

Torturing Databases for Fun and Profit

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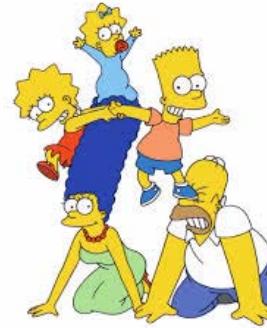
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[‡] *HP Labs*







database



- ACID: atomicity, consistency, isolation, and durability
- even under failures

List of databases survived

Everything is broken under simulated power faults

Database	File System	Workload				
		W-1	W-2	W-3	W-4.1	W-4.2
TokyoCabinet	ext3	D	D	D	ACD	ACD
	XFS	--	D	D	ACD	D
MariaDB	ext3	D	D	D	D	D
	XFS	D	D	D	D	D
LightningDB	ext3	--	--	--	--	--
	XFS	--	--	--	--	--
SQLite	ext3	D	D	--	D	D
	XFS	--	--	D	D	D
KVS-A	ext3	--	--	Hang	--	--
	XFS	--	--	--	--	--
SQL-A	ext3	D	D	D	D	D
	XFS	D	D	D	D	D
SQL-B	ext3	D	D	CD	CD	CD
	XFS	CD	D	CD	CD	CD
SQL-C	NTFS	D	D	D	D	D

**Power faults cannot happen
nowadays, right?**





2013: "... **POWER OUTAGE** during Super Bowl ... because a **RELAY DEVICE MALFUNCTIONED**."



2014: "Internap Data Center **OUTAGE** Takes Down Livestream, StackExchange"



2014: "Data Center **FIRE** Leads to **OUTAGE** ..."



2014: "... **ELECTRICAL FIRE** took down ... primary data center... **ALL POWER** was **OFF** ..."



2013: "**POWER OUTAGE** knocks DreamHost customers offline ..."



2013: "A data center **POWER OUTAGE** is being blamed for ...Visa downtime ..."



2013: "A **POWER OUTAGE** at a key New Jersey data center ..."



2013: "... **POWER OUTAGE** during Super Bowl ... because a **RELAY DEVICE MALFUNCTIONED**."



2012: "... **HUNAM ERROR** was responsible for a data center **POWER OUTAGE** ..."



2012: "**POWER OUTAGE** Hits London Data Center ..."



2012: "Amazon Data Center **LOSES POWER** During **STORM** ..."



2011: "Colocation provider Colo4 experienced a **POWER OUTAGE** ..."



2010: "About 3,000 servers at Montreal web host iWeb experienced an **OUTAGE** ..."

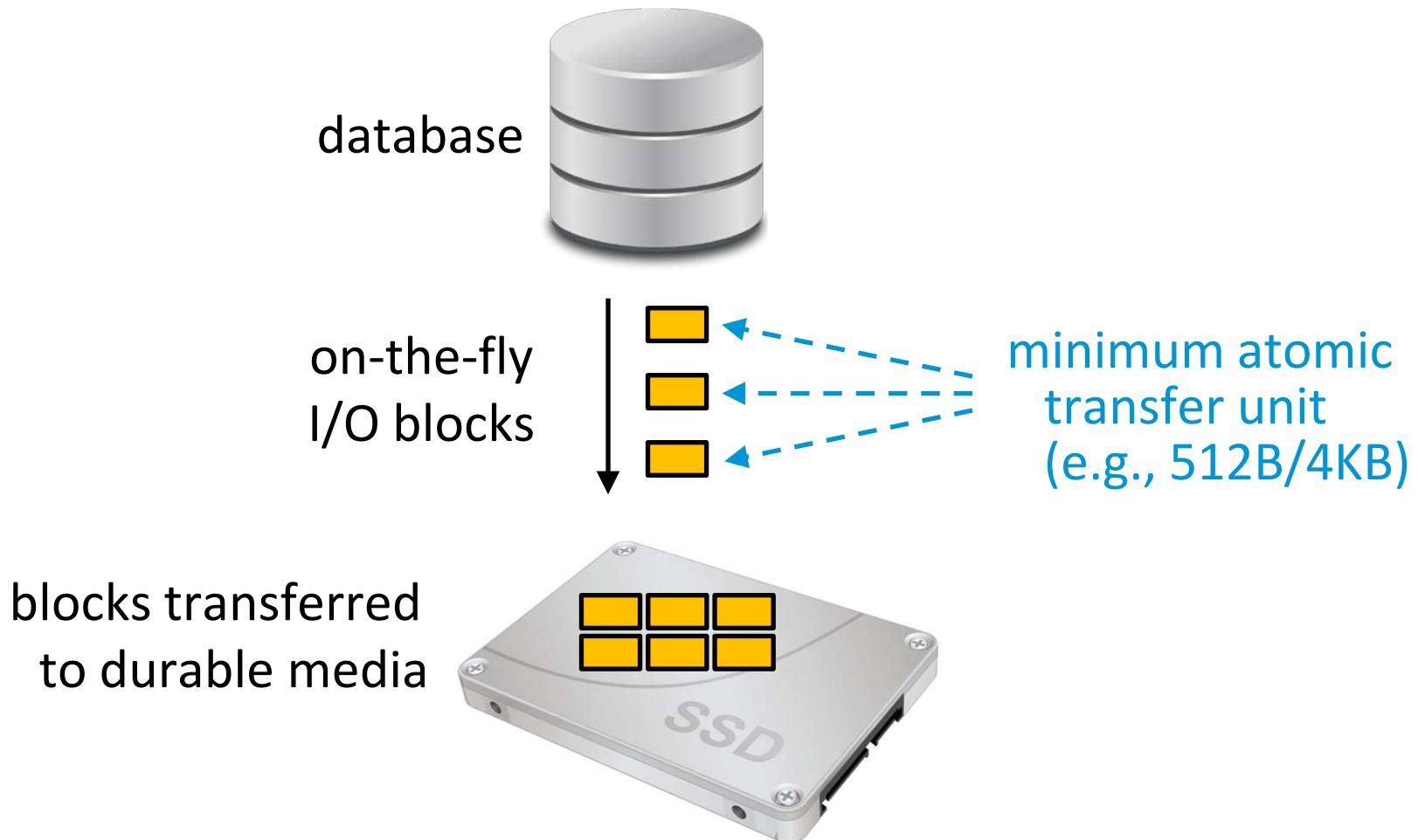


2010: "**CAR CRASH** Triggers Amazon **POWER OUTAGE** ..."

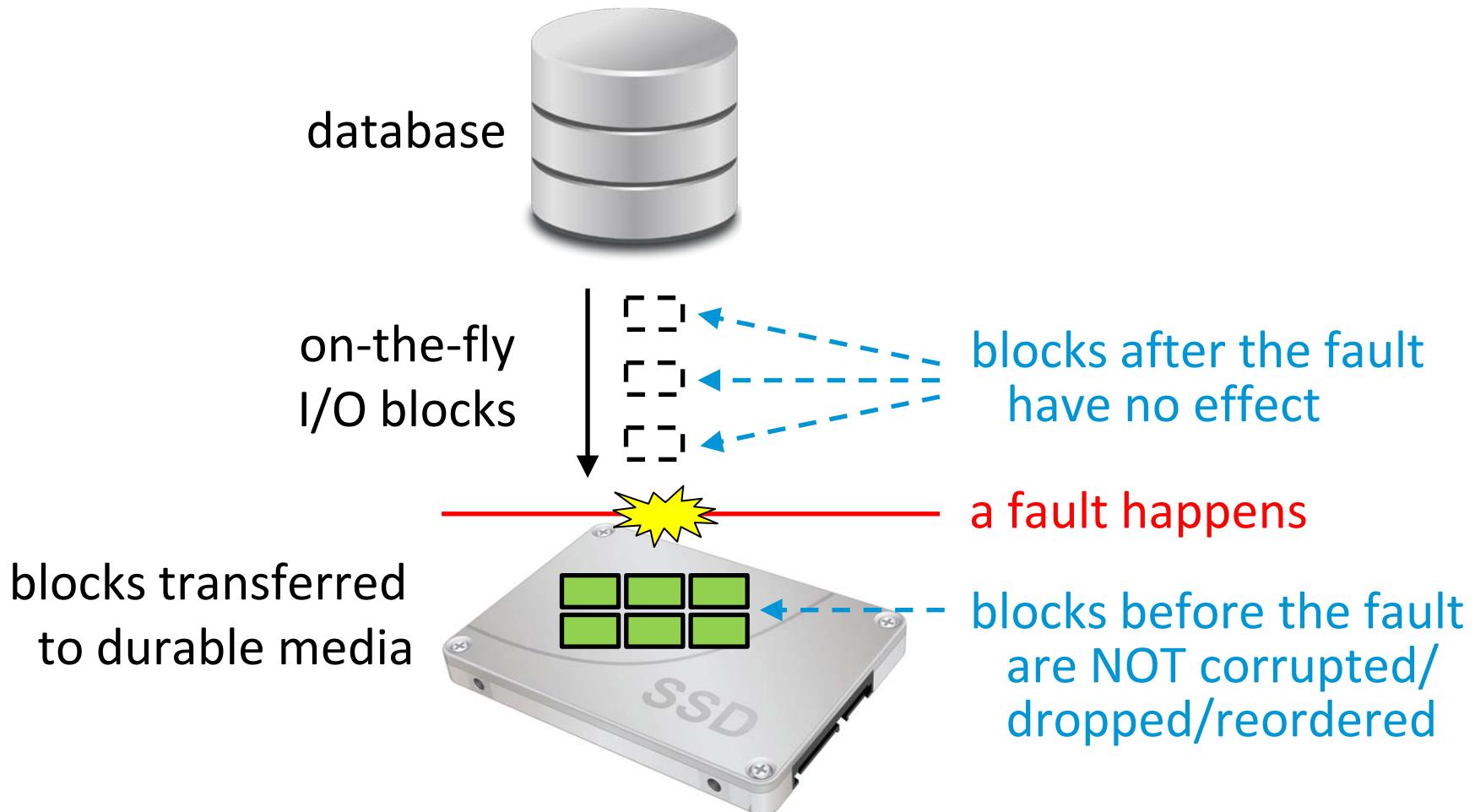
Database Torture 101



Fault Model: Clean termination of I/O stream



Fault Model: Clean termination of I/O stream



Why not introduce corruption/dropping/reordering?

- Unreasonable to require databases to handle arbitrary bad behavior introduced in the lower layers
- Simulated bad behavior w/o verification by real failures may be unrealistic
 - I/O path in kernel & device puts constraints on failure states



How do we test DBs on multiple OSes w/ high fidelity?

database

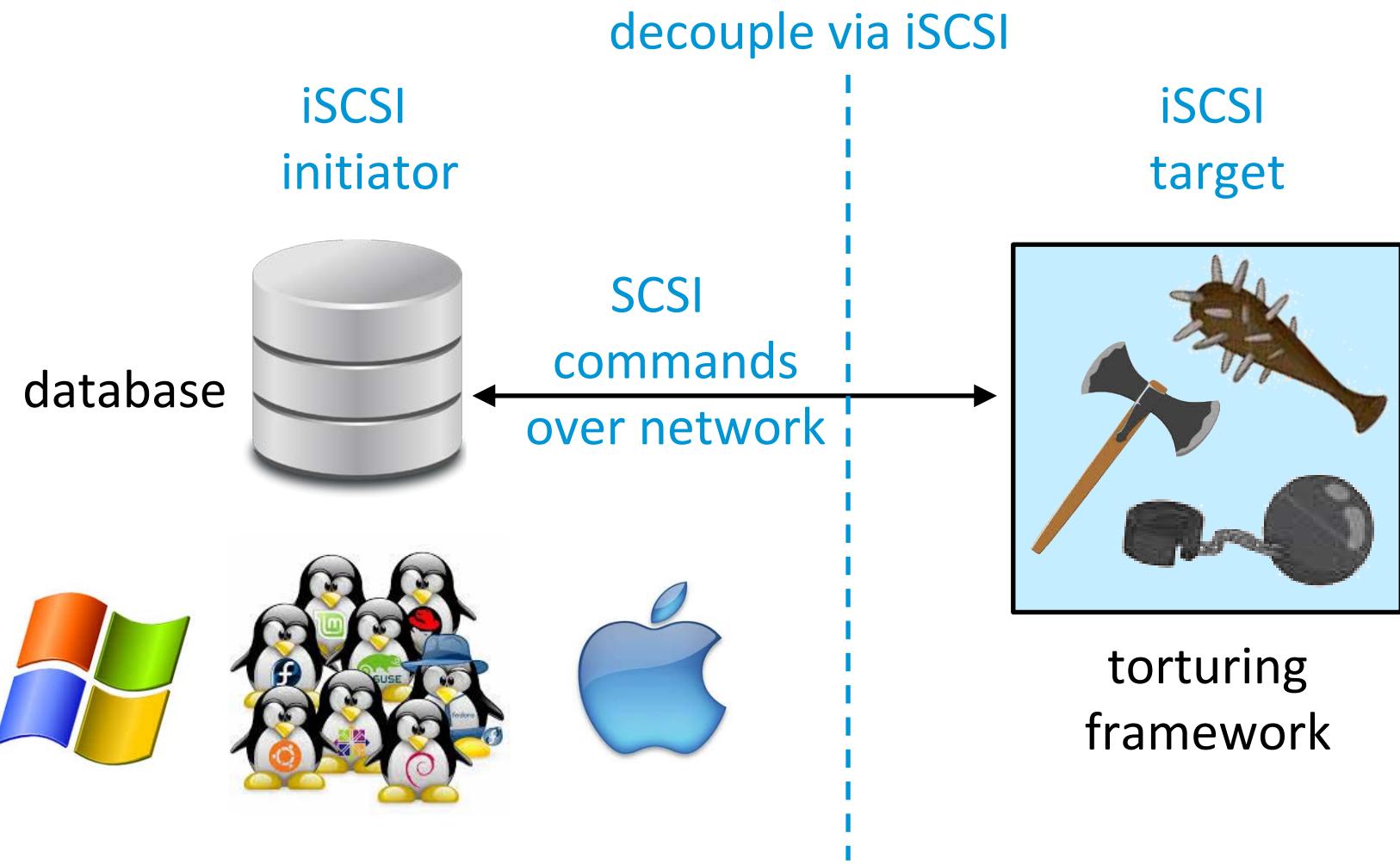


No disturbance on
thread scheduling

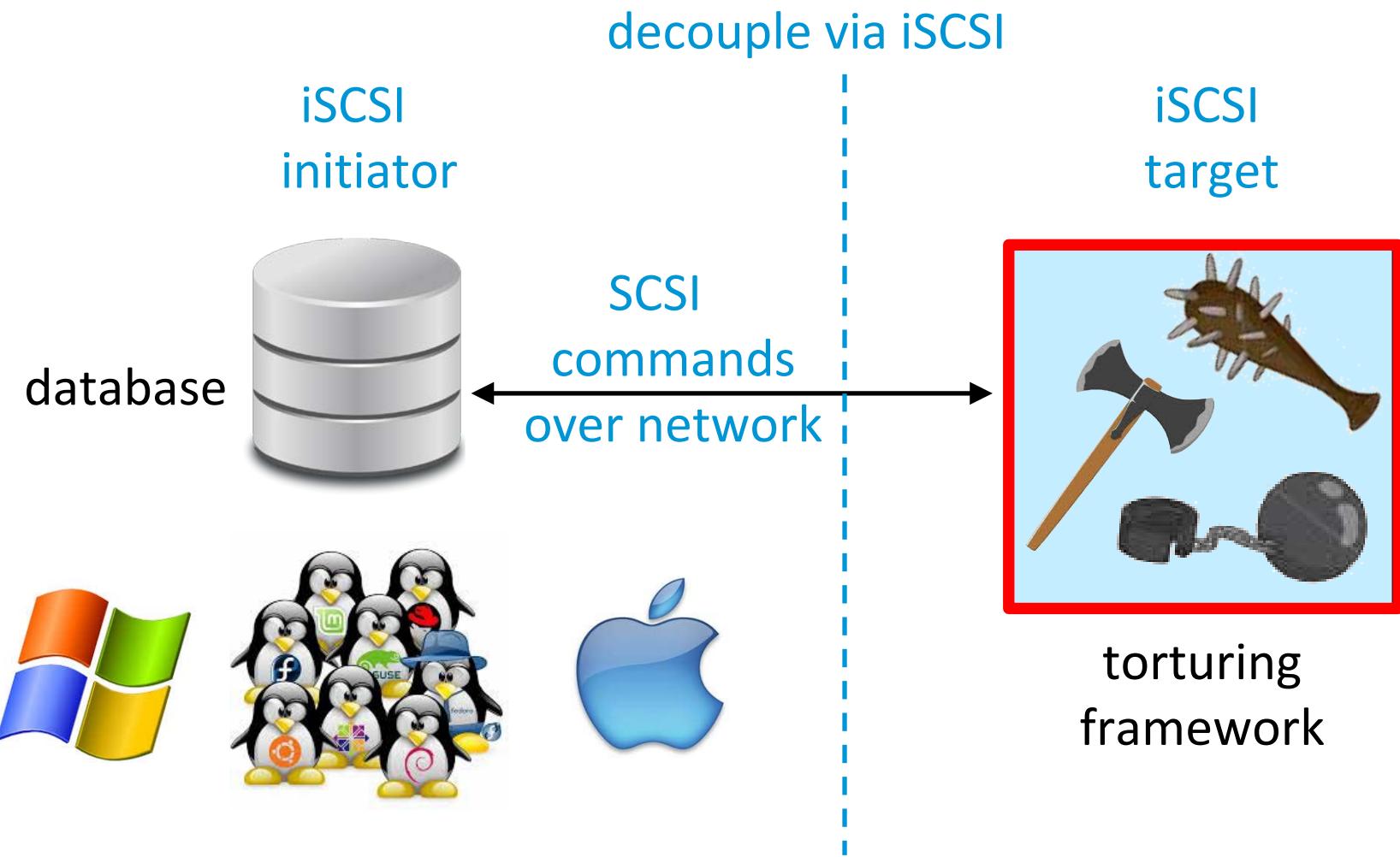
No disturbance on
interactions among
DB, memory manager,
FS, volume manager,
I/O scheduler, ...



