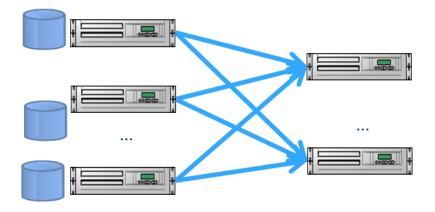


Rhea: Automatic Filtering for Unstructured Cloud Storage

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Cluster design for data analytics: [Traditional] Collocate storage & compute



Hadoop & MapReduce, Dryad/DryadLing, Scope, etc

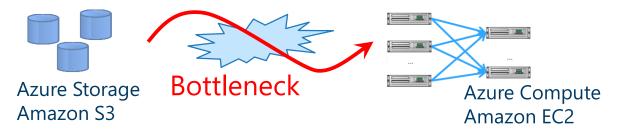
Cloud Analytics: Hadoop in the Cloud **Separate** storage and compute



Azure Storage Amazon S3 Azure Compute Amazon EC2

Examples:Hadoop on AzureAmazon's Elastic MapReduce

Cloud Analytics: Hadoop in the Cloud **Separate** storage and compute

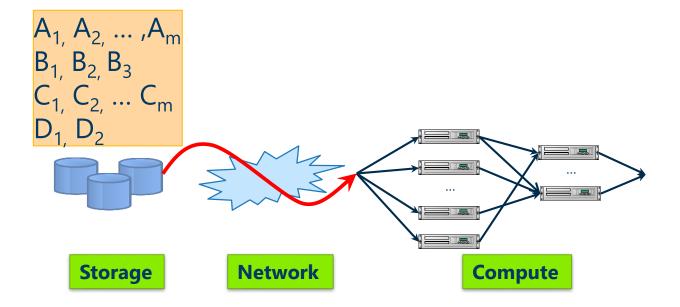


Why separate storage from compute?

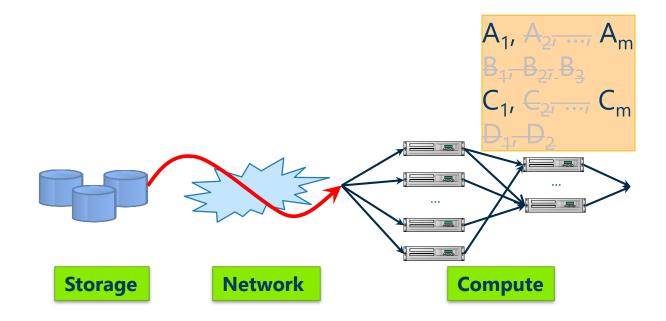
- + (User) Don't pay for compute just to keep data alive
- + (User) Offload storage management to operator
- + (Operator) Evolve compute & storage independently
- + (Operator) Offer services that do not require both

- Network between storage and compute is limited (see paper for details)

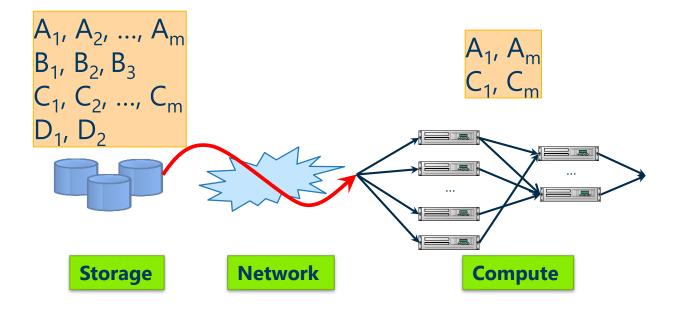
Problem: Transfer lots of data ...



Problem: Transfer lots of data even when only a subset is needed



Problem: Transfer lots of data even when only a subset is needed



Scenario

- Apache Hadoop (Map/Reduce)
- Input data in storage service
- Hadoop running in compute service
- Unstructured data:
 - \cdot text, log files, etc

Goal

Transparently reduce data transfers from storage to compute

How to minimize transfers?

- Strawman: Can we execute mappers on storage nodes?
 - \cdot Intuition: Mappers throw away a lot of data
 - Data reduction not guaranteed
 - Difficult to stop mappers during storage overload
 - Storage nodes have to execute complicated logic (Hadoop system & protocol)
 - Dependencies on runtime environment, libraries, etc
- Better approach: Filter unnecessary data at storage nodes
 - Filters need to be **opportunistic and transparent** i.e. can kill/restart at any time (e.g. during overload)
 - Filters need to be **correct** i.e. always preserve correctness of computation

Challenge: How to filter the data?

Recall: data are typically unstructured text • No external source of structure/schema

Insight:

- The data analytic job knows structure
- ... and what needs to be filtered

Idea: static analysis of job bytecode

public void map(... **value** ...)

