



USENIX ATC'22 Best Paper Award!



Co-opting Linux Processes for High-Performance Network Simulation

Rob Jansen, U.S. Naval Research Laboratory

Jim Newsome, Tor Project

Ryan Wails, Georgetown University & U.S. Naval Research Laboratory

Rob Jansen, Ph.D.

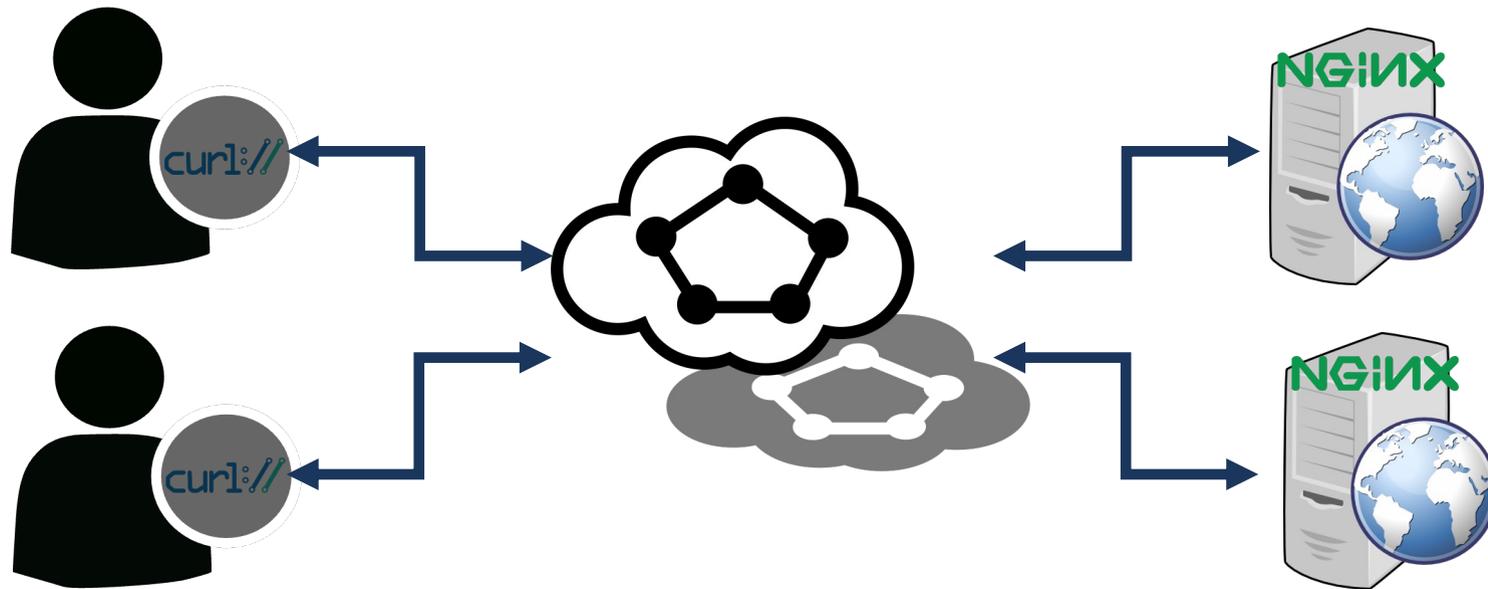
Computer Security Research Scientist
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USENIX Annual Technical Conference
Carlsbad, CA, USA
July 11th, 2022