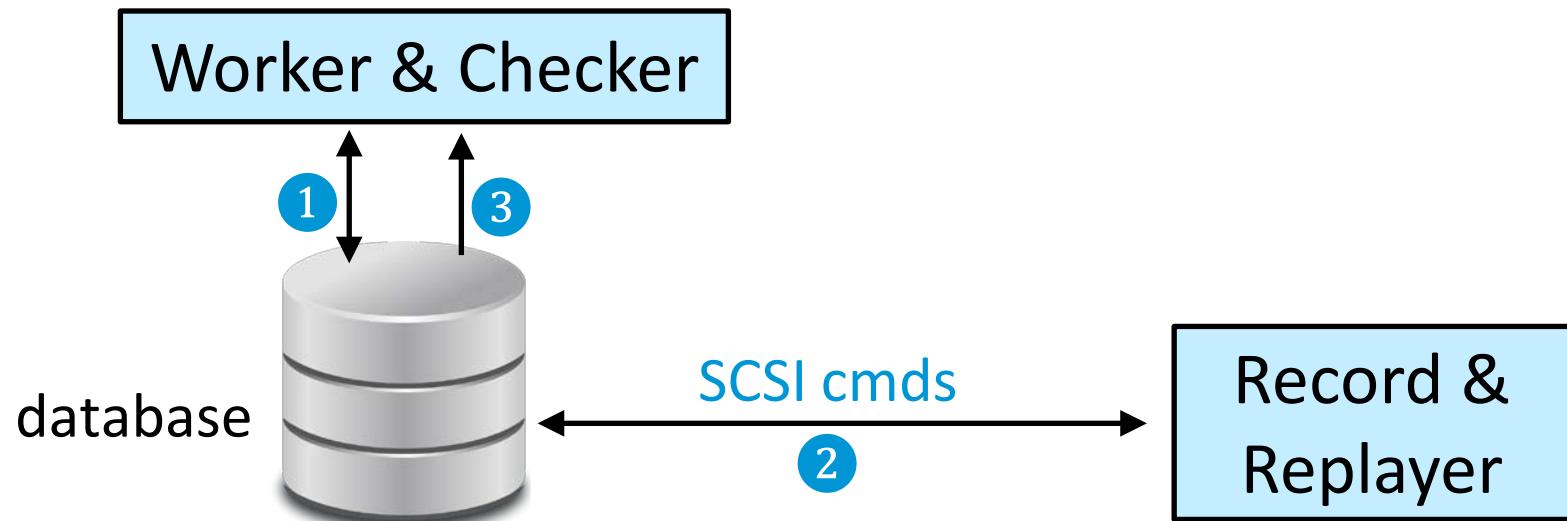
How do we test DBs on multiple OSes w/ high fidelity?



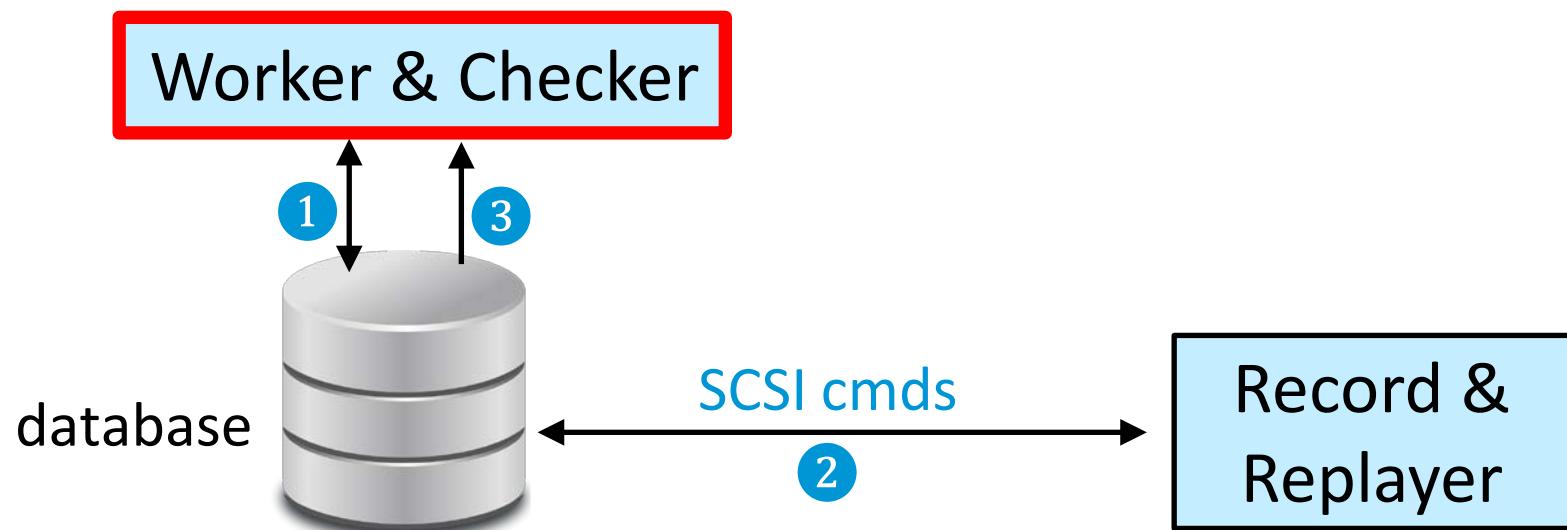
How do we test DBs on multiple OSes w/ high fidelity?



Framework Overview



Framework Overview



Workload Example

	key	value
meta rows	THR-1-TXN-1	v-init-THR-1-TXN-1
	THR-1-TXN-2	v-init-THR-1-TXN-2
	THR-2-TXN-1	v-init-THR-2-TXN-1
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-init-2
	k-3	v-init-3
	k-4	v-init-4
	k-5	v-init-5
	k-6	v-init-6
	k-7	v-init-7
	k-8	v-init-8

DB table

Workload Example

two parts

meta rows

work rows

	key	value
meta rows	THR-1-TXN-1	v-init-THR-1-TXN-1
	THR-1-TXN-2	v-init-THR-1-TXN-2
	THR-2-TXN-1	v-init-THR-2-TXN-1
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-init-2
	k-3	v-init-3
	k-4	v-init-4
	k-5	v-init-5
	k-6	v-init-6
	k-7	v-init-7
	k-8	v-init-8

DB table

2 threads,
2 transactions
per thread

Has known initial state

	key	value
meta rows	THR-1-TXN-1	v-init-THR-1-TXN-1
	THR-1-TXN-2	v-init-THR-1-TXN-2
	THR-2-TXN-1	v-init-THR-2-TXN-1
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-init-2
	k-3	v-init-3
	k-4	v-init-4
	k-5	v-init-5
	k-6	v-init-6
	k-7	v-init-7
	k-8	v-init-8

DB table

Each transaction updates N random work rows + 1 meta row

	key	value
meta rows	THR-1-TXN-1	v-init-THR-1-TXN-1
	THR-1-TXN-2	v-init-THR-1-TXN-2
	THR-2-TXN-1	v-init-THR-2-TXN-1
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-init-2
	k-3	v-init-3
	k-4	v-init-4
	k-5	v-init-5
	k-6	v-init-6
	k-7	v-init-7
	k-8	v-init-8

DB table

Each transaction updates N random work rows + 1 meta row

	key	value
meta rows	THR-1-TXN-1	v-init-THR-1-TXN-1
	THR-1-TXN-2	v-init-THR-1-TXN-2
	THR-2-TXN-1	v-init-THR-2-TXN-1
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-init-2
	k-3	v-init-3
	k-4	v-init-4
	k-5	v-init-5
	k-6	v-init-6
	k-7	v-init-7
	k-8	v-init-8

THR-1-TXN-1

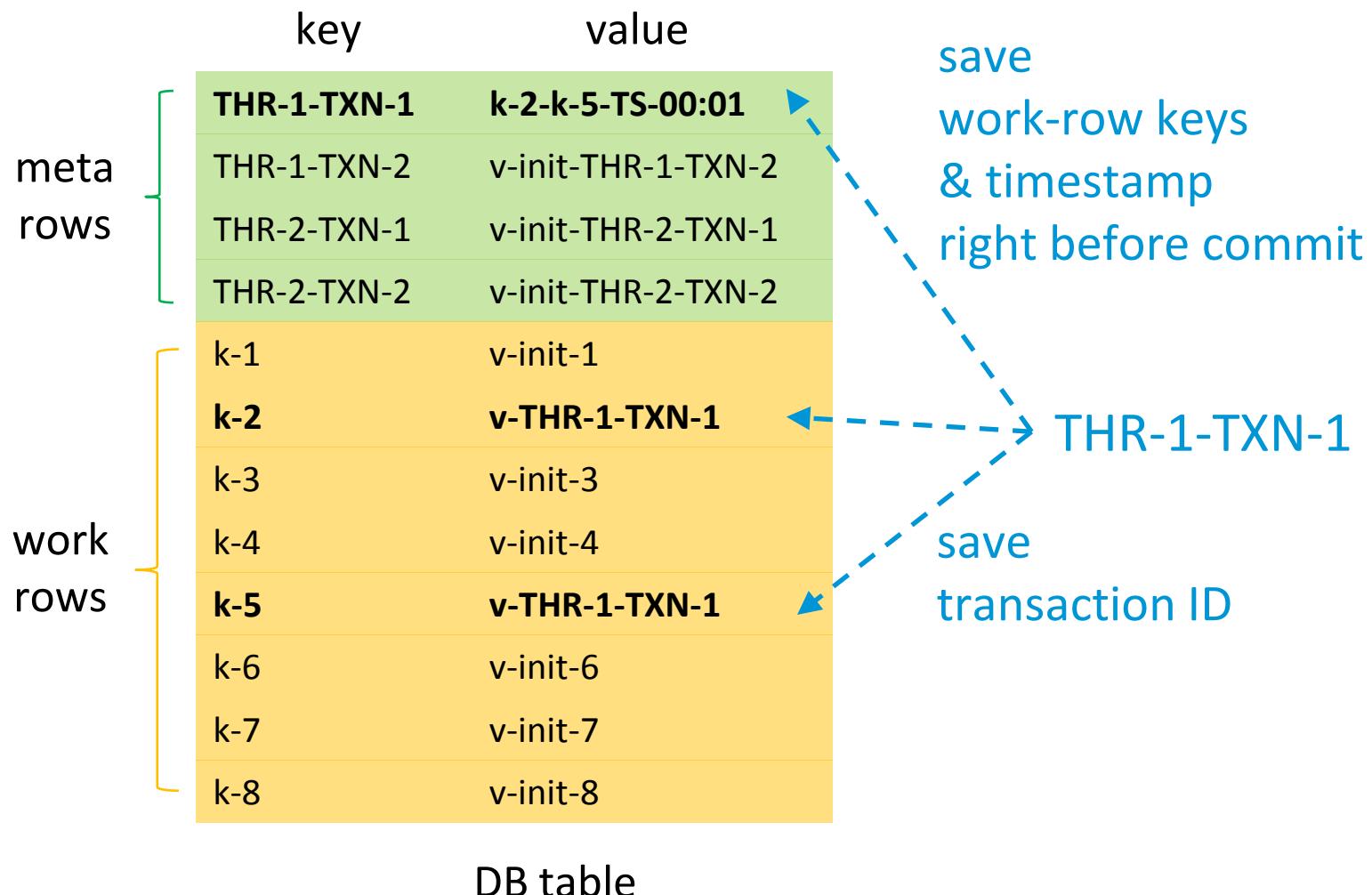
DB table

Each transaction updates N random work rows + 1 meta row

	key	value	
meta rows	THR-1-TXN-1	v-init-THR-1-TXN-1	
	THR-1-TXN-2	v-init-THR-1-TXN-2	
	THR-2-TXN-1	v-init-THR-2-TXN-1	
	THR-2-TXN-2	v-init-THR-2-TXN-2	
work rows	k-1	v-init-1	
	k-2	v-THR-1-TXN-1	THR-1-TXN-1
	k-3	v-init-3	
	k-4	v-init-4	
	k-5	v-THR-1-TXN-1	save transaction ID
	k-6	v-init-6	
	k-7	v-init-7	
	k-8	v-init-8	

DB table

Each transaction updates N random work rows + 1 meta row



Fully exercise concurrency control

	key	value
meta rows	THR-1-TXN-1	k-2-k-5-TS-00:01
	THR-1-TXN-2	v-init-THR-1-TXN-2
	THR-2-TXN-1	v-init-THR-2-TXN-1
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-THR-1-TXN-1
	k-3	v-init-3
	k-4	v-init-4
	k-5	v-THR-1-TXN-1
	k-6	v-init-6
	k-7	v-init-7
	k-8	v-init-8

DB table

The diagram illustrates a database table structure with two main sections: 'meta rows' and 'work rows'. The 'meta rows' section contains four entries, each consisting of a key and a value. The 'work rows' section contains eight entries, also consisting of a key and a value. A transaction, labeled 'THR-1-TXN-1' in blue, is shown interacting with the table. Three dashed blue arrows originate from the transaction and point to specific rows: one arrow points to the first meta row (key 'THR-1-TXN-1', value 'k-2-k-5-TS-00:01'), another arrow points to the second meta row (key 'THR-1-TXN-2', value 'v-init-THR-1-TXN-2'), and a third arrow points to the fifth work row (key 'k-5', value 'v-THR-1-TXN-1').

Fully exercise concurrency control

	key	value
meta rows	THR-1-TXN-1	k-2-k-5-TS-00:01
	THR-1-TXN-2	v-init-THR-1-TXN-2
	THR-2-TXN-1	k-7-k-6-TS-00:03
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-THR-1-TXN-1
	k-3	v-init-3
	k-4	v-init-4
	k-5	v-THR-1-TXN-1
	k-6	v-THR-2-TXN-1
	k-7	v-THR-2-TXN-1
	k-8	v-init-8

DB table

THR-2-TXN-1

Fully exercise concurrency control

	key	value
meta rows	THR-1-TXN-1	k-2-k-5-TS-00:01
	THR-1-TXN-2	k-6-k-8-TS-00:13
	THR-2-TXN-1	k-7-k-6-TS-00:03
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-THR-1-TXN-1
	k-3	v-init-3
	k-4	v-init-4
	k-5	v-THR-1-TXN-1
	k-6	v-THR-1-TXN-2
	k-7	v-THR-2-TXN-1
	k-8	v-THR-1-TXN-2

DB table

THR-1-TXN-2

Fully exercise concurrency control

	key	value
meta rows	THR-1-TXN-1	k-2-k-5-TS-00:01
	THR-1-TXN-2	k-6-k-8-TS-00:13
	THR-2-TXN-1	k-7-k-6-TS-00:03
	THR-2-TXN-2	k-3-k-7-TS-00:14
work rows	k-1	v-init-1
	k-2	v-THR-1-TXN-1
	k-3	v-THR-2-TXN-2
	k-4	v-init-4
	k-5	v-THR-1-TXN-1
	k-6	v-THR-1-TXN-2
	k-7	v-THR-2-TXN-2
	k-8	v-THR-1-TXN-2

DB table

THR-2-TXN-2

**A power fault just happened during
our workload ...**



Is there any ACID violation after recovery?

	key	value
meta rows	THR-1-TXN-1	k-2-k-5-TS-00:01
	THR-1-TXN-2	k-6-k-8-TS-00:13
	THR-2-TXN-1	k-7-k-6-TS-00:03
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-THR-1-TXN-1
	k-3	v-THR-2-TXN-2
	k-4	v-init-4
	k-5	v-THR-1-TXN-1
	k-6	v-THR-1-TXN-2
	k-7	v-THR-2-TXN-2
	k-8	v-THR-1-TXN-2

recovered DB table

Is there any ACID violation after recovery?

	key	value
meta rows	THR-1-TXN-1	k-2-k-5-TS-00:01
	THR-1-TXN-2	k-6-k-8-TS-00:13
	THR-2-TXN-1	k-7-k-6-TS-00:03
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-THR-1-TXN-1
	k-3	v-THR-2-TXN-2
	k-4	v-init-4
	k-5	v-THR-1-TXN-1
	k-6	v-THR-1-TXN-2
	k-7	v-THR-2-TXN-2
	k-8	v-THR-1-TXN-2

recovered DB table

Is there any ACID violation after recovery?

	key	value
meta rows	THR-1-TXN-1	k-2-k-5-TS-00:01
	THR-1-TXN-2	k-6-k-8-TS-00:13
	THR-2-TXN-1	k-7-k-6-TS-00:03
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-THR-1-TXN-1
	k-3	v-THR-2-TXN-2
	k-4	v-init-4
	k-5	v-THR-1-TXN-1
	k-6	v-THR-1-TXN-2
	k-7	v-THR-2-TXN-2
	k-8	v-THR-1-TXN-2

recovered DB table

Is there any ACID violation after recovery?

	key	value
meta rows	THR-1-TXN-1	k-2-k-5-TS-00:01
	THR-1-TXN-2	k-6-k-8-TS-00:13
	THR-2-TXN-1	k-7-k-6-TS-00:03
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-THR-1-TXN-1
	k-3	v-THR-2-TXN-2
	k-4	v-init-4
	k-5	v-THR-1-TXN-1
	k-6	v-THR-1-TXN-2
	k-7	v-THR-2-TXN-2
	k-8	v-THR-1-TXN-2

recovered DB table

Is there any ACID violation after recovery?

