	mutex_unlock (M)	
Apache Bug #21287			
Thread 0	Thread 1	Thread 0	Thread 1
<pre>barrier_wait(B)</pre>	<pre>barrier_wait(B) result +=</pre>	<pre>barrier_wait(B) print(result)</pre>	<pre>barrier_wait(B)</pre>
<pre>print(result)</pre>		_	result +=
FFT in SDLASH2			



- *Schedule*: sequence of communication operations
- *Buggy schedule*: schedule that triggers concurrency bug

#### Challenges caused by nondeterminism



- Testing: less effective
- Debugging: more challenging

# Challenges caused by too many schedules



multithreading

- m threads, k lock() → more than (m!)^k schedules
- Even more schedules aggregated over all inputs
- Finding buggy schedules is like finding needles in a haystack

# Determinism: neither sufficient nor necessary for reliability

### Deterministic multithreading (DMT): one input → one schedule



- One testing execution validates all future executions on the same input
- Reproducing a concurrency bug requires only the input

### Determinism: not sufficient



Traditional multithreading

Deterministic multithreading

- Determinism is a narrow property
  - Same input + same program → same behavior
  - Input or program changes slightly? Unstable

#### Determinism: not necessary



multithreading

Deterministic multithreading

Stable multithreading

- Determinism is a binary property
  - Nondeterministic if one input  $\rightarrow$  n > 1 schedules
- Small n (e.g., 2) → challenges caused by nondeterminism are easy to solve

# Improving reliability with stable multithreading

# Are all exponentially many schedules necessary?

 Insight 1: for many programs, a wide range of inputs share the same equivalent class of schedules

• Insight 2: the overhead of enforcing a schedule on different inputs is low (e.g., 15%)

# Stable multithreading (StableMT): all inputs $\rightarrow$ a small set of schedules



- Vastly shrink the haystack → needles much easier to find
- Provide anticipated robustness and stability

# StableMT and DMT: orthogonal





Traditional multithreading

Deterministic multithreading



Stable deterministic multithreading

### Implementing and applying StableMT

- How can we compute the schedules to map inputs to?
  - Tern [OSDI 10]
- How can we enforce schedules deterministically and efficiently?
  - Peregrine [SOSP 11]
- How can we apply StableMT to effectively analyze multithreaded programs?
   – Schedule specialization [PLDI 12]

# Conclusion

- Determinism: neither sufficient nor necessary for reliability
- StableMT: map all inputs to vastly reduced set of schedules, greatly improving reliability

### Future work

- Systems level: more efficient, lightweight, scalable
- Application level: more applications
- Conceptual level: StableMT programming languages, models, and methods

# Applying StableMT to better analyze multithreaded programs

Analyzing multithreaded programs: hard



coverage (# of analyzed schedules / # of total schedules)

# Schedule Specialization [PLDI 12]

- Precision: Analyze the program over a small set of schedules.
- Coverage: Enforce these schedules at runtime.



coverage (# of analyzed schedules / # of total schedules)