

# How Sharp is SHARP?

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CARS@CSE-IITK

Secure hierarchy-aware cache replacement policy

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#### Mitigation for side-channel attacks

## Side Channel Attacks











Step 0: Attacker *fills* the entire shared cache (set)





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Step 1: Victim *evicts* cache blocks while running





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Step 0: Attacker *fills* the entire shared cache (set)

Step 1: Victim *evicts* cache blocks while running

Step 2: Attacker *probes* the cache set

If *misses* then victim has accessed the set





#### Cache Layout [HPCA '16]



Cache Layout [HPCA '16]

Fuzzing the timer [ISCA '12]



Cache Layout [HPCA '16]

Fuzzing the timer [ISCA '12]

Cache Addressing [MICRO '18]



Cache Layout [HPCA '16]

Fuzzing the timer [ISCA '12]

Cache Addressing [MICRO '18]

Cache replacement policy [ISCA '17]

Secure hierarchy-aware cache replacement policy

Mitigation for side-channel attacks

Secure hierarchy-aware cache replacement policy

Mitigation for side-channel attacks

Prevents cross-core back invalidation

#### Cross-core Back-Invalidation - I









#### **Cross-core Back-Invalidation - II**



#### **Cross-core Back-Invalidation - III**



LLC Miss





Shared Cache







Cache







Snarec



X





Shared Cache



X







Cache



X









Shared Cache







.





Ζ










Ζ

Stage-2







# How SHARP Works?



Shared Cache































# How SHARP Works?







Counter per core

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Increments on inter-core eviction

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For 1 billion cycles, the threshold value is 2000

Counter per core

Increments on inter-core eviction

For 1 billion cycles, the threshold value is 2000

On exceeding threshold, SHARP triggers OS interrupt

# Questions That We Ask?

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Does SHARP mitigate all attacks?








































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### **Questions That We Ask?**

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Does SHARP facilitate few more attacks?



#### Multi-threaded Attacker























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Does SHARP mitigate all attacks?

Does SHARP facilitate few more attacks?

Does threshold affect benign applications?





ChampSim, a trace driven simulator

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Simulated SHARP on a 16-core system with three levels of caches and huge pages

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Used different combinations of LLC thrashing and LLC fitting applications Example, 16:0 denotes 16 thrashing and zero fitting

# LLC Thrashing Benchmarks [SPEC CPU 2017]

Mix No	Thrashing Benchmarks
1	605.mcf-484B
2	605.mcf-665B
3	605.mcf-994B
4	607.cactubssn-2421B
5	620.omnetpp-141B
6	620.omnetpp-874B
7	621.wrf-6673B
8	623.xalancbmk-10B
9	649.fotonik-10881B
10	654.roms-523B

### Interbackhit Rate



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#### Interbackhit Counter



#### Interbackhit Counter



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#### Interbackhit Counter



## **Questions That We Ask?**

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### **Questions That We Ask?**

Does SHARP mitigate all attacks?

Does SHARP facilitate few more attacks?

Does threshold affect benign applications?

What does OS do when it receives an interrupt?






To deschedule

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To migrate to another socket

### Migration to Another Socket



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Causes performance overhead









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To migrate to an another socket

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causes slowdown 🙁

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causes significant slowdown 😕

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16-0 Mix, 100% apps got killed 😕

To deschedule

causes slowdown 🙁

To migrate to an another socket

causes significant slowdown 😕

To kill

16-0 Mix, 100% apps got killed 😕

Does mitigation strategy facilitates any new attack?

#### Threshold Aware Attack - I



#### **Threshold Aware Attack - II**



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causes slowdown 🙁

To migrate to another socket

causes significant slowdown 🙁

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16-0 Mix, 100% apps got killed 🙁

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Is SHARP secure in terms of information leakage?







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#### Conclusion

SHARP is not that sharp

Facilitates new attacks

Don't mitigates all attacks

Role of OS is not defined

Performance overhead to benign applications

# Thank You!



#### Semiconductor Research Corporation