	key	value
meta rows	THR-1-TXN-1	k-2-k-5-TS-00:01
	THR-1-TXN-2	k-6-k-8-TS-00:13
	THR-2-TXN-1	k-7-k-6-TS-00:03
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-THR-1-TXN-1
	k-3	v-THR-2-TXN-2
	k-4	v-init-4
	k-5	v-THR-1-TXN-1
	k-6	v-THR-1-TXN-2
	k-7	v-THR-2-TXN-2
	k-8	v-THR-1-TXN-2

recovered DB table

Atomicity violation!
should have been
updated w/ work rows

Is there any ACID violation after recovery?

	key	value
meta rows	THR-1-TXN-1	k-2-k-5- TS-00:01
	THR-1-TXN-2	k-6-k-8- TS-00:13
	THR-2-TXN-1	k-7-k-6- TS-00:03
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-THR-1-TXN-1
	k-3	v-THR-2-TXN-2
	k-4	v-init-4
	k-5	v-THR-1-TXN-1
	k-6	v-THR-1-TXN-2
	k-7	v-THR-2-TXN-2
	k-8	v-THR-1-TXN-2

recovered DB table

allow checking time & order related properties

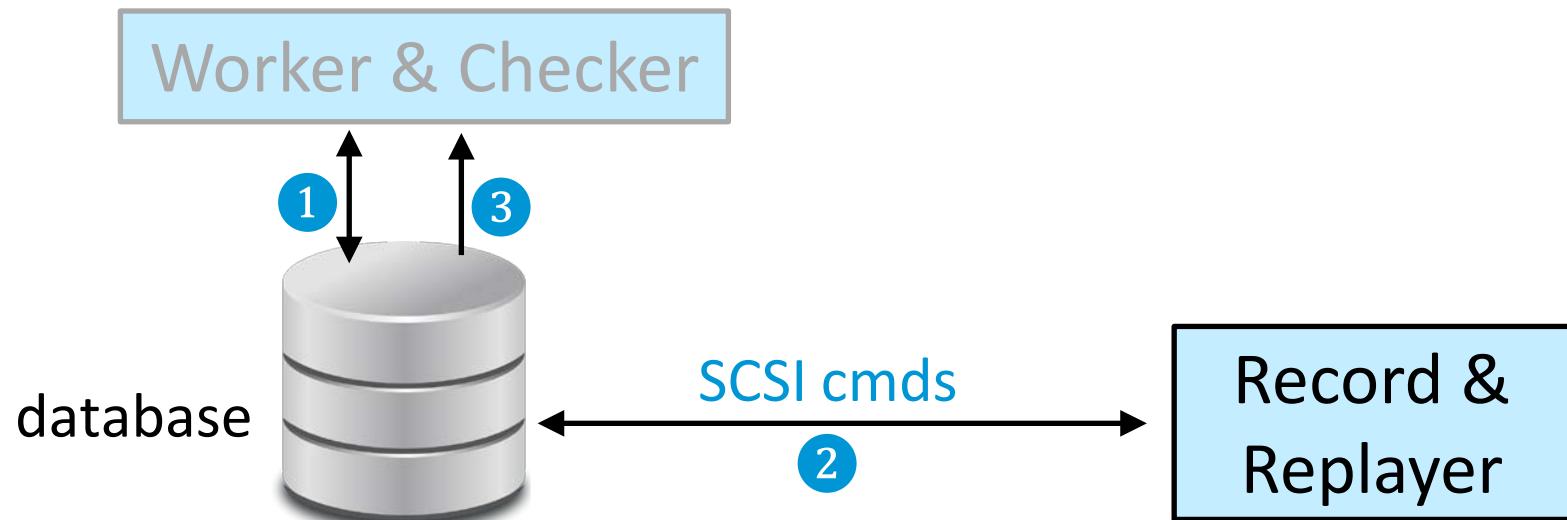
Is there any ACID violation after recovery?

	key	value
meta rows	THR-1-TXN-1	k-2-k-5-TS-00:01
	THR-1-TXN-2	k-6-k-8-TS-00:13
	THR-2-TXN-1	k-7-k-6-TS-00:03
	THR-2-TXN-2	v-init-THR-2-TXN-2
work rows	k-1	v-init-1
	k-2	v-THR-1-TXN-1
	k-3	v-THR-2-TXN-2
	k-4	v-init-4
	k-5	v-THR-1-TXN-1
	k-6	v-THR-1-TXN-2
	k-7	v-THR-2-TXN-2
	k-8	v-THR-1-TXN-2

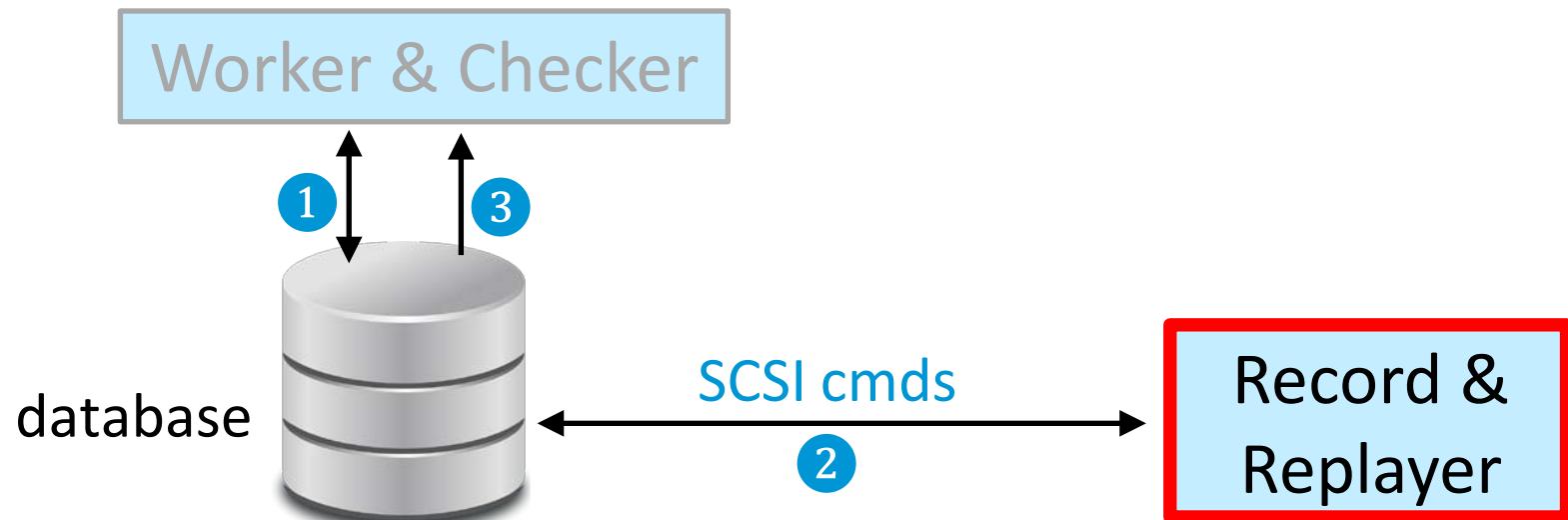
recovered DB table

More workloads
& ACID checking
in the paper

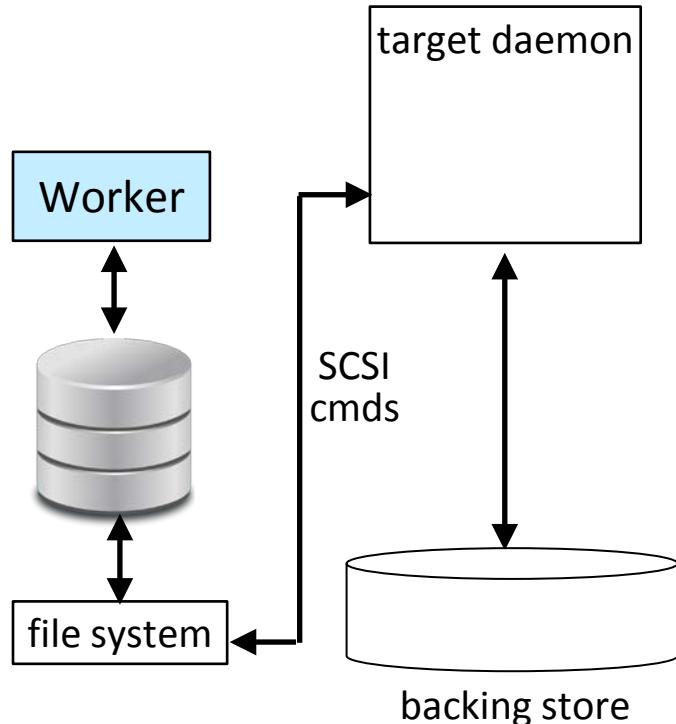
Framework Overview



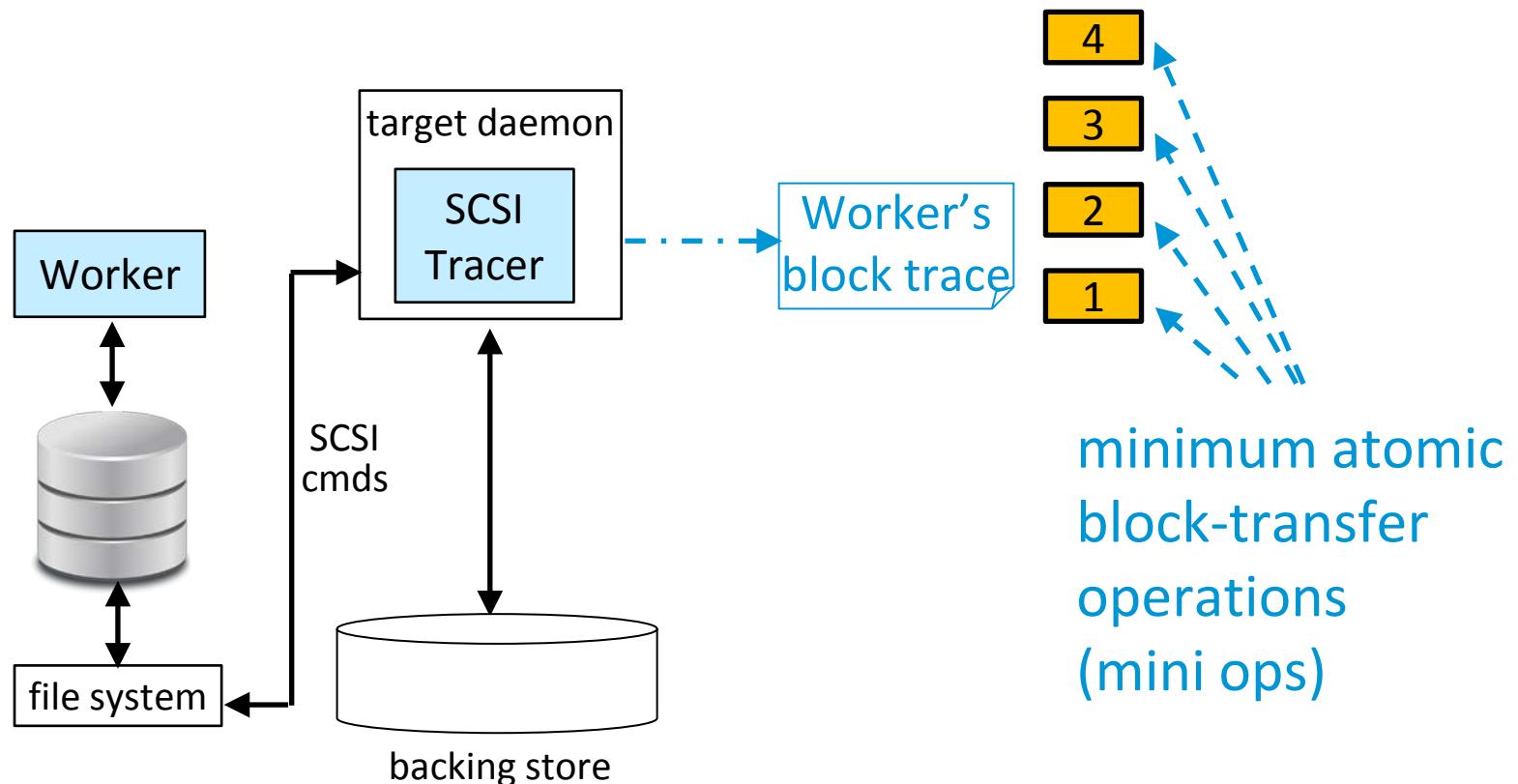
Framework Overview



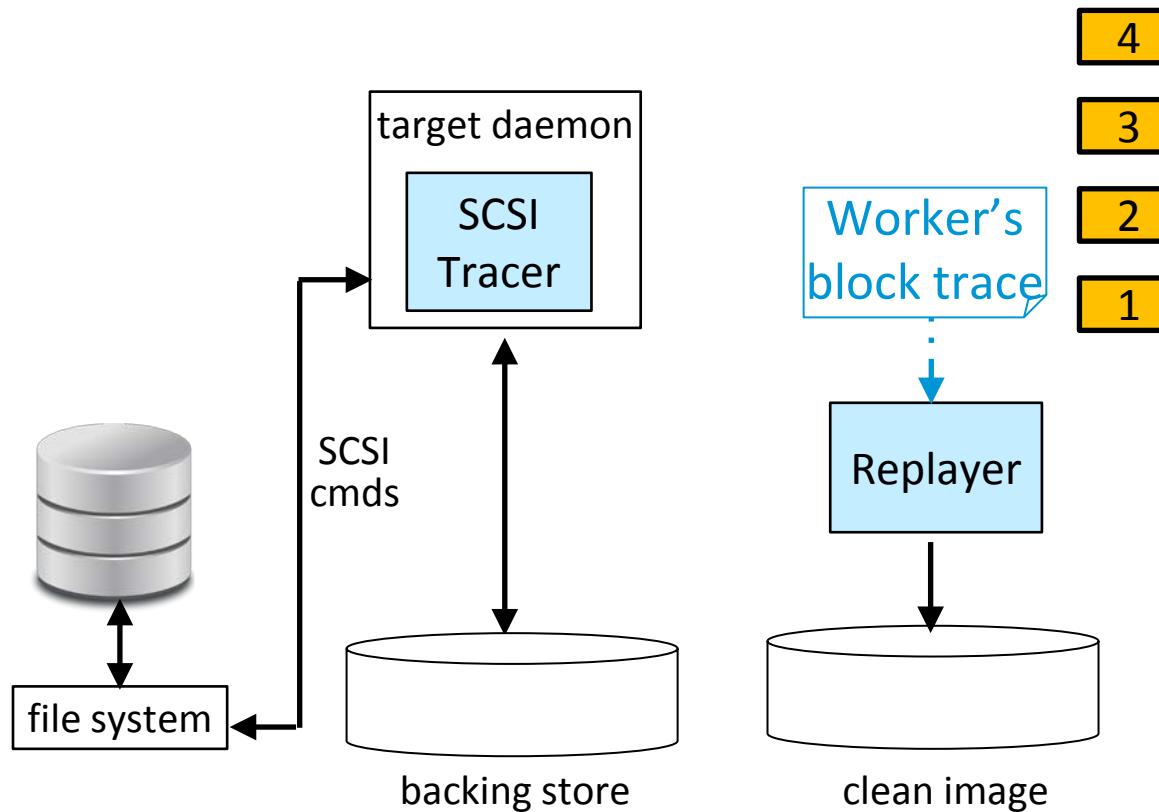
Capturing I/O trace without kernel modification



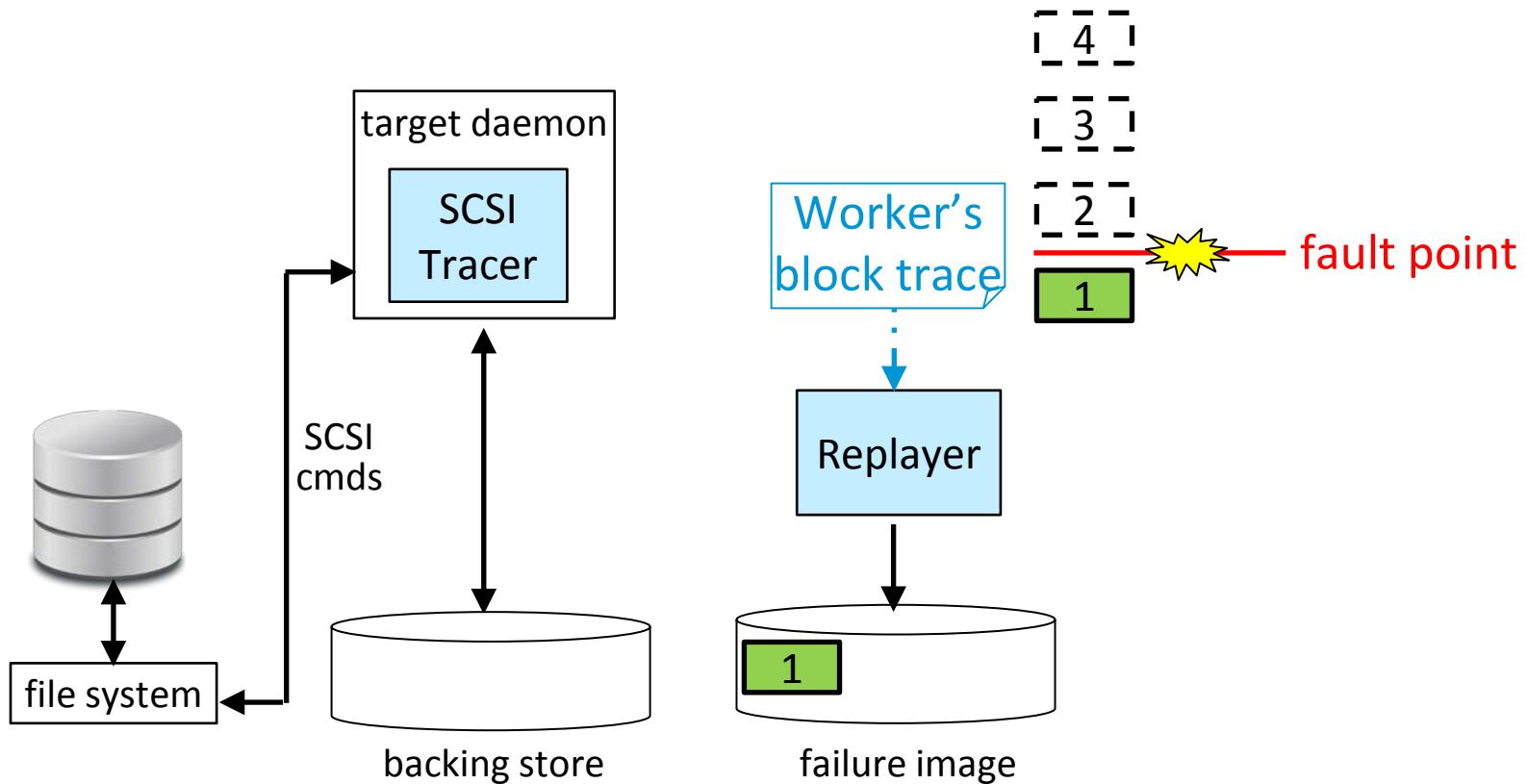
Capturing I/O trace without kernel modification