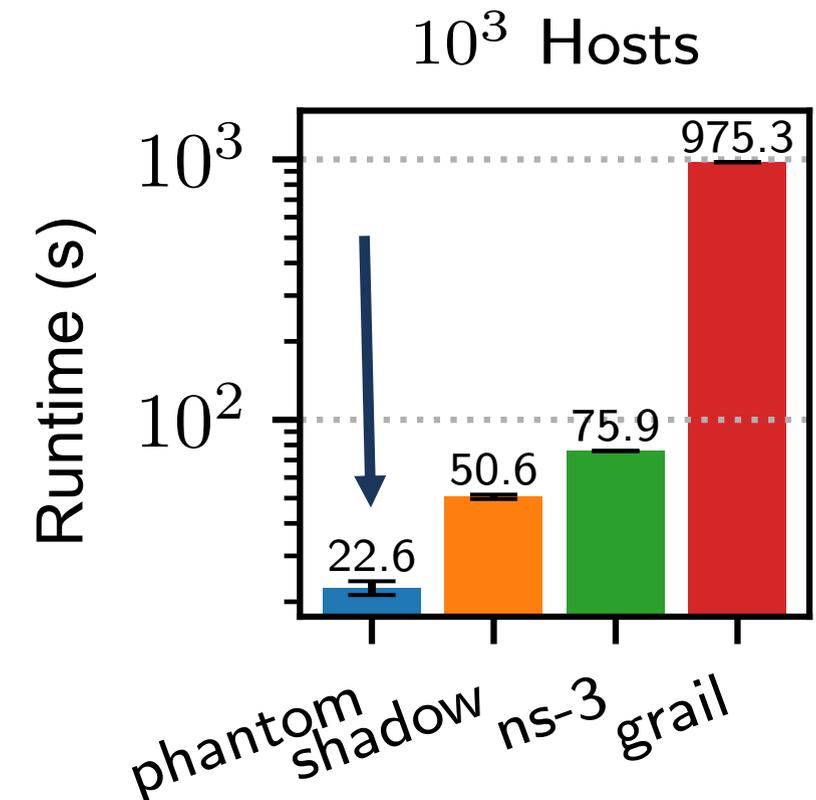
Designed a new, hybrid network simulator/emulator

- co-opts Linux processes into a discrete-event network simulation that emulates kernel functionality
- enables large-scale, distributed system experiments



- Merged into the open-source Shadow project and synonymous with **Shadow v2**
- Artifacts: <https://netsim-atc2022.github.io>

- **2.3x** faster than Shadow v1
- **3.4x** faster than NS-3
- **43x** faster than gRaIL [ToN'19]