Input Value

String[] entries = value.toString().split("\t"); *String* articleName = entries[0]; *String* pointType = entries[1]; String geoPoint = entries[2];

Projection operation

3 "columns" interesting (out of 4 for this job)

if (GEO_RSS_URI.equals(pointType)) {

StringTokenizer st = new StringTokenizer(geoPon String strLat = st.nextToken();

String strLong = st.nextToken();

double lat = Double.parseDouble(strLat);

double lang = Double.parseDouble "selects"/"projects"

} }

String locationName = geoLocationKey.set(locationKey);

geoLocationName.set(locationName);

outputCollector.collect(geoLocationKey, geoLocationName);

Output operation

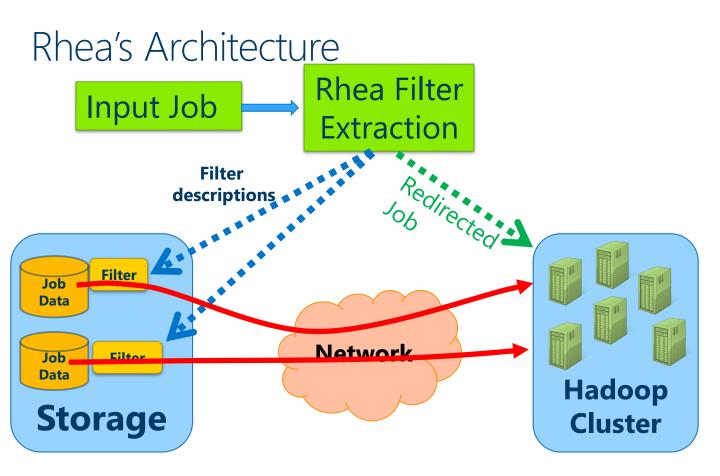
Selection operation

roughly 1/3 of rows are of the interesting type

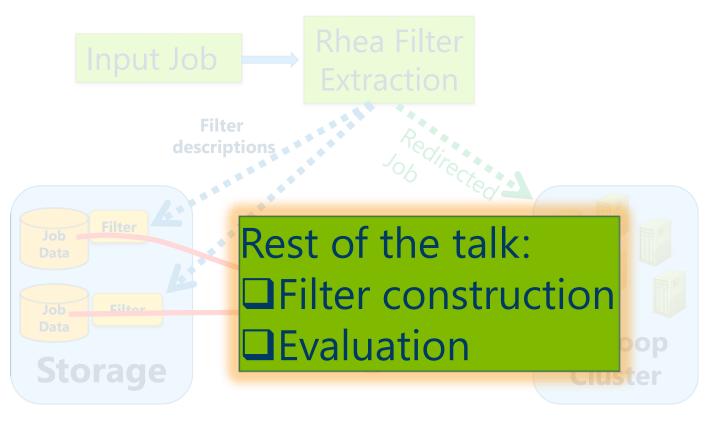
implicit in Java byte code

Rhea

- Static analysis of Java byte code
- Extract row (select) & column (project) filters
 - · as **executable** Java methods
 - · column filters can also be C, regular expressions, etc.
- Filters are **conservative**:
 - \cdot May accept more data than strictly necessary
- Filters are **opportunistic**
 - \cdot kill/restart at any time (e.g. during storage overload)
- Filters are transparent
 - \cdot no change to Hadoop job



Rhea's Architecture



Filters: Identify bits of data that affect output of mapper

- Row Filters:
 - \cdot Given an input row:

Does it lead to output?

- \cdot Row corresponds to one invocation of map
- · Approach: Path Slicing
- · Challenge: Deal with mutable state

• Column Filters:

- Given a row that leads to output:
 Which substrings of the row affect output?
- · Approach: Abstract interpretation
- · Challenge: Deal with loops

Row Filter Generation via Path Slicing

public void map(... value ...)

```
String[] entries = value.toString().split("\t");
String articleName = entries[0];
String pointType = entries[1];
String geoPoint = entries[2];
```

if (GEO_RSS_URI.equals(pointType)) {

StringTokenizer st = new StringTokenizer(geoPoint, " "); String strLat = st.nextToken(); String strLong = st.nextToken(); double lat = Double.parseDouble(strLat); double lang = Double.parseDouble(strLong); String locationKey = String locationName = geoLocationName = geoLocationName.set(locationKey); geoLocationName.set(locationName); outputCollector.collect(geoLocationKey, geoLocationName);

- 1. Tag "observable" instructions
- 2. Identify path conditions that lead to observable instructions
- Perform dataflow analysis to identify all instructions that affect path conditions
- 4. Emit code



public boolean filter(Text bcvar2) {
 String[] bcvar5 = bcvar2.toString().split("\t");
 String bcvar7 = bcvar5[1];
 boolean irvar0_1 =
GEO_RSS_URI.equals(bcvar7);
 if (irvar0_1 == 1) { return true; }
 return false;
}

Challenge: Taming State

- Map-Reduce program are often NOT pure functions \rightarrow M/R programmers use state (i.e. objects in heap):
 - \cdot ... to avoid frequent initializations
 - · ... to pass job parameters
 - · ... to optimize temporary storage (e.g. with dictionaries)
- Filters cannot rely on mutable state:
 - · Recall: output of filtered data = output of original data
- Solution: Tag all access to mutable fields as "observable" (i.e. output) instructions.