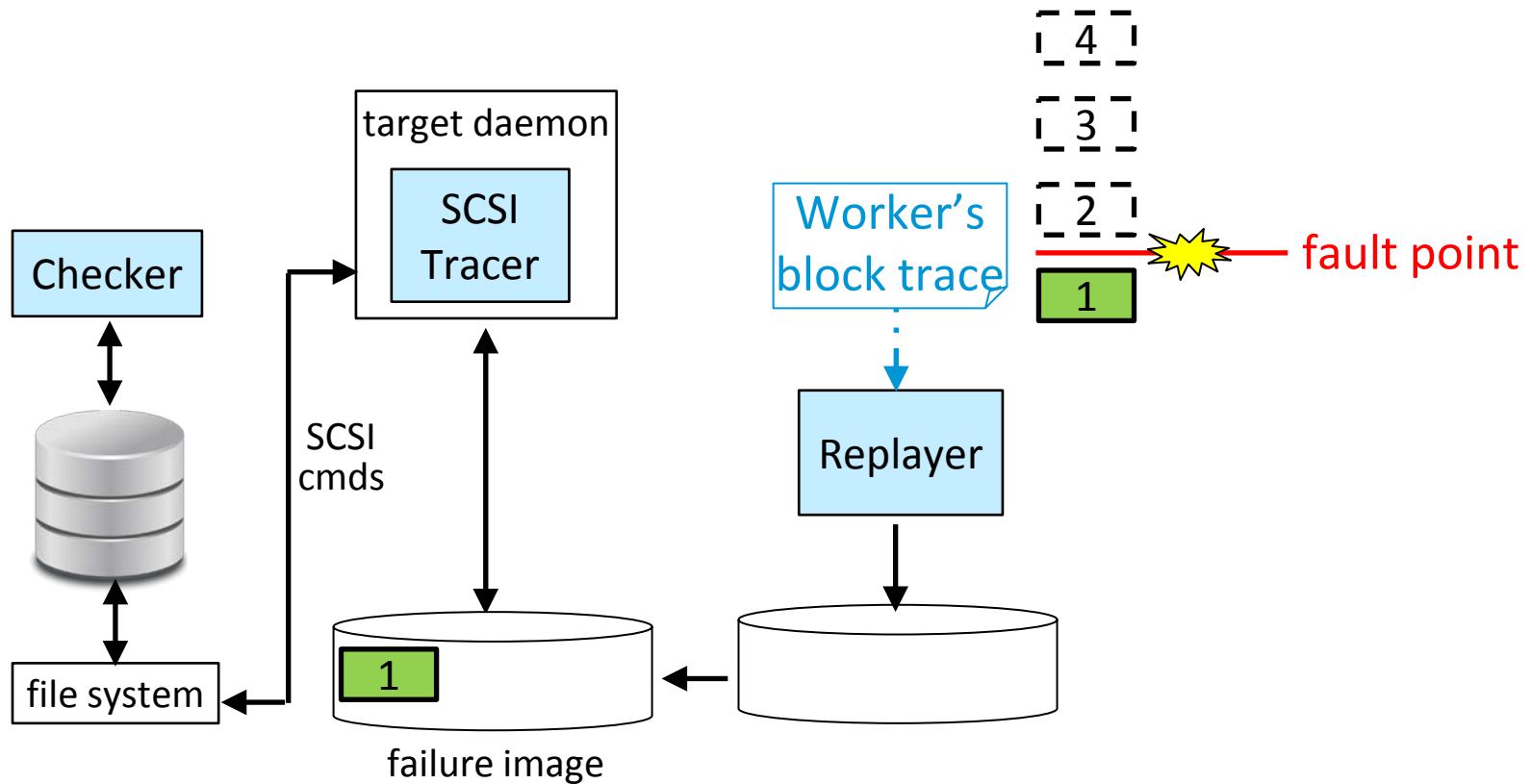
Constructing a post-fault disk image



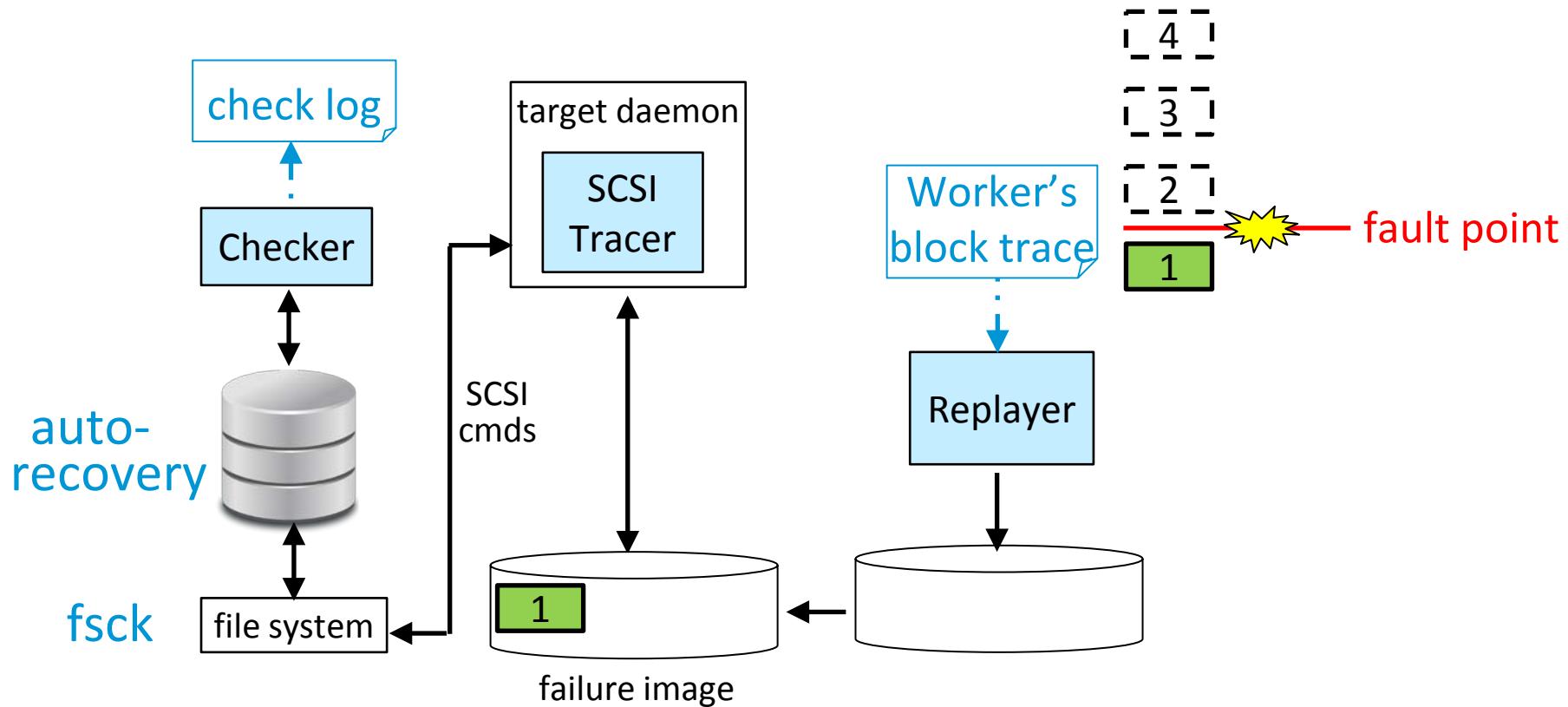
Constructing a post-fault disk image



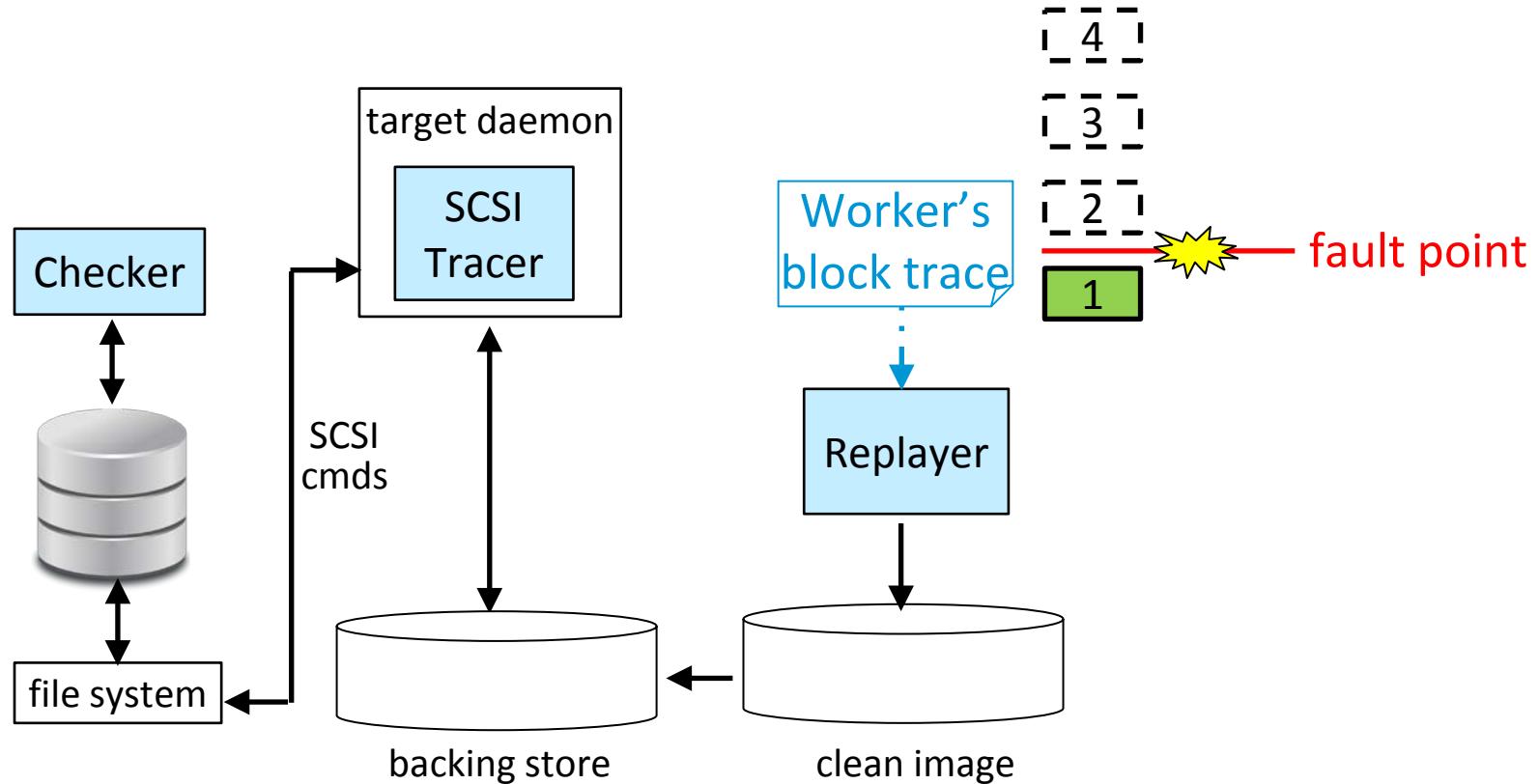
Checking the post-fault DB



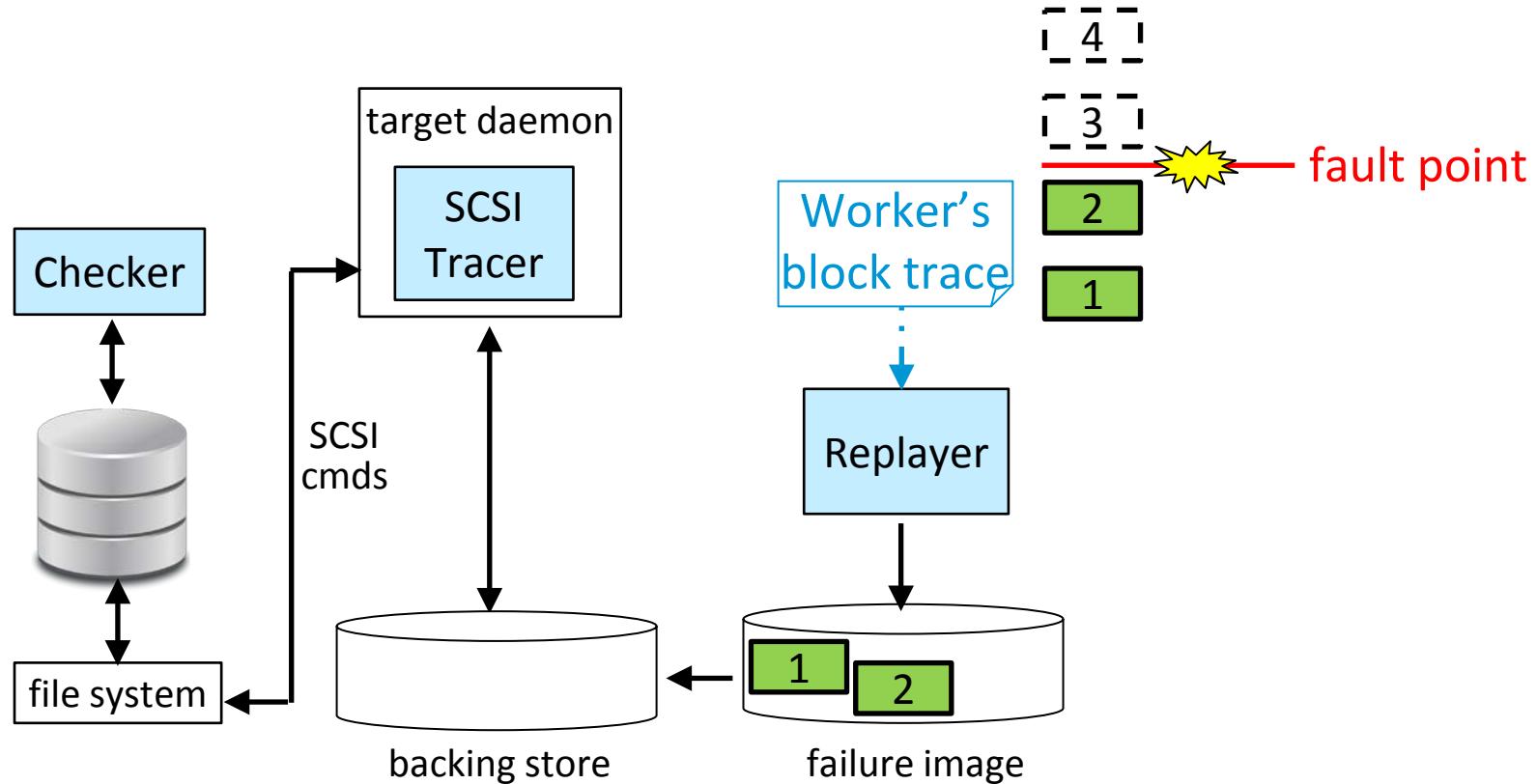
Checking the post-fault DB



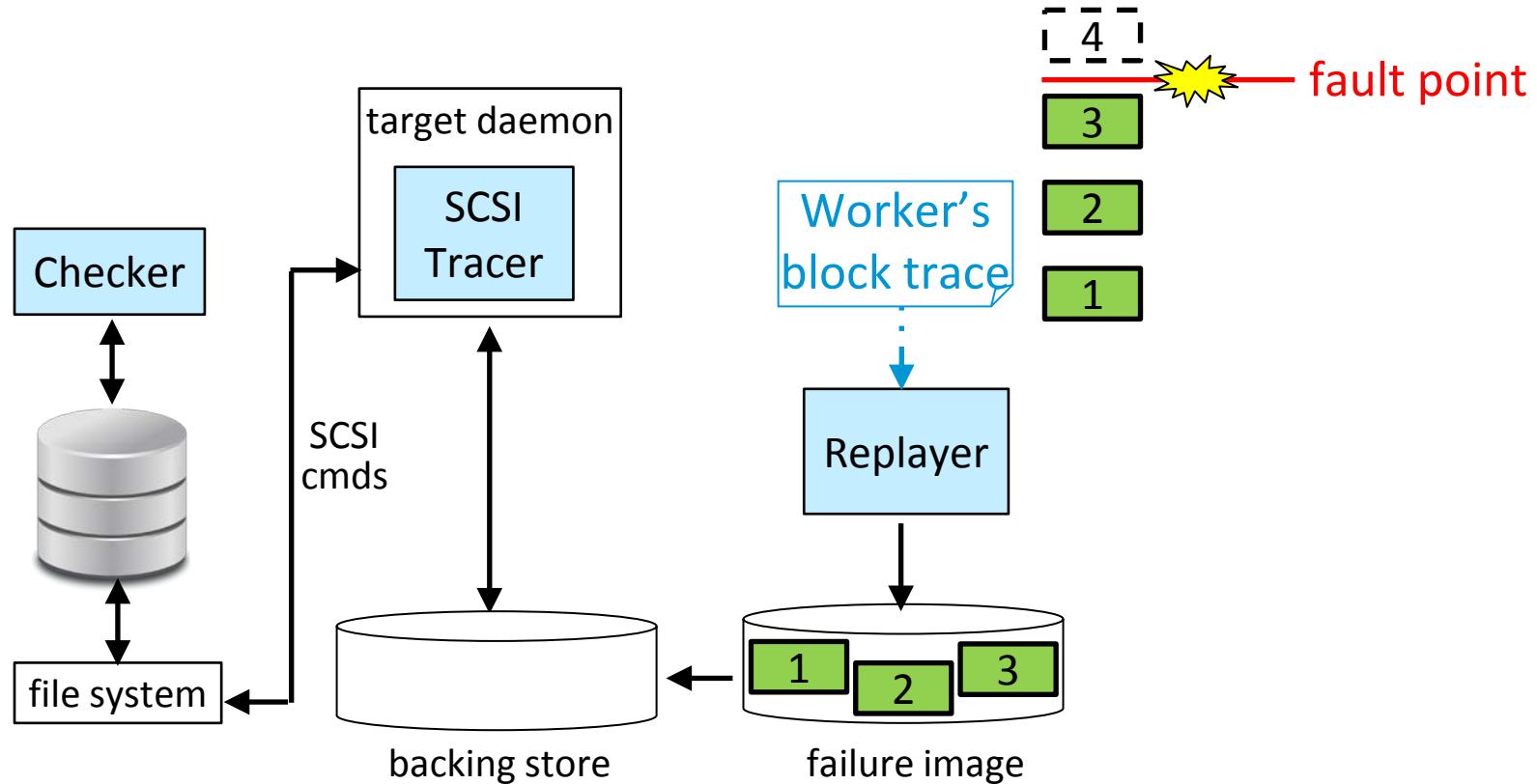
Testing different fault points easily



Testing different fault points easily

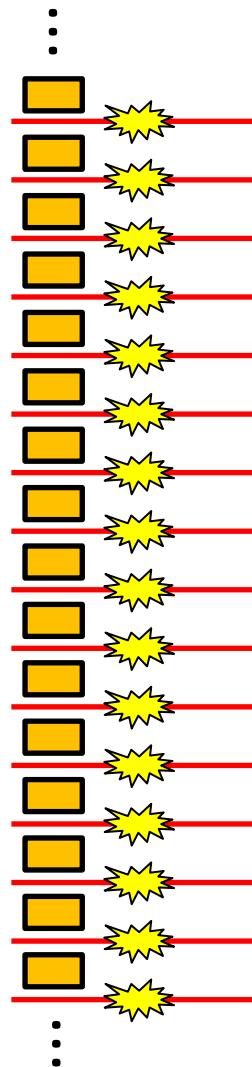


Testing different fault points easily



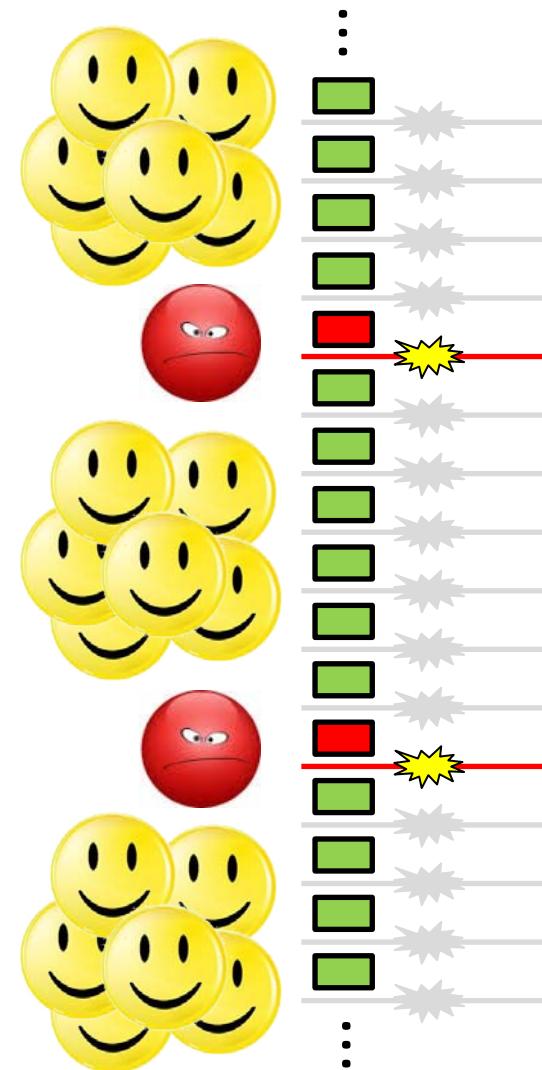
The framework is not good enough