Outline

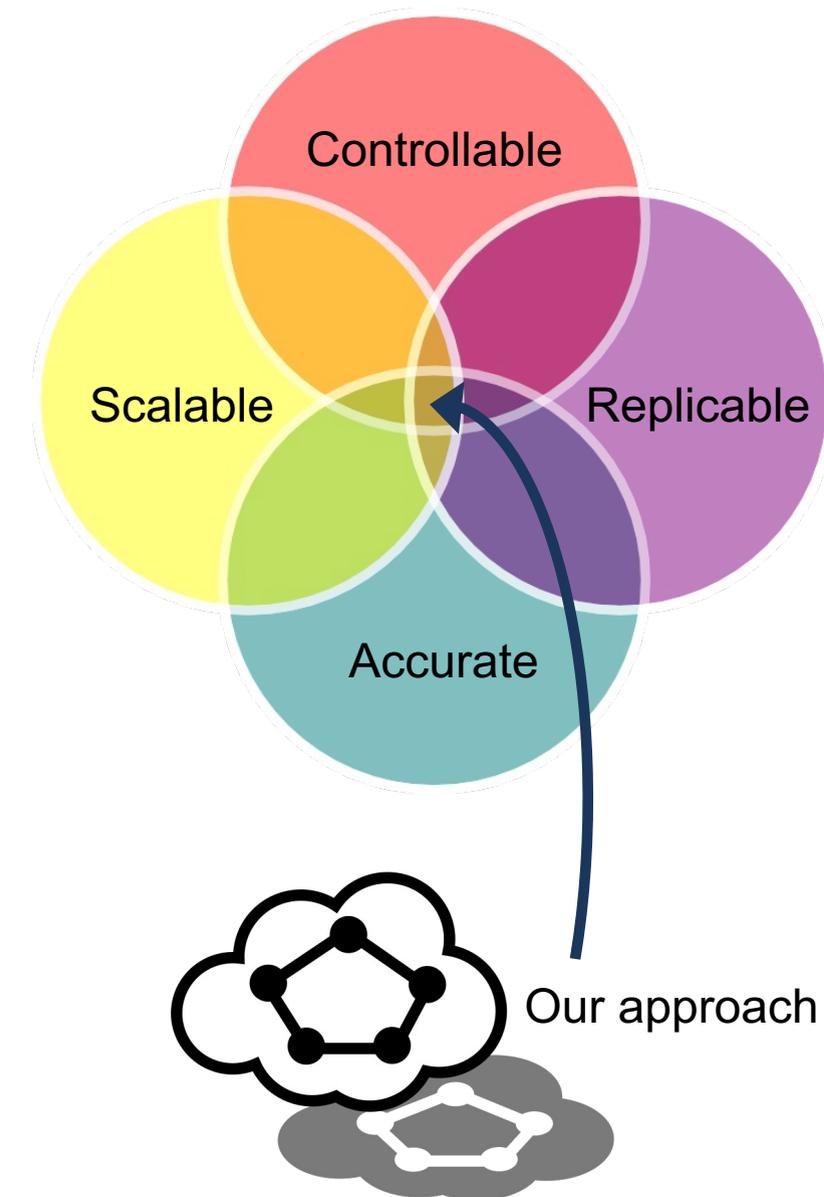
motivation

design

evaluation

Requirements for Large Distributed System Experimentation

- Important properties of test networks:
 - **Controllable**: isolate important factors
 - **Replicable**: identically replicate experiments (determinism)
- Requirements for large distributed systems:
 - **Accurate**: directly execute system software (not an abstraction)
 - **Scalable**: decouple from time, computational constraints of host



Problems with Traditional Approaches

- Simulation (e.g., ns-3)
 - Not realistic: runs abstractions instead of real applications
 - Hard to maintain and can lead to invalid results
- Emulation (e.g., mininet)
 - Not controllable: results will not be identical
 - Not scalable: CPU overload → time distortion



