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HiveD:

Sharing a GPU Cluster for Deep Learning with Guarantees

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Multi-Tenant GPU Clusters Today

- Shared by multiple tenants
 - Built with budgets/hardware from tenants
- **Reservation** is necessary for *guaranteed* resource availability and user experiences
- Reserve *number of GPUs (quota)* for each tenant
 - A tenant expects to access at least its contributed share



Affinity Matters, but NOT Reserved

- Jobs usually have affinity requirements
- Quota reservation + global affinity optimization (defragmentation)



They are all equivalent in quota usage!

Anomaly: Sharing Leads to Suffering!

- "External" fragmentation across tenants
 - Fragmentation from *other* tenants, even everyone is within the quota
 - Quota cannot *isolate* fragmentation across tenants
- External fragmentation makes sharing harmful to tenants
 - Worse performance in the *shared* cluster than in *private* clusters
 - Global defrag might sometimes hurt job performance (a complex multi-objective optimization)
- Real users are reverting to private clusters!



HiveD

- Primary goal: Sharing Safety
 - Any sequence of GPU requests (possibly with affinity constraints) can be satisfied on the shared cluster if satisfied on the tenant's private cluster

- Key idea

- Reserve GPU affinity explicitly
- Separate the concern of sharing safety from other scheduling goals

Two-Layer Architecture

- Virtual Private Cluster
 - New resource abstraction: cell, captures quota and affinity
 - Compatibility with SOTA deep learning schedulers in VCs
- From Virtual to Physical
 - Dynamic cell allocation with proven sharing safety
 - Natural support for low-priority jobs



Separation of Concerns

Sharing Safety	\bigcirc
Scheduling Efficiency	/ 🕢
Utilization	\bigcirc

Two-Layer Architecture

- Virtual Private Cluster
 - New resource abstraction: cell, captures quota and affinity
 - Compatibility with SOTA deep learning schedulers in VCs
- From Virtual to Physical
 - Dynamic cell allocation with proven sharing safety
 - Natural support for low-priority jobs
- Best of both worlds
 - A private cluster: guaranteed resource availability regardless of other tenants
 - A shared cluster: more resources when available

Separation of Concerns

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- A cell is a set of GPUs at a certain level of affinity



A cell can be split into multiple equivalent **buddy cells**

Virtual Private Clusters

- Cells at each level, modeling a tenant's private cluster
- Dynamic cell binding
 - Reducing preemptions and fragmentation, handling faulty hardware
 - Handled by Buddy Cell Allocation algorithm



Dynamic binding via Buddy Cell Allocation

Tenant A	Cell Level	Α	В	С
······································	L4 cell (8-GPU)	0	0	2
Tenant B	L3 cell (4-GPU)	1	1	0
Turne Ci	L2 cell (2-GPU)	1	1	1
Tenant C	L1 cell (1-GPU)	1	1	0

VC Assignment

- For a cell request at level-k:
 - Allocate a free level-k cell if any
 - Split a free level-(k+1) cell otherwise



Cell request for level-4 (node)

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- For a cell request at level-k:
 - Allocate a free level-k cell if any
 - Split a free level-(k+1) cell otherwise
- Cell release (and merge) works oppositely
- Keep as many higher-level cells as possible
- Proven safety guarantee
 - Satisfies any cell request within a VC, if the initial VC assignment is feasible

Algorithm 1 Buddy Cell Allocation Algorithm

- 1: // Initial state of free_cells: only top level has cells
- 2: **procedure** ALLOCATECELL(cell_level)
- 3: **if** free_cells[cell_level].size() == 0 **then**
- 4: $c = AllocateCell(cell_level+1)$
- 5: cells = Split(c) \triangleright Split cells are buddies
- 6: free_cells[cell_level].extend(cells)
- 7: **Return** free_cells[cell_level].pop()
- 8:
- 9: procedure RELEASECELL(cell)
- 10: **if** cell.buddies \subseteq free_cells[cell.level] **then**
- 11: higher_cell = Merge(cell, cell.buddies)
- 12: free_cells[cell.level].remove(cell.buddies)
- 13: ReleaseCell(higher_cell)
- 14: **else**
- 15: free_cells[cell.level].add(cell)

- Two cell views
 - High-priority guaranteed jobs with VC safety
 - Low-priority opportunistic jobs to improve utilization





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Open-Source Implementation

https://github.com/microsoft/hivedscheduler

- Implemented on Kubernetes



- Integrated with Microsoft OpenPAI

https://github.com/microsoft/pai

- Deployed at Microsoft for 12+ months
 - Managing 1000+ heterogeneous GPUs
 - Serving research and production workloads at scale
- More implementation details and operation experiences in the paper

Evaluation: 96-GPU Cluster Experiment

- Schedulers
 - YARN-CS (Philly) [ATC 19]
 - Gandiva [OSDI 18]
 - Tiresias [NSDI 19]
 - A VC preserves the precise affinity structure, making STOA schedulers applicable
- HiveD achieves the best of both worlds
 - Sharing anomalies identified in all schedulers with quota; HiveD eliminates them all!
 - Significantly shorter queuing delay than in the private cluster
 - Similar overall JCT compared to applying the schedulers globally



Evaluation: Trace-Driven Simulation

- 2-month trace from a 2232-GPU cluster with 11 tenants
- More sharing anomalies under high load
 - Up to 8,000+ minutes of excessive queuing delay
 - 7x on average
 - Again HiveD eliminates them all



Evaluation: Trace-Driven Simulation

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 - Up to 8,000+ minutes of excessive queuing delay
 - 7x on average
 - Again HiveD eliminates them all
- Sharing anomaly leads to diminishing benefits of sharing
 - Decommission a tenant with higher avg. queuing delay in the shared cluster with quota than in the private cluster
 - Tenants owning 37% quota (two large tenants) decommissioned!
 - Significantly longer queuing delay of the other 9 tenants in this smaller cluster

Evaluation: Buddy Cell Allocation

- Reducing preemptions
 - Avoid dirty cells with dynamic binding
- Reducing fragmentation
 - Pack cells across VCs



Conclusion

- HiveD addresses the challenge of sharing a GPU cluster with
 - Sharing safety: simple and practical guarantee easily appreciated by tenants
 - *Cell:* new resource abstraction for defining tenants' affinity structures
 - Buddy cell allocation: proven safety and support for low-priority jobs
 - *Two-layer architecture* to incorporate other scheduling goals while guaranteeing sharing safety

Thank you!

Contact

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Code released at https://github.com/microsoft/hivedscheduler