- Sometimes need several days
 - too many mini operations, too many potential fault points



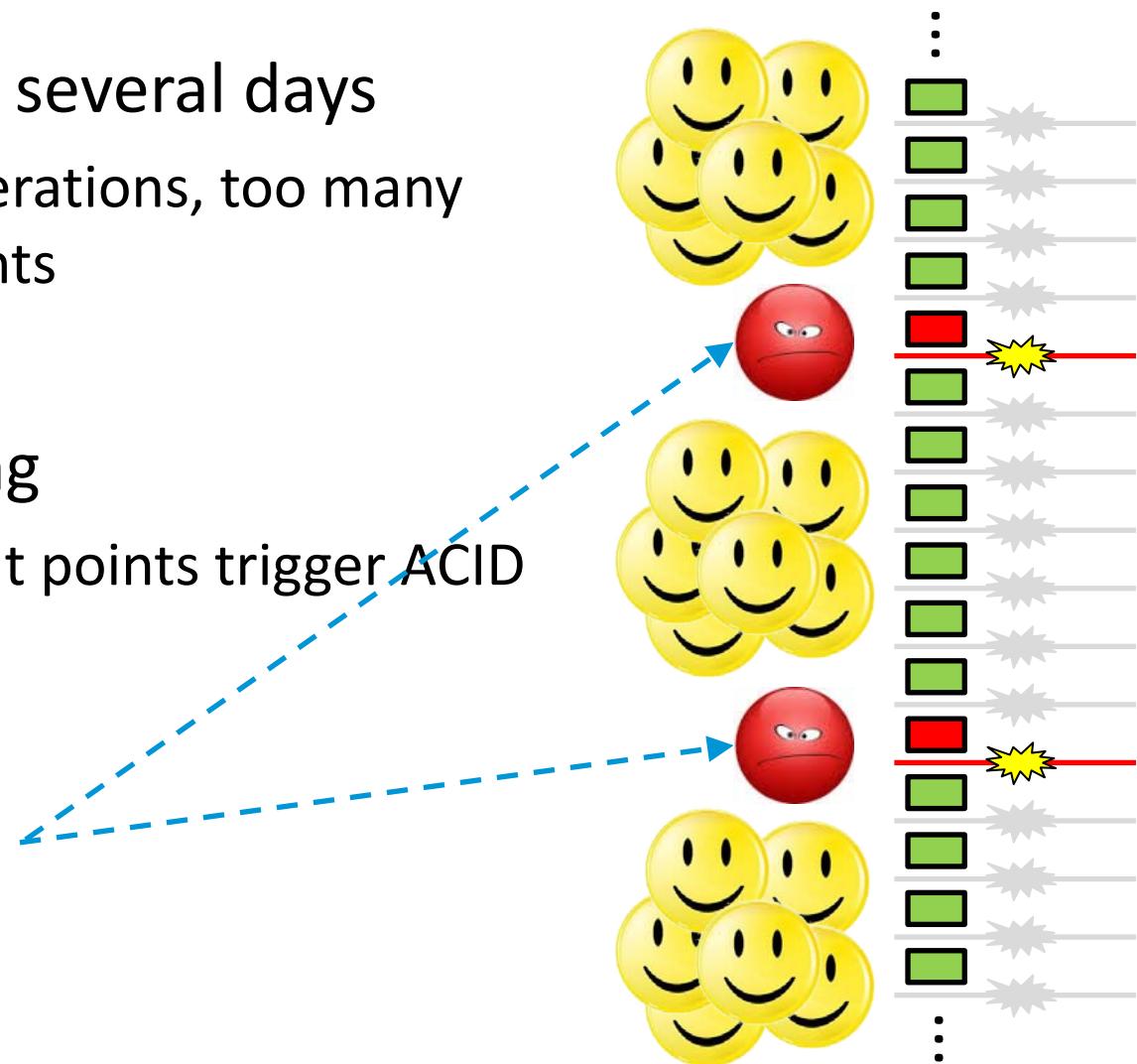
The framework is not good enough

- Sometimes need several days
 - too many mini operations, too many potential fault points
- We tried sampling
 - but only a few fault points trigger ACID violations



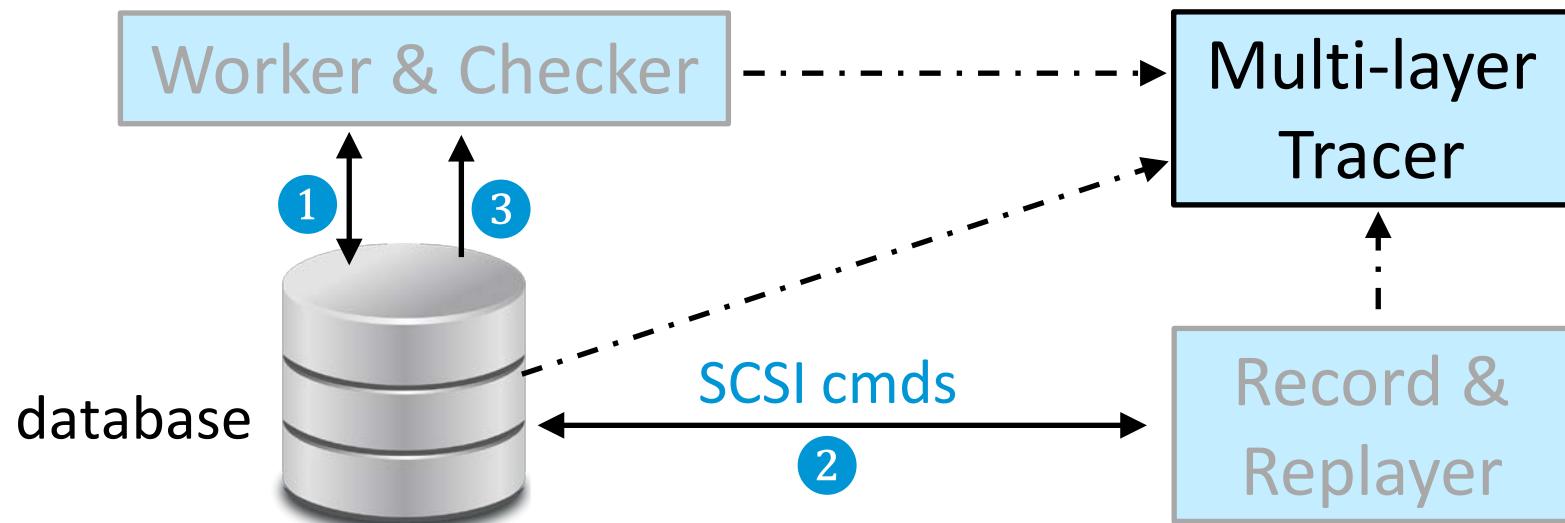
The framework is not good enough

- Sometimes need several days
 - too many mini operations, too many potential fault points
- We tried sampling
 - but only a few fault points trigger ACID violations
- Don't know why