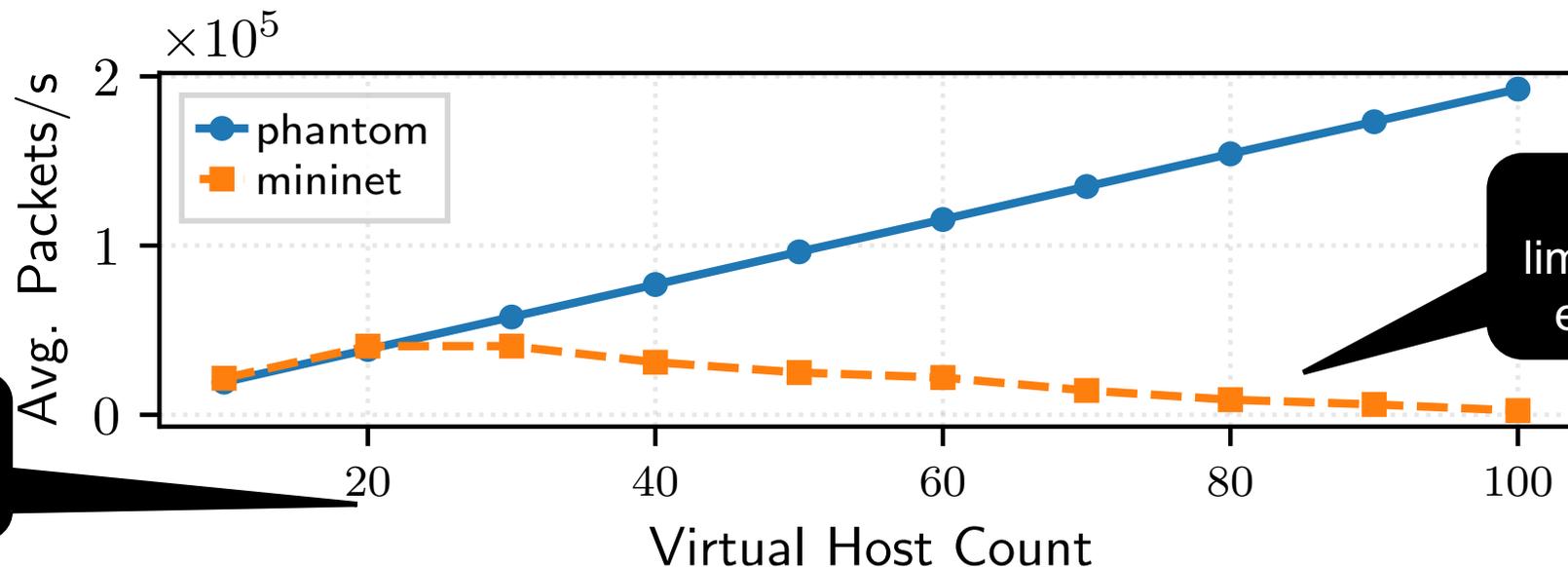
Mininet > sudo mn

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Mininet > sudo mn



CPU overload with >20 hosts

Forwarding capacity limited: fewer packets than expected are forwarded

Hybrid Architectures and Challenges

- Hybrid architecture
 - Network simulation, but directly execute application code
 - Enjoys advantages of both simulation and emulation
 - Best opportunity to scale to large-scale distributed systems

Architecture	Example Tool	Scalability	Realism	Control
Emulation	Mininet	✗	✓	✗
Simulation	NS-3	✓	✗	✓
→ Hybrid	This work	✓	✓	✓

Executing application code via plugin (link-map) namespaces

- `appid = dlmopen(app.so)`
- `func = dlsym(appid, "main")`
- `func()`

NS-3-DCE, Shadow

Limitations

- Compatibility (must build PIC/PIE)
- Correctness (intercept libcalls only)
- Maintainability (custom ld, threading)

Limitations of Hybrid Architectures (1)

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NS-3-DCE, Shadow

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Table 2: Application Properties Supported in Hybrid Simulators

Application Property	Shadow	Phantom
Multiple threads (e.g., support for pthreads)	●	●
Multiple processes (e.g., support for fork)	◐	◐
Not position-independent (i.e., PIC or PIE)	○	●
Not dynamically linked to libc	○	●
Symbols not exported to dynamic symbol table	○	●
System calls made in statically linked code	○	●
System calls made in assembly (i.e., avoiding libc)	○	●
100% statically linked (e.g., some go programs)	○	◐

- Does not work in tool or architecture ● Works in tool & architecture
 ◐ Not implemented in tool (as of writing) but supported by architecture

Limitations of Hybrid Architectures (2)

