

FlightTracker

Consistency across Read-Optimized Online Stores at Facebook

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The challenges of providing Read-Your-Writes (RYW) consistency for the social graph

Agenda

Our solution: FlightTracker

Lessons learned and production experiences



Q&A

TAO Read-optimized data store for the social graph



User devices

[Usenix ATC'13] TAO: Facebook's Distributed Data Store for the Social Graph

Stateless web servers

Application logic

Graph store TAO



Social graph consistency model

 Applications can query latest data *if necessary*

• Reading fresher data is OK i.e., per-item at-or-after / lower bound semantics

• End users get Read-Your-Writes

• Eventual consistency as a baseline

TAO 2013: two-layer write-through cache



[Usenix ATC'13]



TAO 2013: fixed communication patterns for RYW





Evolution since 2013



















Scalability: add global indexes and other data stores





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Rethinking social graph consistency



ENABLE

Dynamic communication paths

Extend uniform semantics to global indexes and new Database types



Our solution: FlightTracker

Our solution: decompose the consistency problem

Consistency: what writes are visible to a read?



FlightTracker Identify missing writes

- Data-store agnostic
- Reusable and extensible
- Write metadata only



Ticket Encapsulated set of writes





Ticket-inclusive reads

Ensure visibility: read results reflect missing writes

- Data-store specific strategies
- Ticket attached on each query

RYW: User writes span web requests









L1 cache



FlightTracker maps user_ids to recent write metadata

RYW: Read flow using FlightTracker



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RYW: Write flow using FlightTracker





Ticket

• Write set: metadata that identifies a set of writes Joinable, i.e., set union Encapsulated Most code paths treat Tickets as opaque tokens Serialized and compressed on the wire

about its semantics



- Named "Ticket" (vs. timestamp / version) to reduce potential preconception



Ticket representation

```
Ticket {
  RepForDatabaseA databaseA;
  RepForDatabaseB databaseB;
  • • •
  Timestamp globalTs;
   Example database-specific representation
| |
RepForDatabaseA {
 map<WriteKey, pair<Version, Timestamp>> perKeyMap;
 map<ShardId, pair<TxnId, Timestamp>> perShardMap;
```





Ticket-inclusive read

Data-store specific implementation strategies

Fix data store first e.g., consistency miss for caches





Fix stale results

e.g., client read repair for indexes

Reevaluate query e.g., on a diff replica; at a later time



Ticket-inclusive read for caches







Challenges for global indexes









Beyond RYW

Beyond RYW: additional FlightTracker session types

- The default session is an end user, which is sticky to a region.
- Select applications need write visibility guarantees other than user-centric RYW.
- Flexible definition of "session"
 - E.g., async job, particular TAO object (see paper)
 - Reads and writes can belong to multiple sessions.
- Customizable FlightTracker quorum config
 - E.g., write to FlightTracker in all regions, read locally





Beyond RYW: external Ticket handling

- Systems at the product infrastructure layer may handle Tickets explicitly
 - Especially when we can piggyback on existing communication
 - Still hidden from applications



Example: pub-sub notification system



Pub-sub notification system: the problem







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Lessons learned production experiences

Ticket internals are encapsulated from applications.

Can safely include additional write metadata while honoring RYW. e.g., FlightTracker server or client are free to join Tickets whenever new writes happen.

Ticket-inclusive reads only targets per-item at-or-after / lower-bound semantics.





Ticket compaction for writes older than 60s

e.g., can replace write metadata with a single global timestamp



Can safely include additional write metadata while honoring RYW. e.g., joining Tickets









Can safely include additional write metadata while honoring RYW. e.g., joining Tickets



Single-round protocol for FlightTracker Only need to provide **durability** but NOT atomicity



Lessons learned (Cont'd)





Identifying logged-in user_id was more difficult than we expected.

Constraints on FlightTracker design are not based on the average case, but the extreme ones, such as hot spots or disaster scenarios.





The ability to opt into alternative write visibility guarantees late in product dev cycle enabled us to make RYW a good default.



Lessons learned (Cont'd)



The applications that cause the most operational trouble often need RYW the least.

Ticket-inclusive reads established a contract that revealed latent bugs in our existing eventual consistency protocols.





The decomposition in the FlightTracker design allowed us to incrementally provide RYW for 2 caches, 3 global indexes, and 2 database technologies.



It's real and it works

99.9999%

FlightTracker read availability measured from the client



FlightTracker write availability compared with underlying data stores



CPU/RAM overhead on existing data stores and web servers



It's real and it works

4 Yrs

In production

FlightTracker write QPS

20M

10^{15}

FlightTracker read QPS

Social graph queries per day





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