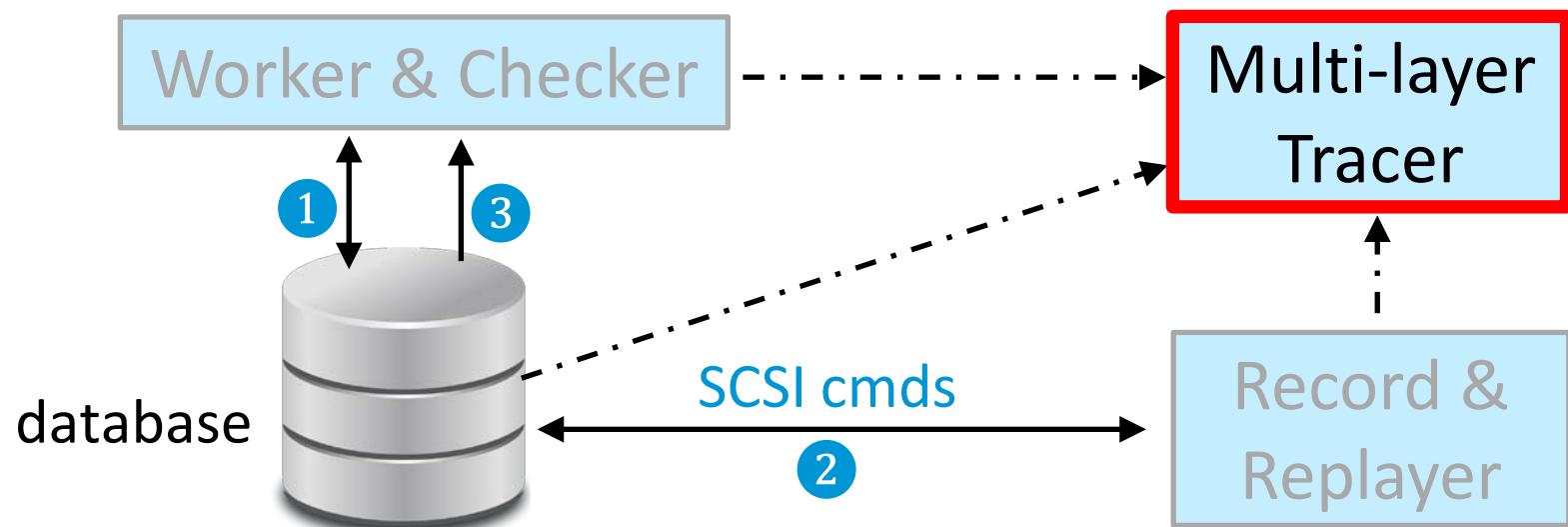


Enhanced Design

Framework Overview

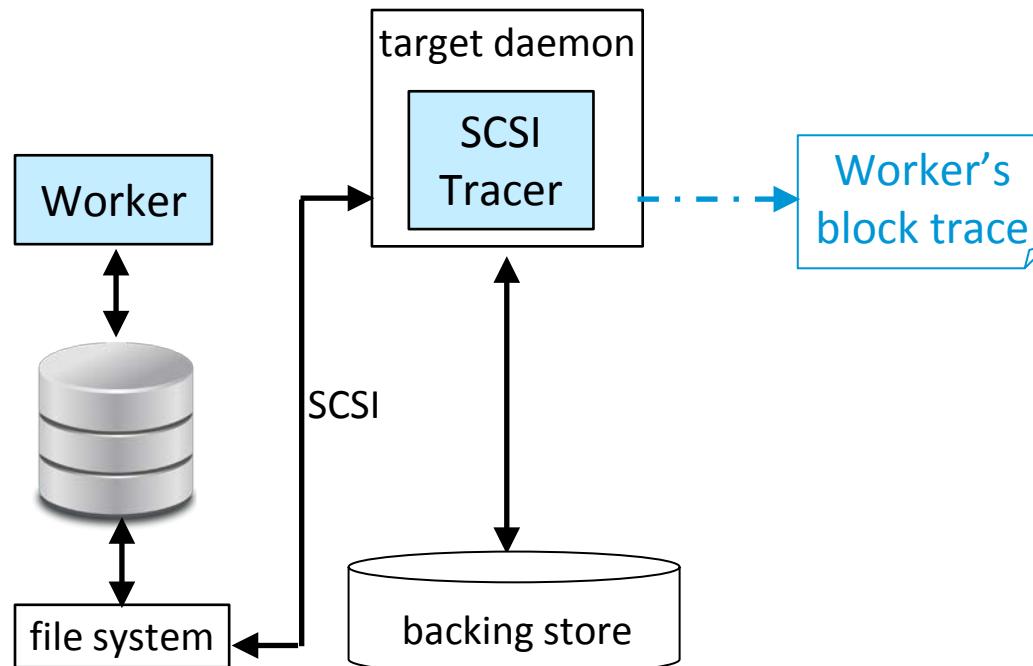


Framework Overview



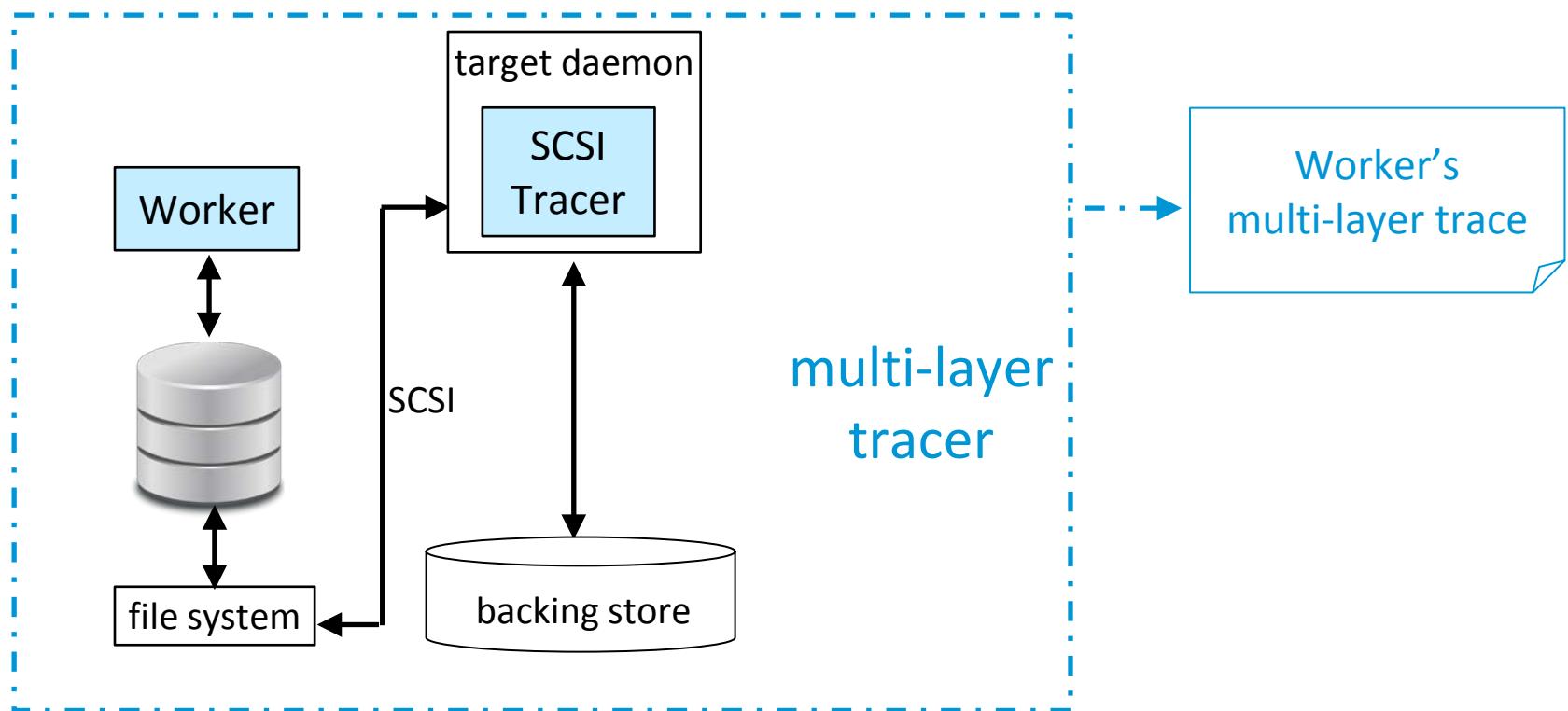
Original trace provides little semantics

op#	content	LBA
1	0a080101 ...	1012
2	0a080001 ...	6541
3	98393bc0 ...	9598
4	00000100 ...	9602



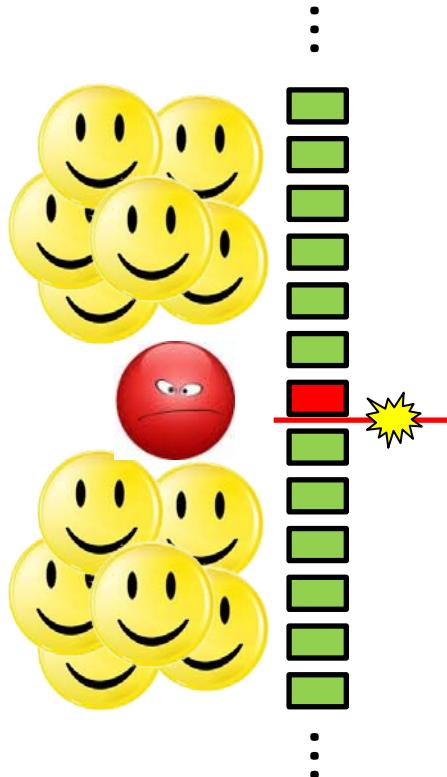
Enhancing w/ more context

op#	content	LBA	timestamp	SCSI cmd#	file	syscall
1	0a080101 ...	1012	139...013065	1	x.db	msync(x.db)
2	0a080001 ...	6541	139...210438	2	x.log	fsync(x.log)
3	98393bc0 ...	9598	139...355253	3	fs-j	fsync(x.log)
4	00000100 ...	9602	139...506097	3	fs-j	fsync(x.log)



What makes some fault points special?

Checking result



Worker's multi-layer trace

anything special?

5 patterns found from 2 databases

- MMAP_p : unintended update to mmap'ed blocks

		op#	LBA	file	syscall
	⋮
😊	█	463	1012	x.db	fsync(x.log)
😡	⋮
😡	⋮
😡	⋮
😊	█	564	1012	x.db	msync(x.db)
	⋮

5 patterns found from 2 databases

- MMAP_p : unintended update to mmap'ed blocks

	op#	LBA	file	syscall
⋮	…	…	…	…
😊	463	1012	x.db	fsync(x.log)
🔴	⋮	…	…	…
🔴	⋮	…	…	…
🔴	⋮	…	…	…
😊	564	1012	x.db	msync(x.db)
⋮	…	…	…	…

5 patterns found from 2 databases

- MMAP_p : unintended update to mmap'ed blocks

	op#	LBA	file	syscall	
⋮	…	…	…	…	
😊	463	1012	x.db	fsync(x.log)	
😡	…	…	…	…	
😡	…	…	…	…	
😡	…	…	…	…	
😊	564	1012	x.db	msync(x.db)	intended
⋮	…	…	…	…	

5 patterns found from 2 databases

- MMAP_p : unintended update to mmap'ed blocks

	op#	LBA	file	syscall	
⋮	⋮	⋮	⋮	⋮	⋮
😊	463	1012	x.db	fsync(x.log)	
😡	⋮	⋮	⋮	⋮	⋮
😡	⋮	⋮	⋮	⋮	⋮
😡	⋮	⋮	⋮	⋮	⋮
😊	564	1012	x.db	msync(x.db)	intended
⋮	⋮	⋮	⋮	⋮	⋮

5 patterns found from 2 databases

- MMAP_p : unintended update to mmap'ed blocks

	op#	LBA	file	syscall	
⋮	…	…	…	…	…
😊	463	1012	x.db	fsync x.log	unintended
😡	…	…	…	…	…
😡	…	…	…	…	…
😡	…	…	…	…	…
😊	564	1012	x.db	msync(x.db)	intended
⋮	…	…	…	…	…

implicit flush of dirty blocks by kernel or FS
under heavy transactions

5 patterns found from 2 databases

- MMAP_p : unintended update to mmap'ed blocks

	op#	LBA	file	syscall	
⋮	
😊	463	1012	x.db	fsync(x.log)	unintended
😡	special!
😡	
😡	
😊	564	1012	x.db	msync(x.db)	intended
⋮	

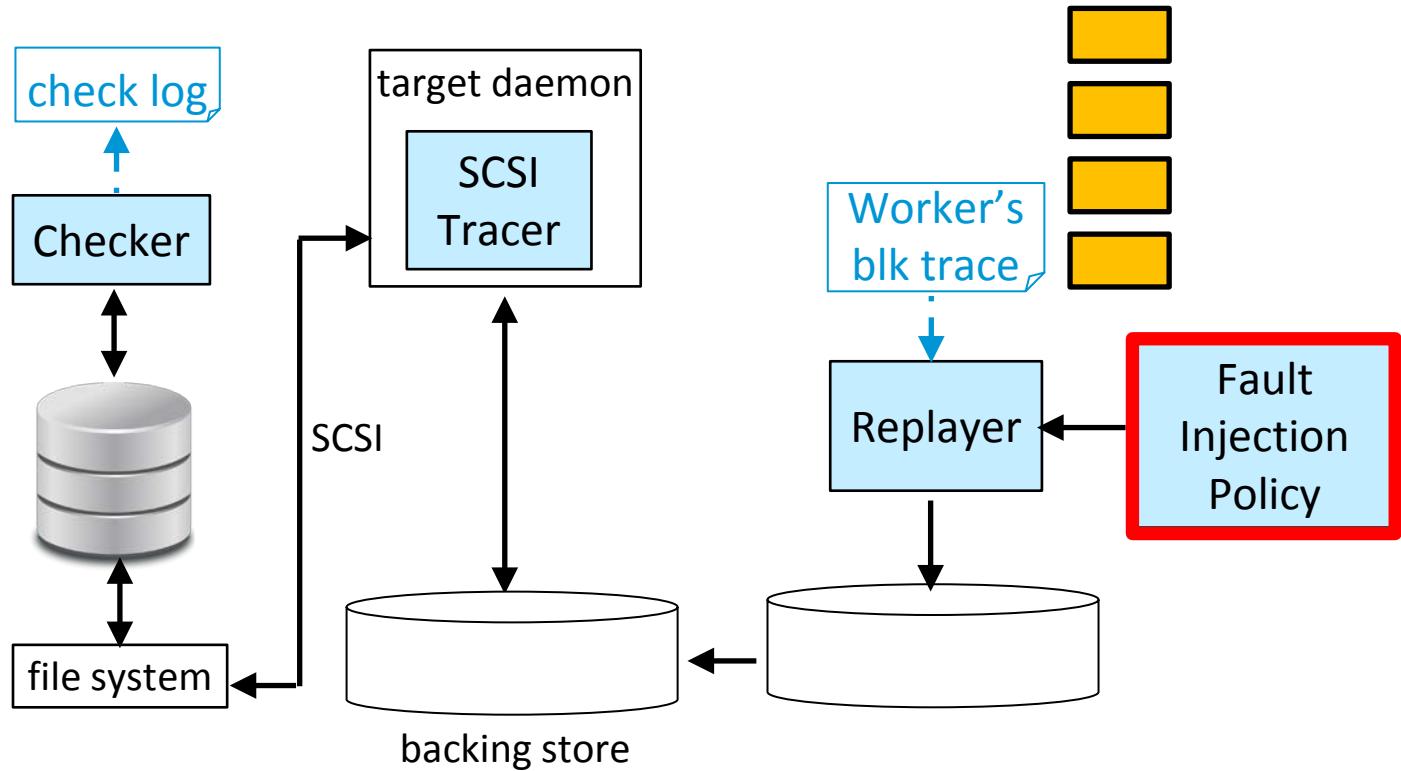
5 patterns found from 2 databases

- MMAP_p : unintended update to mmap'ed blocks

	op#	LBA	file	syscall	
⋮	
😊	463	1012	x.db	fsync(x.log)	unintended
😡	special!
😡	
😡	
😊	564	1012	x.db	msync(x.db)	intended
⋮	

- Four more patterns: REP_p , JUMP_p , HEAD_p , TRAN_p

Add fault injection policy to determine where to inject faults



Alternative to Exhaustive: Pattern-based Ranking

Worker's multi-layer trace					Scoreboard				total score	
op#	LBA	cmd#	file	syscall	MMAP _p	REP _p	JUMP _p	HEAD _p	TRAN _p	
1	348	1	x.db	msyc(x.db)	1					
2	352	2	x.log	fsync(x.log)	2					
3	356	2	x.log	fsync(x.log)	3					
4	360	2	x.log	fsync(x.log)	4					
5	364	2	x.log	fsync(x.log)	5					
6	370	3	x.log	fsync(x.log)	6					
7	348	4	x.db	fsync(x.log)	7					
8	906	5	fs-j	fsync(x.log)	8					

Alternative to Exhaustive: Pattern-based Ranking

Worker's multi-layer trace					Scoreboard				total score	
op#	LBA	cmd#	file	syscall	MMAP _p	REP _p	JUMP _p	HEAD _p	TRAN _p	
1	348	1	x.db	msyc(x.db)	1	0	0	0	0	0
2	352	2	x.log	fsync(x.log)	2	0	0	0	0	0
3	356	2	x.log	fsync(x.log)	3	0	0	0	0	0
4	360	2	x.log	fsync(x.log)	4	0	0	0	0	0
5	364	2	x.log	fsync(x.log)	5	0	0	0	0	0
6	370	3	x.log	fsync(x.log)	6	0	0	0	0	0
7	348	4	x.db	fsync(x.log)	7	1	0	0	0	1
8	906	5	fs-j	fsync(x.log)	8	1	0	0	0	1

Alternative to Exhaustive: Pattern-based Ranking

Worker's multi-layer trace					Scoreboard					total score	
op#	LBA	cmd#	file	syscall		MMAP _p	REP _p	JUMP _p	HEAD _p	TRAN _p	
1	348	1	x.db	msyc(x.db)	1	0	1				
2	352	2	x.log	fsync(x.log)	2	0	0				
3	356	2	x.log	fsync(x.log)	3	0	0				
4	360	2	x.log	fsync(x.log)	4	0	0				
5	364	2	x.log	fsync(x.log)	5	0	0				
6	370	3	x.log	fsync(x.log)	6	0	0				
7	348	4	x.db	fsync(x.log)	7	1	1				
8	906	5	fs-j	fsync(x.log)	8	1	0				

Alternative to Exhaustive: Pattern-based Ranking

Worker's multi-layer trace					Scoreboard					total score	
op#	LBA	cmd#	file	syscall		MMAP _p	REP _p	JUMP _p	HEAD _p	TRAN _p	
1	348	1	x.db	msyc(x.db)	1	0	1	0			
2	352	2	x.log	fsync(x.log)	2	0	0	0			
3	356	2	x.log	fsync(x.log)	3	0	0	0			
4	360	2	x.log	fsync(x.log)	4	0	0	0			
5	364	2	x.log	fsync(x.log)	5	0	0	0			
6	370	3	x.log	fsync(x.log)	6	0	0		1		
7	348	4	x.db	fsync(x.log)	7	1	1	1			
8	906	5	fs-j	fsync(x.log)	8	1	0		1		

Alternative to Exhaustive: Pattern-based Ranking

Worker's multi-layer trace					Scoreboard					total score
op#	LBA	cmd#	file	syscall		MMAP _p	REP _p	JUMP _p	HEAD _p	TRAN _p
1	348	1	x.db	msyc(x.db)	1	0	1	0	0	
2	352	2	x.log	fsync(x.log)	2	0	0	0	1	
3	356	2	x.log	fsync(x.log)	3	0	0	0	0	
4	360	2	x.log	fsync(x.log)	4	0	0	0	0	
5	364	2	x.log	fsync(x.log)	5	0	0	0	0	
6	370	3	x.log	fsync(x.log)	6	0	0	1	0	
7	348	4	x.db	fsync(x.log)	7	1	1	1	0	
8	906	5	fs-j	fsync(x.log)	8	1	0	1	0	

Alternative to Exhaustive: Pattern-based Ranking

Worker's multi-layer trace					Scoreboard					total score
op#	LBA	cmd#	file	syscall	MMAP _p	REP _p	JUMP _p	HEAD _p	TRAN _p	
1	348	1	x.db	msyc(x.db)	1	0	1	0	0	1
2	352	2	x.log	fsync(x.log)	2	0	0	0	1	1
3	356	2	x.log	fsync(x.log)	3	0	0	0	0	0
4	360	2	x.log	fsync(x.log)	4	0	0	0	0	0
5	364	2	x.log	fsync(x.log)	5	0	0	0	0	0
6	370	3	x.log	fsync(x.log)	6	0	0	1	0	1
7	348	4	x.db	fsync(x.log)	7	1	1	1	0	1
8	906	5	fs-j	fSync(x.log)	8	1	0	1	0	1

Alternative to Exhaustive: Pattern-based Ranking

Worker's multi-layer trace					Scoreboard					total score	
op#	LBA	cmd#	file	syscall	MMAP _p	REP _p	JUMP _p	HEAD _p	TRAN _p		
1	348	1	x.db	msync(x.db)	1	0	1	0	0	1	2
2	352	2	x.log	fsync(x.log)	2	0	0	0	1	1	2
3	356	2	x.log	fsync(x.log)	3	0	0	0	0	0	0
4	360	2	x.log	fsync(x.log)	4	0	0	0	0	0	0
5	364	2	x.log	fsync(x.log)	5	0	0	0	0	0	0
6	370	3	x.log	fsync(x.log)	6	0	0	1	0	1	2
7	348	4	x.db	fsync(x.log)	7	1	1	1	0	1	4
8	906	5	fs-j	fsync(x.log)	8	1	0	1	0	1	3

Alternative to Exhaustive: Pattern-based Ranking

Worker's multi-layer trace					Scoreboard					total score
op#	LBA	cmd#	file	syscall	MMAP _p	REP _p	JUMP _p	HEAD _p	TRAN _p	
1	348	1	x.db	msyc(x.db)	1	0	1	0	1	2
2	352	2	x.log	fsync(x.log)	2	0	0	0	1	1
3	356	2	x.log	fsync(x.log)	3	0	0	0	0	0
4	360	2	x.log	fsync(x.log)	4	0	0	0	0	0
5	364	2	x.log	fsync(x.log)	5	0	0	0	0	0
6	370	3	x.log	fsync(x.log)	6	0	0	1	0	1
7	348	4	x.db	fsync(x.log)	7	1	1	1	0	1
8	906	5	fs-j	fsync(x.log)	8	1	0	1	0	1

1st-rank:  ←----- predicted most error-prone
 2nd-rank: 
 3rd-rank:   
 4th -rank:   