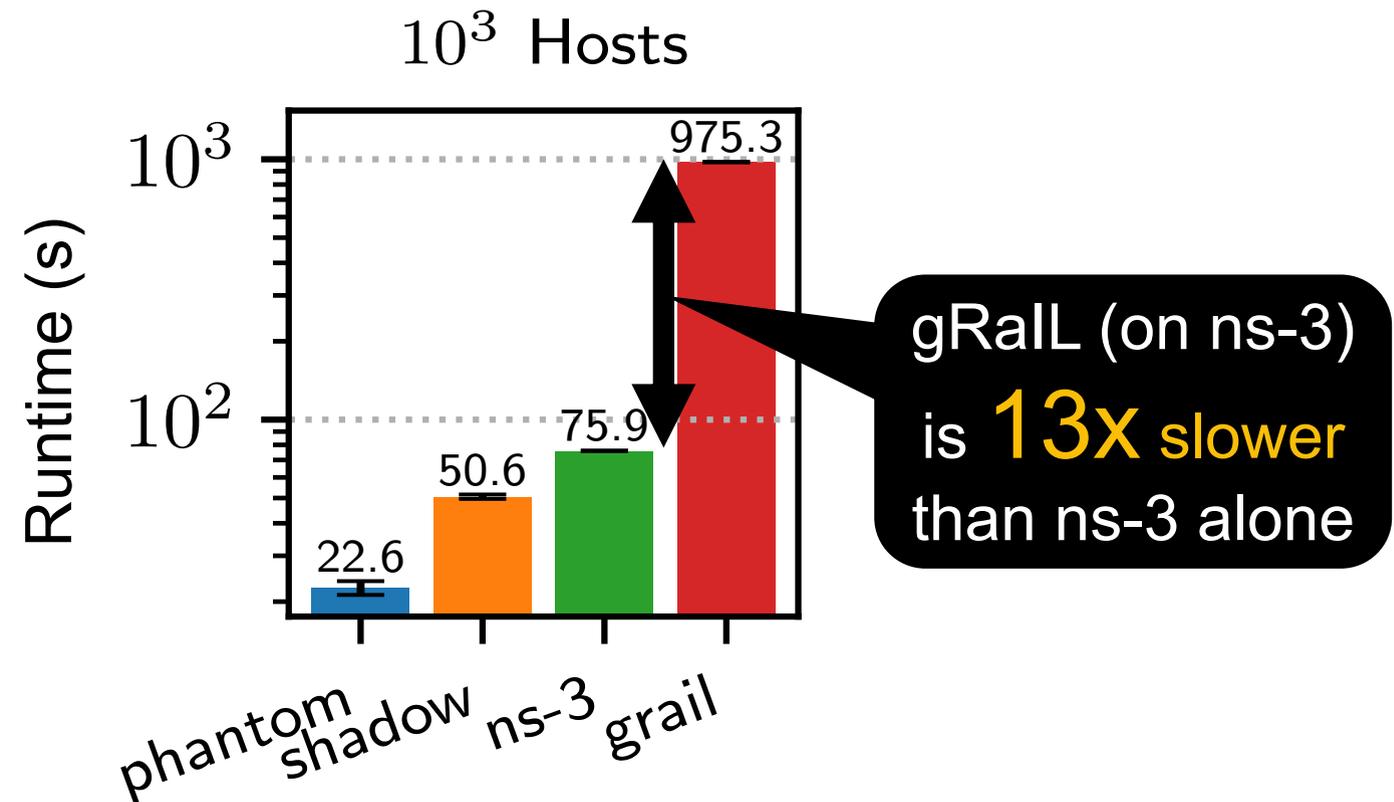
Executing code via Linux processes

- `fork() + exec() → ptrace()`

gRaIL [ToN'19]

Limitations: ptrace is slow!

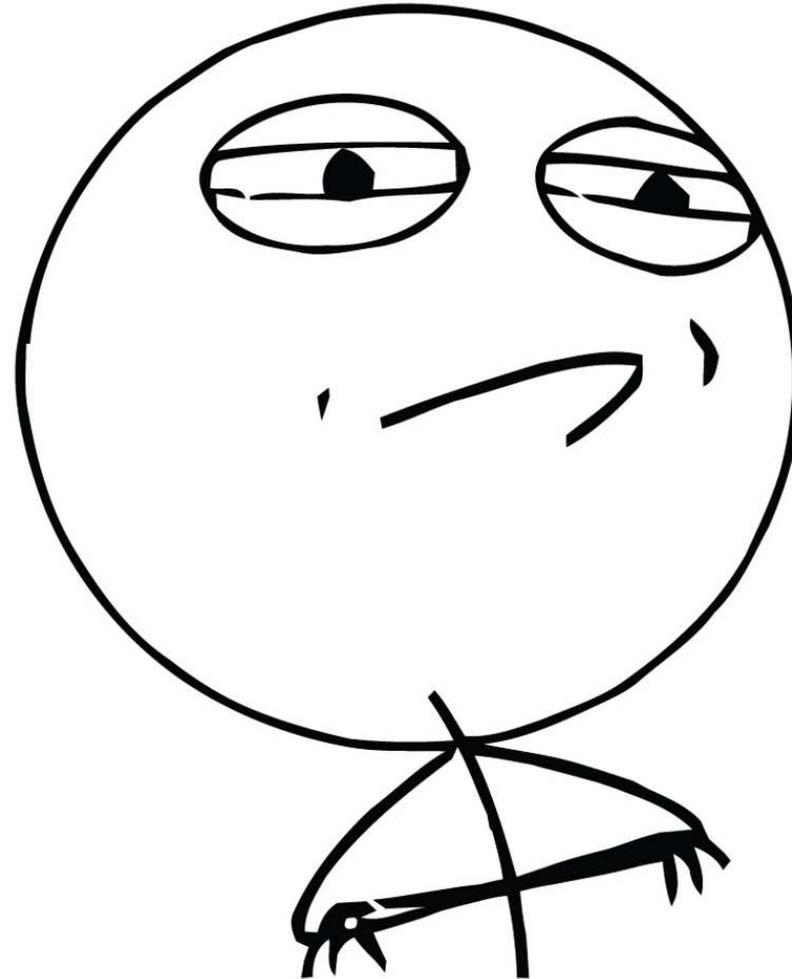
- Process control: overhead **quadratic** in total number of processes
- Syscall interception: 4 context switches **for every syscall**
- Data transfer: extra syscall + mode change **for every word of memory**



