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### NVMeDirect: A User-Space I/O Framework for Application-specific Optimization on NVMe SSD

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### **Evolution of storage device**



### **Evolution of storage I/O stack**



### **Previous approaches** – Optimizing the kernel storage stack

- Use of polling for the fast I/O completion [Yang et al. FAST 2012]
- Optimization of a low-level hardware abstraction layer [Shin et al. ATC 2014]
  - Reducing the translation overhead between abstraction layers
- Optimizations to fully exploit the performance of fast storage devices [Yu et al. ACM TOCS 2014]
  - Polling, request merging, double buffering and reducing context switches

### **Limitations -** Optimizing the kernel storage stack

- Kernel should be general to provide an abstraction layer
- Kernel cannot implement any policy that favors a certain application
- Updating kernel requires a constant effort to port applicationspecific optimizations

### Previous approaches – Direct access to storage device

- Direct access to the special storage device [Caulfield et al. ASPLOS 2012]
  - Special hardware is required
- Direct access to NVMe device
  - Intel Storage Performance Development Kit SPDK (Sep 2015)
  - Micron Userspace NVMe driver project UNVMe (Feb 2016)
  - Device dedicated to a single user process
  - Provides just simple read & write interface based on polling
  - Not sufficient to port existing applications

### **Our Approach: NVMeDirect**



### **NVMeDirect Overview**

- Allows user-space applications to directly access NVMe SSDs without any hardware modifications
  - Achieves high performance by avoiding storage stack overhead
- Supports various I/O policies
  - Applications can be optimized according to their I/O characteristics
  - Selective use of block cache, I/O scheduler, or I/O completion thread
  - Asynchronous I/O vs. Synchronous I/O
  - Buffered I/O vs. Direct I/O
- Designed to maximize performance for trusted applications
  - Storage appliance, private clouds, etc.

### **NVMeDirect Design**



• NVMeDirect Management

- Kernel driver
- Admin tool

#### NVMeDirect I/O

- I/O Handles
- User-space I/O Queues

#### • NVMeDirect I/O Framework

- Block Cache
- I/O Scheduler
- I/O Completion Thread

### **NVMeDirect Design – Queues and Handles**

#### User-space I/O Queues

- Memory-mapped address space for NVMe I/O Queues created in the kernel address space
- I/O Handles
  - Used to send I/O requests to NVMe I/O Queue(s)
  - A thread can create one or more I/O Handles
  - Each Handle can be configured to use different features : caching, I/O scheduling, I/O completion, etc.



### **NVMeDirect Design - APIs**

Driver Connection	<pre>struct nvmed* nvmed_open(char*); int nvmed_close(struct nvmed*);</pre>
Queue Management	<pre>struct nvmed_queue* nvmed_create_queue(struct nvmed*); int nvmed_destroy_queue(struct nvmed_queue*);</pre>
I/O Handle Management	<pre>struct nvmed_handle* nvmed_create_handle(struct nvmed_queue*); struct nvmed_handle* nvmed_create_mq_handle(struct nvmed_queue**); int nvmed_destroy_handle(struct nvmed_handle*); int nvmed_set_param(struct nvmed_handle*, int, int);</pre>
Buffer Management	<pre>void* nvmed_get_buffer(struct nvmed_handle*, unsigned int); int nvmed_put_buffer(void*);</pre>
<b>I/O</b>	<pre>off_t nvmed_lseek(struct nvmed_handle*, off_t, int); ssize_t nvmed_read(struct nvmed_handle*, void*, unsigned int); ssize_t nvmed_write(struct nvmed_handle*, void*, unsigned int); int nvmed_flush(struct nvmed_handle*); int nvmed_discard(struct nvmed_handle*, off_t, size_t);</pre>

# Example of I/O using NVMeDirect



#### 1) Open device

nvmed = nvmed\_open("/proc/nvme0/n1");

#### 2) Create queue

queue = nvmed\_queue\_create(nvmed);

#### 3) Create handle

handle = nvmed\_handle\_create(queue);

#### 4) Perform I/O

size = nvmed read(handle, buf, len);

#### 5) Configure Handle

ret = nvmed set param (handle, BUFFERED IO, TRUE);

# **Advantages of NVMeDirect**

- Enables high performance I/O
  - Low latency and high throughput
- Easy to support new interfaces
  - Weighted queue, multi-stream, etc.
- Easy to develop and debug
- Provides various I/O policies
- Free from kernel update
- Co-exists with legacy kernel I/O

# **Evaluation**

- Implementation on the Linux kernel 4.3.3
- Experimental setup
  - Ubuntu 14.04 LTS
  - 3.3GHz Intel Core i7 CPU (6 cores) & 64GB of DRAM
  - Intel 750 Series 400GB NVMe SSD
- Comparison with
  - Kernel I/O
  - SPDK
  - NVMeDirect

### **Baseline performance**

• Asynchronous random I/O performance using FIO



# Impact of the Polling Period

- Polling is not efficient on bandwidth sensitive workload due to the significant increase in the CPU load
- Significant performance degradation occurs in a certain polling period



Polling Period (µs)

 Control Polling Period dynamically based on I/O size or hints from applications

### **Latency-sensitive Application**

- Redis: in-memory data structure store
- Logging every operation for persistency
- Logs are 10 to 100 bytes in size
- Write buffer is required due to small-size data
  - Difficult to run on SPDK without significant code modification



# **Latency-sensitive Application**

- Using workload-A in YCSB on Redis
- Update-heavy workload with Zipf distribution



# **Differentiated I/O Service**

- NVMeDirect supports prioritized I/O without H/W features
  - Prioritized I/O without a weighted round-robin scheduler
  - Using flexible binding between Handles and Queues
  - Sharing a single Queue with multiple Handles



# **Differentiated I/O Service**

- One prioritized thread with a dedicated queue, Three threads with a shared queue
- Each thread performs 4KB random write



# Conclusion

- NVMeDirect
  - First full framework for
     I/O in the user-space based on stock NVMe devices
  - Can be easily applied to many applications
  - Useful for emerging storage devices, e.g. 3D XPoint<sup>™</sup>, etc.
  - Available as open-source at https://github.com/nvmedirect (July 2016)
- Future work
  - User-level file systems
  - Porting diverse data-intensive applications over NVMeDirect
  - Protecting the system from illegal access

# Thank you

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