Alternative to Exhaustive: Pattern-based Ranking



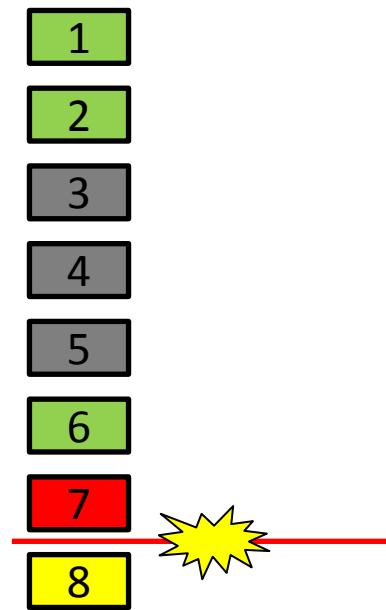
1st-rank: 7 ←----- predicted most error-prone

2nd-rank: 8

3rd-rank: 1 2 6

4th -rank: 3 4 5

Alternative to Exhaustive: Pattern-based Ranking



1st-rank: 7 ←----- predicted most error-prone

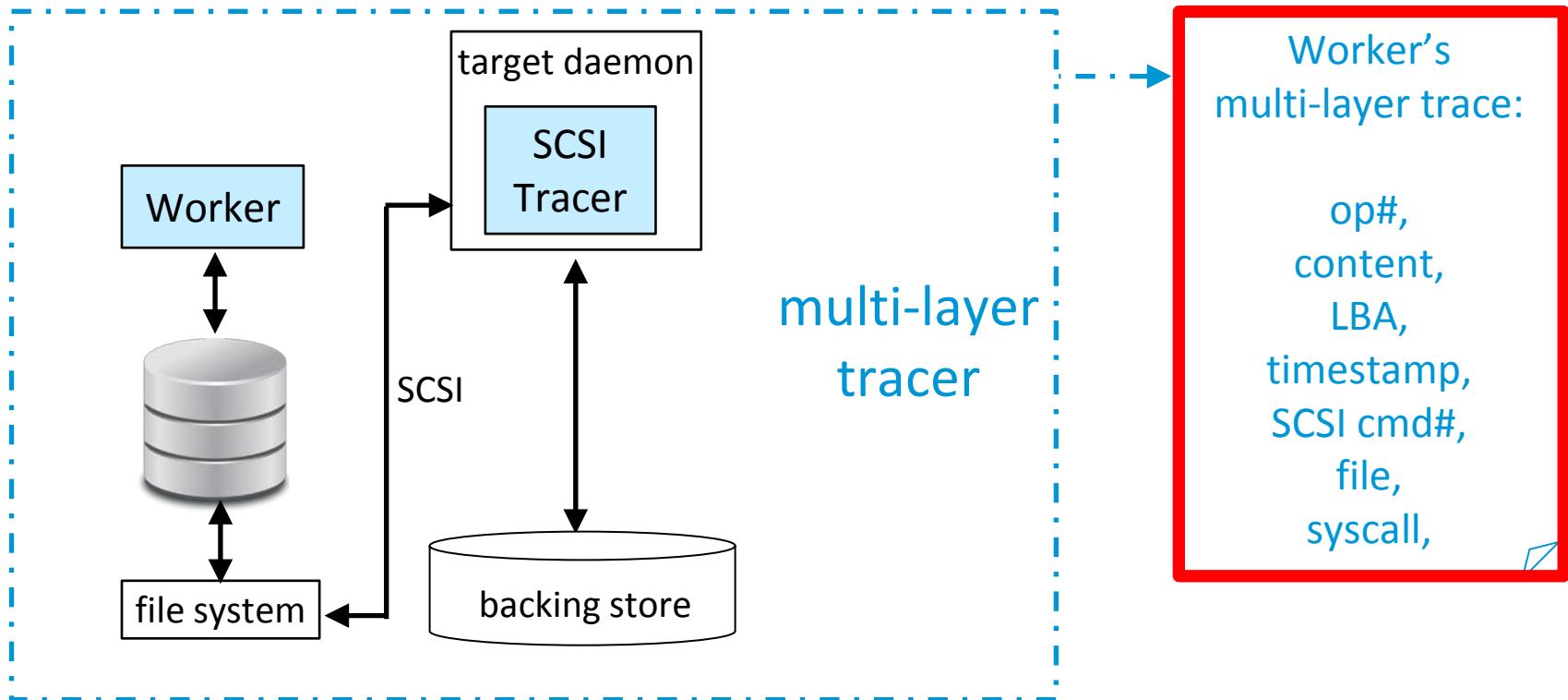
2nd-rank: 8

3rd-rank: 1 2 6

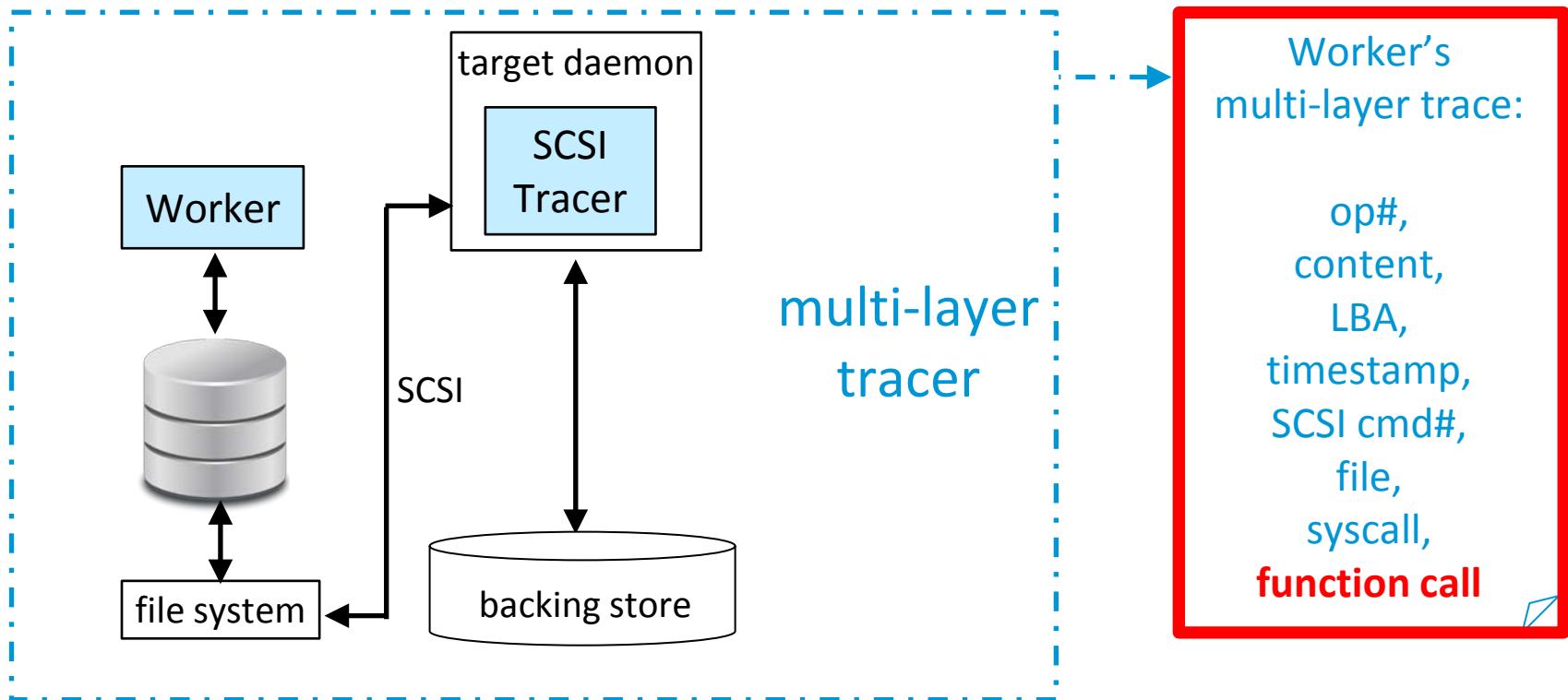
4th -rank: 3 4 5

Diagnosis Support

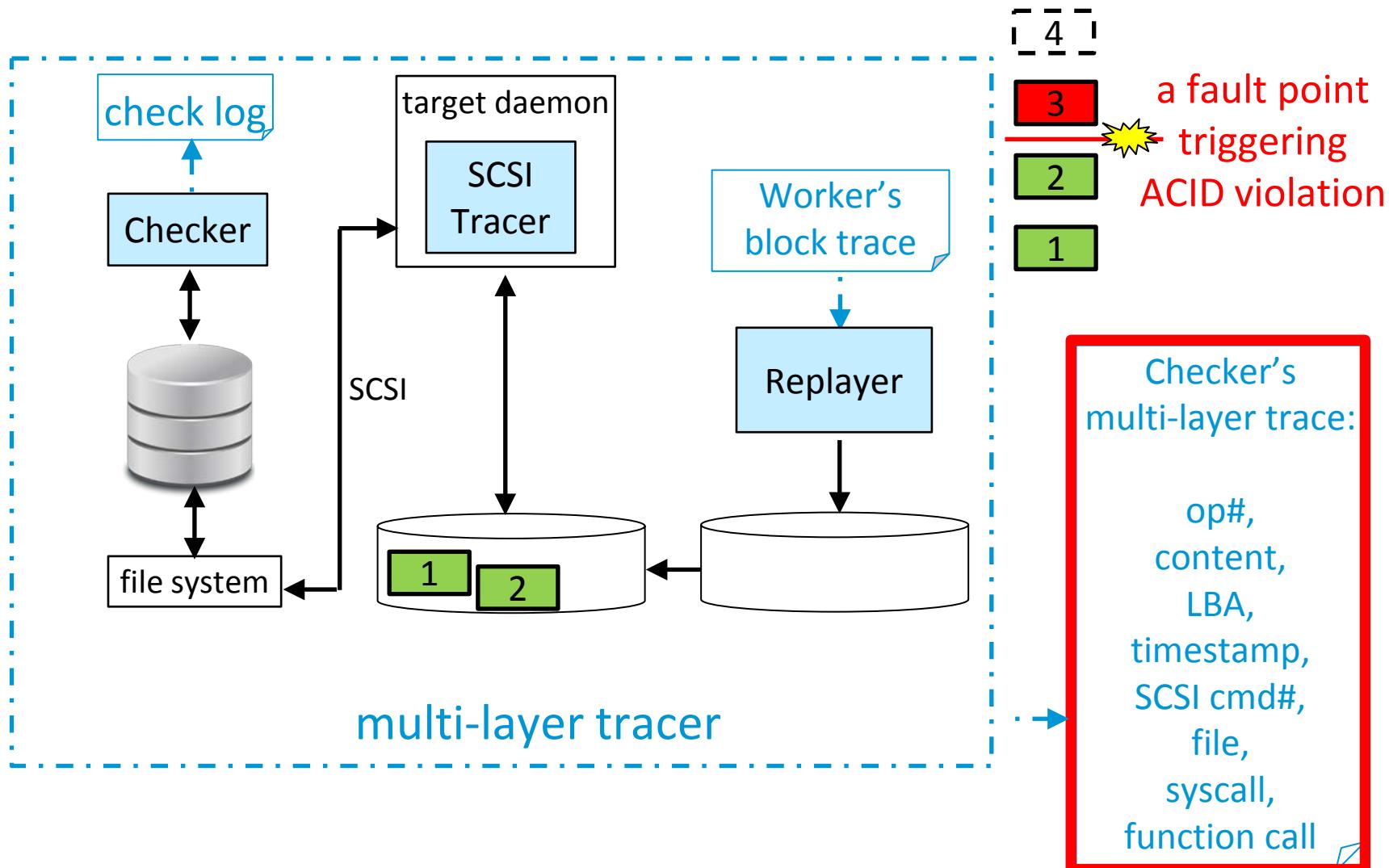
Worker's multi-layer trace helps understand what happened at fault time



Add function-call tracing to disclose more semantics for diagnosis



Enable same tracing during checking to see why recovery didn't work



Results

Experimental Environment

- 8 databases
 - Open-source: TokyoCabinet, MariaDB, LightningDB, SQLite
 - Commercial: KVS-A, SQL-A, SQL-B, SQL-C
- 4 workloads
- 3 file systems
 - ext3, XFS, NTFS
- Several operating systems
 - Linux: RHEL 6, Debian6, Ubuntu 12 LTS
 - Windows 7 Enterprise

Not a single DB can survive all tests

DB	FS	W-1	W-2	W-3	W-4.1	W-4.2	W-4.3	A	C	I	D
TokyoCabinet	ext3	D	D	D	ACD	ACD	ACD	0.15%	0.14%	0	16.05%
	XFS	--	D	D	ACD	D	ACD	<0.01%	0.01%	0	4.38%
MariaDB	ext3	D	D	D	D	D	D	0	0	0	1.36%
	XFS	D	D	D	D	D	D	0	0	0	0.49%
LightningDB	ext3	--	--	--	--	--	D	0	0	0	0.05%
	XFS	--	--	--	--	--	--	0	0	0	0
SQLite	ext3	D	D	--	D	D	D	0	0	0	19.15%
	XFS	--	--	D	D	D	D	0	0	0	10.60%
KVS-A	ext3	--	--	Hang	--	--	--	0	0	0	0
	XFS	--	--	--	--	--	--	0	0	0	0
SQL-A	ext3	D	D	D	D	D	D	0	0	0	3.31%
	XFS	D	D	D	D	D	D	0	0	0	0.92%
SQL-B	ext3	D	D	CD	CD	CD	CD	0	8.96%	0	3.24%
	XFS	CD	D	CD	CD	CD	CD	0	7.77%	0	3.90%
SQL-C	NTFS	D	D	D	D	D	D	0	0	0	8.08%

Durability violation is most common

DB	FS	W-1	W-2	W-3	W-4.1	W-4.2	W-4.3	A	C	I	D
TokyoCabinet	ext3	D	D	D	ACD	ACD	ACD	0.15%	0.14%	0	16.05%
	XFS	--	D	D	ACD	D	ACD	<0.01%	0.01%	0	4.38%
MariaDB	ext3	D	D	D	D	D	D	0	0	0	1.36%
	XFS	D	D	D	D	D	D	0	0	0	0.49%
LightningDB	ext3	--	--	--	--	--	D	0	0	0	0.05%
	XFS	--	--	--	--	--	--	0	0	0	0
SQLite	ext3	D	D	--	D	D	D	0	0	0	19.15%
	XFS	--	--	D	D	D	D	0	0	0	10.60%
KVS-A	ext3	--	--	Hang	--	--	--	0	0	0	0
	XFS	--	--	--	--	--	--	0	0	0	0
SQL-A	ext3	D	D	D	D	D	D	0	0	0	3.31%
	XFS	D	D	D	D	D	D	0	0	0	0.92%
SQL-B	ext3	D	D	CD	CD	CD	CD	0	8.96%	0	3.24%
	XFS	CD	D	CD	CD	CD	CD	0	7.77%	0	3.90%
SQL-C	NTFS	D	D	D	D	D	D	0	0	0	8.08%

Some violations are difficult to trigger

DB	FS	W-1	W-2	W-3	W-4.1	W-4.2	W-4.3	A	C	I	D
TokyoCabinet	ext3	D	D	D	ACD	ACD	ACD	0.15%	0.14%	0	16.05%
	XFS	--	D	D	ACD	D	ACD	<0.01%	0.01%	0	4.38%
MariaDB	ext3	D	D	D	D	D	D	0	0	0	1.36%
	XFS	D	D	D	D	D	D	0	0	0	0.49%
LightningDB	ext3	--	--	--	--	--	D	0	0	0	0.05%
	XFS	--	--	--	--	--	--	0	0	0	0
SQLite	ext3	D	D	--	D	D	D	0	0	0	19.15%
	XFS	--	--	D	D	D	D	0	0	0	10.60%
KVS-A	ext3	--	--	Hang	--	--	--	0	0	0	0
	XFS	--	--	--	--	--	--	0	0	0	0
SQL-A	ext3	D	D	D	D	D	D	0	0	0	3.31%
	XFS	D	D	D	D	D	D	0	0	0	0.92%
SQL-B	ext3	D	D	CD	CD	CD	CD	0	8.96%	0	3.24%
	XFS	CD	D	CD	CD	CD	CD	0	7.77%	0	3.90%
SQL-C	NTFS	D	D	D	D	D	D	0	0	0	8.08%

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	XFS	--	D	D	ACD	D	ACD	<0.01%	0.01%	0	4.38%
MariaDB	ext3	D	D	D	D	D	D	0	0	0	1.36%
	XFS	D	D	D	D	D	D	0	0	0	0.49%
LightningDB	ext3	--	--	--	--	--	D	0	0	0	0.05%
	XFS	--	--	--	--	--	--	0	0	0	0
SQLite	ext3	D	D	--	D	D	D	0	0	0	19.15%
	XFS	--	--	D	D	D	D	0	0	0	10.60%
KVS-A	ext3	--	--	Hang	--	--	--	0	0	0	0
	XFS	--	--	--	--	--	--	0	0	0	0
SQL-A	ext3	D	D	D	D	D	D	0	0	0	3.31%
	XFS	D	D	D	D	D	D	0	0	0	0.92%
SQL-B	ext3	D	D	CD	CD	CD	CD	0	8.96%	0	3.24%
	XFS	CD	D	CD	CD	CD	CD	0	7.77%	0	3.90%
SQL-C	NTFS	D	D	D	D	D	D	0	0	0	8.08%

Case Study: A TokyoCabinet Bug

- Failure symptoms
 - faults injected in a region of operations cause:
 - A violation: a transaction is partially committed
 - D violation: some rows are irretrievable
 - C violation: retrievable rows by range query and point queries are different

Case Study: A TokyoCabinet Bug

Why recovery didn't work?

```
...
tchdbopenimpl(x.tcb)
...
open(x.tcb) = 3
read(x.tcb) = 256
//op#, LBA, content, file
op#1, 480, ...100..., x.tcb
tcbdbget()
...
```

Checker's trace
when ACID violations
were found

```
...
tchdbopenimpl(x.tcb)
...
open(x.tcb) = 3
read(x.tcb) = 256
//op#, LBA, content, file
op#1, 480, ...101..., x.tcb
tchdbwalrestore()
tcbdbget()
...
```

Checker's trace
when no violation
was found

Delta Debugging
[Zeller, SIGSOFT'02/FSE-10]

Case Study: A TokyoCabinet Bug

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//op#, LBA, content, file
op#1, 480, ...100..., x.tcb
//no tchdbwalrestore()
tcldbget()
...
```

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//op#, LBA, content, file
op#1, 480, ...101..., x.tcb
tchdbwalrestore()
tcldbget()
...
```

Checker's trace
when no violation
was found

Delta Debugging
[Zeller, SIGSOFT'02/FSE-10]

Case Study: A TokyoCabinet Bug

What happened at fault time?