Column Filter Generation (aka projects)

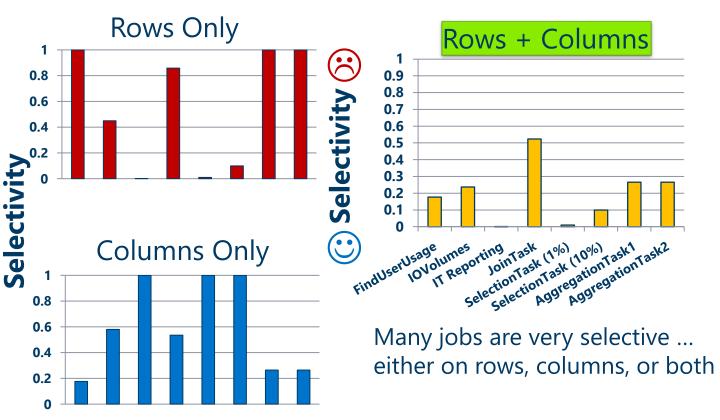
Goal: Identify substrings that affect output

- Based on *abstract interpretation*
 - Captures common patterns for "reading" fields: e.g. string tokenizers, regular expressions, etc.
 - \cdot Guarantees termination by using numerical constraints
 - \cdot Important to deal with loops
- Output:
 Tokenization method and sep
 List of indices of interesting to
 Evaluation

Experimental setup

- •Hadoop on Azure:
- Input data in Azure Storage
- •Compute on Azure Compute
- •8 jobs with both code and data
 - ·200 jobs code only (in paper)
- Same data-center
 - ·Also, cross data-center (in paper)

Job Selectivity



Job Selectivity

- Rows + Columns 0.9 High selectivity \rightarrow 0.8 0.7 Selectivit less bytes to transfer 0.6 0.5 ✓ Good for operators 0.4 0.3 ✓ Cheaper for users 0.2 0.1 for cross-data centers Selection Task (10%)-AggregationTask1 .-yererationTask2 $\overline{\bigcirc}$ selectionTask (1%) FindUserUsage IOVolumes IT Reporting scenarios [see paper]
- reduce runtime ✓ Good for users

Many jobs are very selective ... either on rows, columns, or both

Measuring runtime benefits

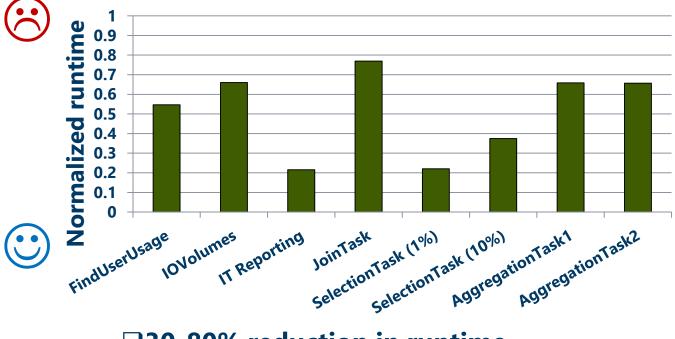
 We cannot extend Azure Storage or Amazon S3 with filters ☺

- Instead, we use pre-filtered data
 and compare with unfiltered data
- We assume storage with: (a) scalable I/O, and
 (b) enough processing power for filtering

Diversion: Do we have enough processing power?

- Row & Column filtering in Java: ~100MBytes/sec per core
 Scales linearly with multiple cores
- $\cdot \leq 2$ cores for filtering enough for all but 1 job
- Runtime always reduces runtime, even with fewer cores
- Performance dominated by string input/output, not filter
- Column filtering in optimized C: 5-17x faster than Java

Runtime benefits



 30-80% reduction in runtime
 Runtime reductions less than selectivity due to Hadoop overheads

Conclusions

- Hadoop in the cloud: separation of storage and compute.
- Rhea minimizes transfers from storage to compute
 - \cdot Uses static analysis on the job bytecode
 - Extracts **selection** and **projection** operators from code
 - · Generates filters to run in the storage layer
 - · Runs transparently to user (and is safe for provider)
 - Potential benefits to the user (time, money) and cloud provider (bandwidth)



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