Worker's trace
around the bug-triggering
fault points **op#30–#90**

```
...
mmap(8192, x.tcb, 0)
...
fsync(x.tcb.wal)
//op#, LBA, content, file,      syscall
op#26, 630, ......., x.tcb.wal, fsync(x.tcb.wal)
op#27, 960, ......., fs-j ,    fsync(x.tcb.wal)
op#28, 964, ......., fs-j ,    fsync(x.tcb.wal)
op#29, 480, ...100..., x.tcb , fsync(x.tcb.wal)
...
msync(x.tcb)
//op#, LBA, content, file, syscall
op#91, 480, ...101..., x.tcb, msync(x.tcb)
...
```

Why recovery didn't work?

```
...
tchdbopenimpl(x.tcb)
...
open(x.tcb) = 3
read(x.tcb) = 256
//op#, LBA, content, file
op#1, 480, ...100..., x.tcb
//no tchdbwalrestore()
tcbdbget()
...
...
tchdbopenimpl(x.tcb)
...
open(x.tcb) = 3
read(x.tcb) = 256
//op#, LBA, content, file
op#1, 480, ...101..., x.tcb
tchdbwalrestore()
tcbdbget()
...
```

Case Study: A TokyoCabinet Bug

What happened at fault time?

Worker's trace
around the bug-triggering
fault points **op#30–#90**

```
...
mmap(8192, x.tcb, 0)
...
fsync(x.tcb.wal)
//op#, LBA, content, file,      syscall
op#26, 630, ......., x.tcb.wal, fsync(x.tcb.wal)
op#27, 960, ......., fs-j ,    fsync(x.tcb.wal)
op#28, 964, ......., fs-j ,    fsync(x.tcb.wal)
op#29, 480, ...100..., x.tcb , fsync(x.tcb.wal)
...
msync(x.tcb)
//op#, LBA, content, file, syscall
op#91, 480, ...101..., x.tcb, msync(x.tcb)
...

```

Why recovery didn't work?

```
...
tchdbopenimpl(x.tcb)
...
open(x.tcb) = 3
read(x.tcb) = 256
//op#, LBA, content, file
op#1, 480, ...100..., x.tcb
//no tchdbwalrestore()
tcbdbget()
...

```

```
...
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op#1, 480, ...101..., x.tcb
tchdbwalrestore()
tcbdbget()
...

```

Case Study: A TokyoCabinet Bug

What happened at fault time?

Worker's trace
around the bug-triggering
fault points **op#30–#90**

```
...
mmap(8192, x.tcb, 0)
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fsync(x.tcb.wal)
//op#, LBA, content, file,      syscall
op#26, 630, ......., x.tcb.wal, fsync(x.tcb.wal)
op#27, 960, ......., fs-j ,    fsync(x.tcb.wal)
op#28, 964, ......., fs-j ,    fsync(x.tcb.wal)
op#29, 480, ...100..., x.tcb , fsync(x.tcb.wal)
...
msync(x.tcb)
//op#, LBA, content, file, syscall
op#91, 480, ...101..., x.tcb, msync(x.tcb)
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Why recovery didn't work?

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← - - Intended update

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← -- Unintended update!

← -- Intended update

Case Study: A TokyoCabinet Bug

What happened at fault time?

Worker's trace
around the bug-triggering
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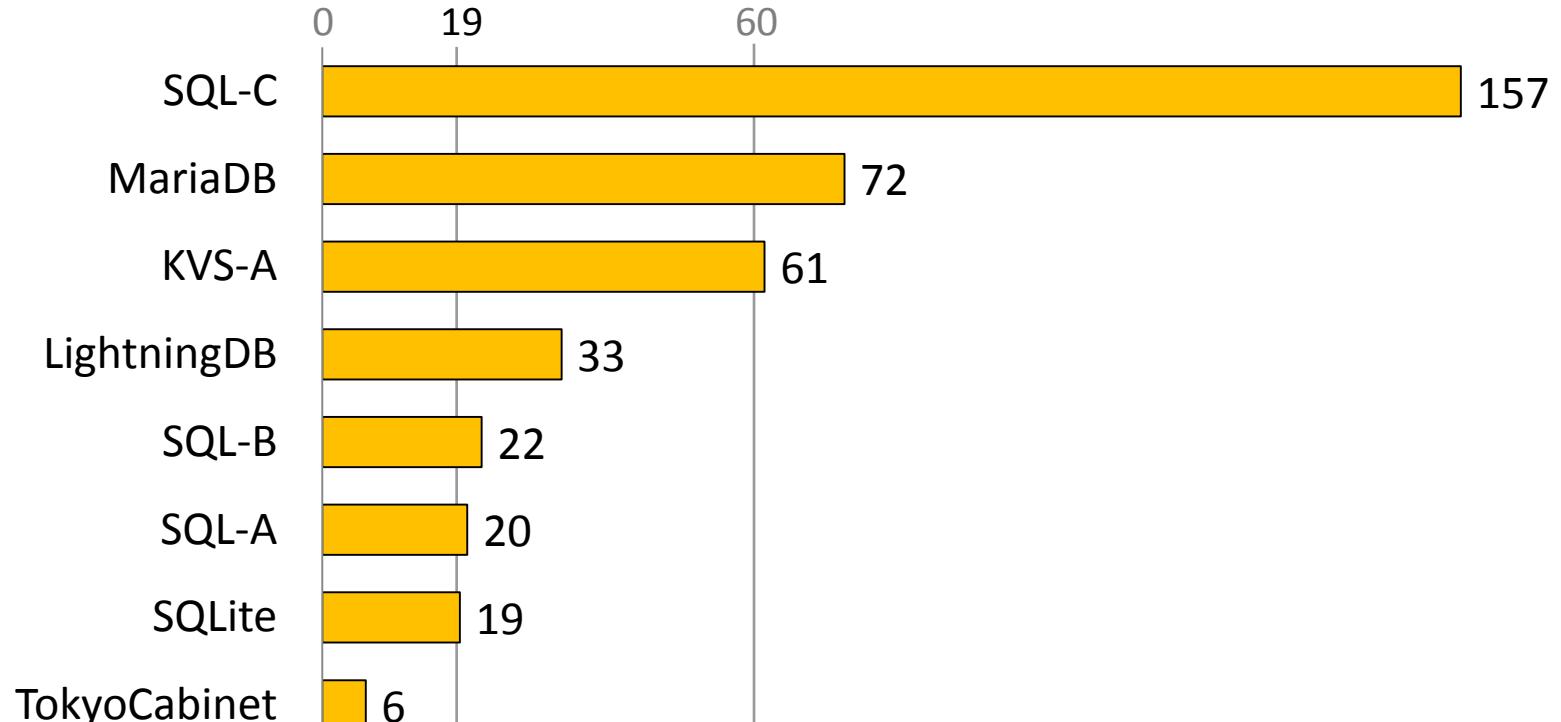
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One solution:
Failure-atomic msync()
[Park et.al., EuroSys'13]

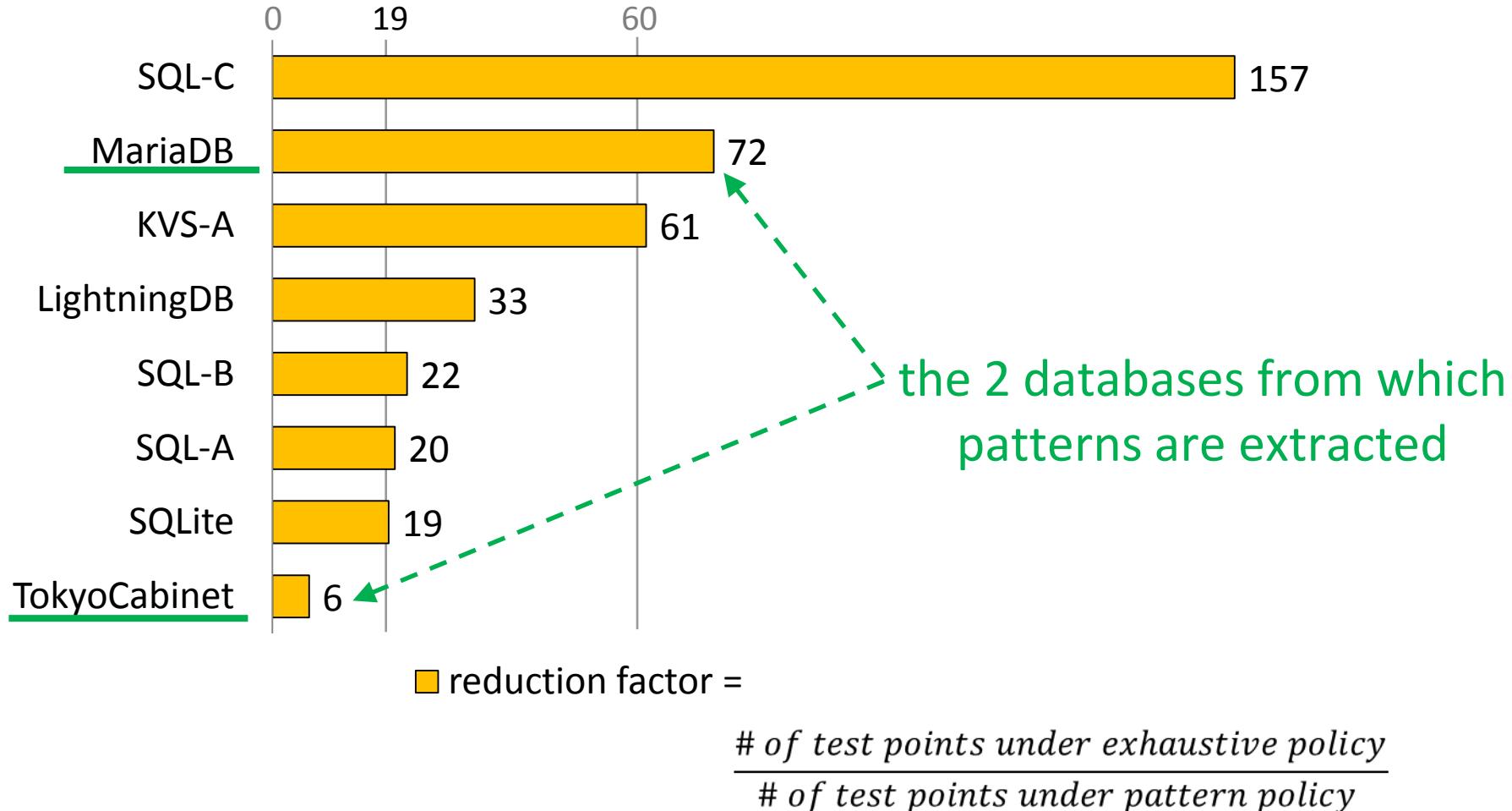
Patterns reduce required test points greatly while achieving similar coverage



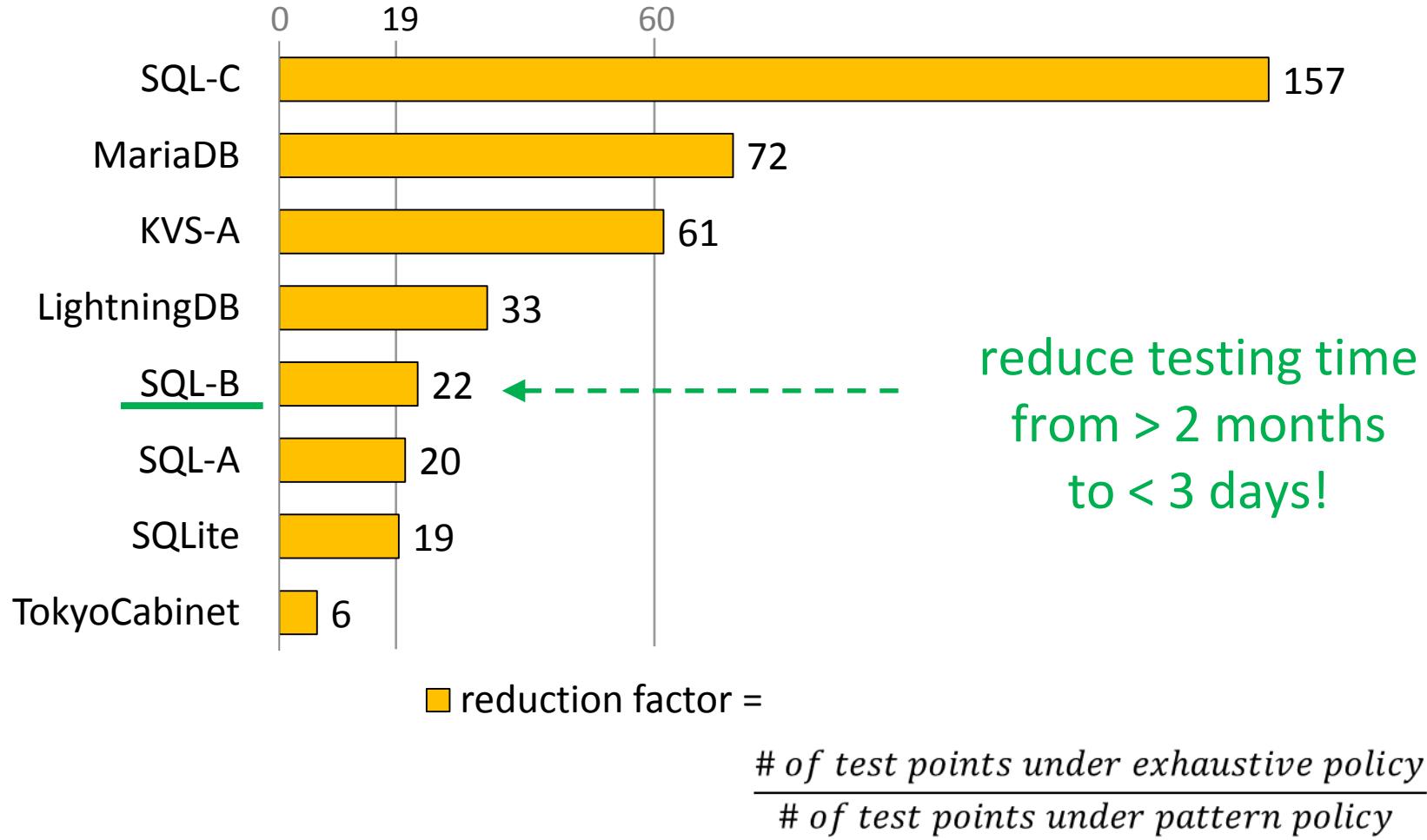
■ reduction factor =

$$\frac{\text{\# of test points under exhaustive policy}}{\text{\# of test points under pattern policy}}$$

Patterns reduce required test points greatly while achieving similar coverage



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Conclusions & Future Work

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 - traditional testing methodology may not be enough for today's complex storage systems
 - thorough testing requires purpose-built workloads and intelligent fault injection techniques

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 - traditional testing methodology may not be enough for today's complex storage systems
 - thorough testing requires purpose-built workloads and intelligent fault injection techniques
- Different layers in OS can help in different ways
 - iSCSI: fault injection w/ high portability & high fidelity
 - LBA & syscall: generic behavior patterns
 - combined multi-layer info: clear whole picture of complicated scenarios

Conclusions & Future Work

- We should bridge the gaps of understanding/assumptions!

between
DB & OS

between
User & DB

Conclusions & Future Work

- We should bridge the gaps of understanding/assumptions!

Pop Quiz: true or false?

between
DB & OS
between
User & DB

“mmap'ed files are not updated until msync()”

“file-length update are persistent after fdatasync()”

“durability is provided by the default configuration”

“transactions are durable after COMMIT”

Conclusions & Future Work

- We should bridge the gaps of understanding/assumptions!

Pop Quiz: true or false?

between DB & OS	“mmap’ed files are not updated until msync()”	false
between User & DB	“file-length update are persistent after fdatasync()”	
	“durability is provided by the default configuration”	
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Conclusions & Future Work

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Pop Quiz: true or false?

between DB & OS	“mmap’ed files are not updated until msync()”	false
between User & DB	“file-length update are persistent after fdatasync()”	depends!
	“durability is provided by the default configuration”	depends!
	“transactions are durable after COMMIT”	depends!

Conclusions & Future Work

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Pop Quiz: true or false?

between DB & OS	“mmap’ed files are not updated until msync()”	false
between User & DB	“file-length update are persistent after fdatasync()”	depends!
	“durability is provided by the default configuration”	depends!
	“transactions are durable after COMMIT”	depends!



Thank you!