Our Research Challenge

Can we design a tool with the
performance benefits of a uni-process
plugin-based architecture
AND
the improved **modularity** and **isolation**
of a mutli-process architecture?

Our Research Challenge



CHALLENGE ACCEPTED

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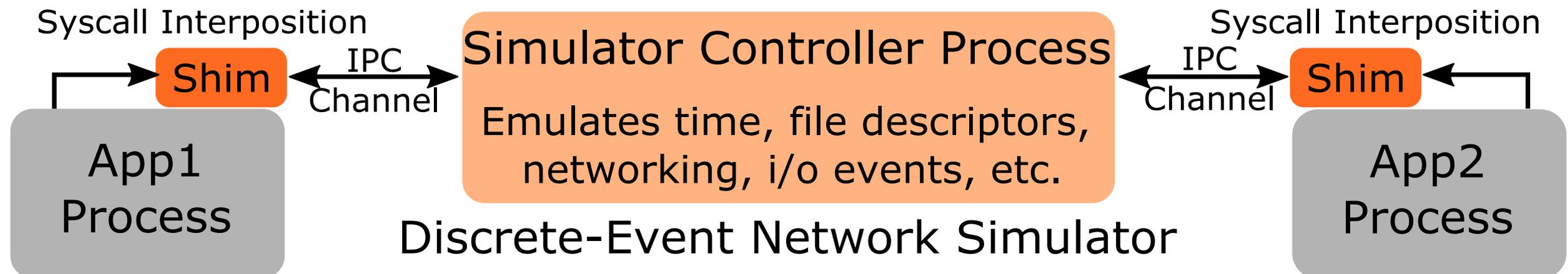
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- Discrete-event packet-level network simulator
- Directly executes apps as standard Linux processes
- Intercepts all system calls made by apps and emulates them

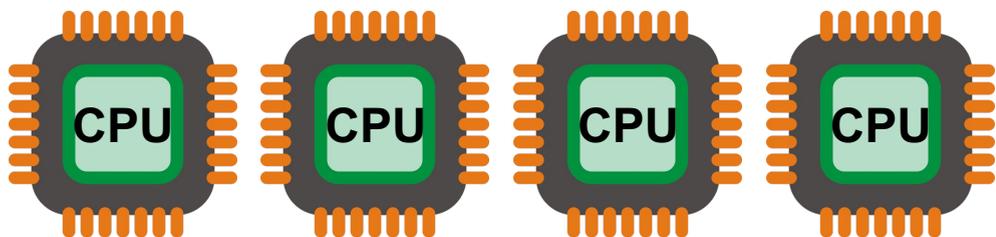
- Simulates system call behavior and networking
 - File descriptors (files, sockets, pipes)
 - Event notification (poll, epoll, select)
 - Networking (buffers, protocols, ifaces)
 - DNS and routing (latency, bandwidth)



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Goal: efficiently parallelize simulation workload

Sim Controller Process

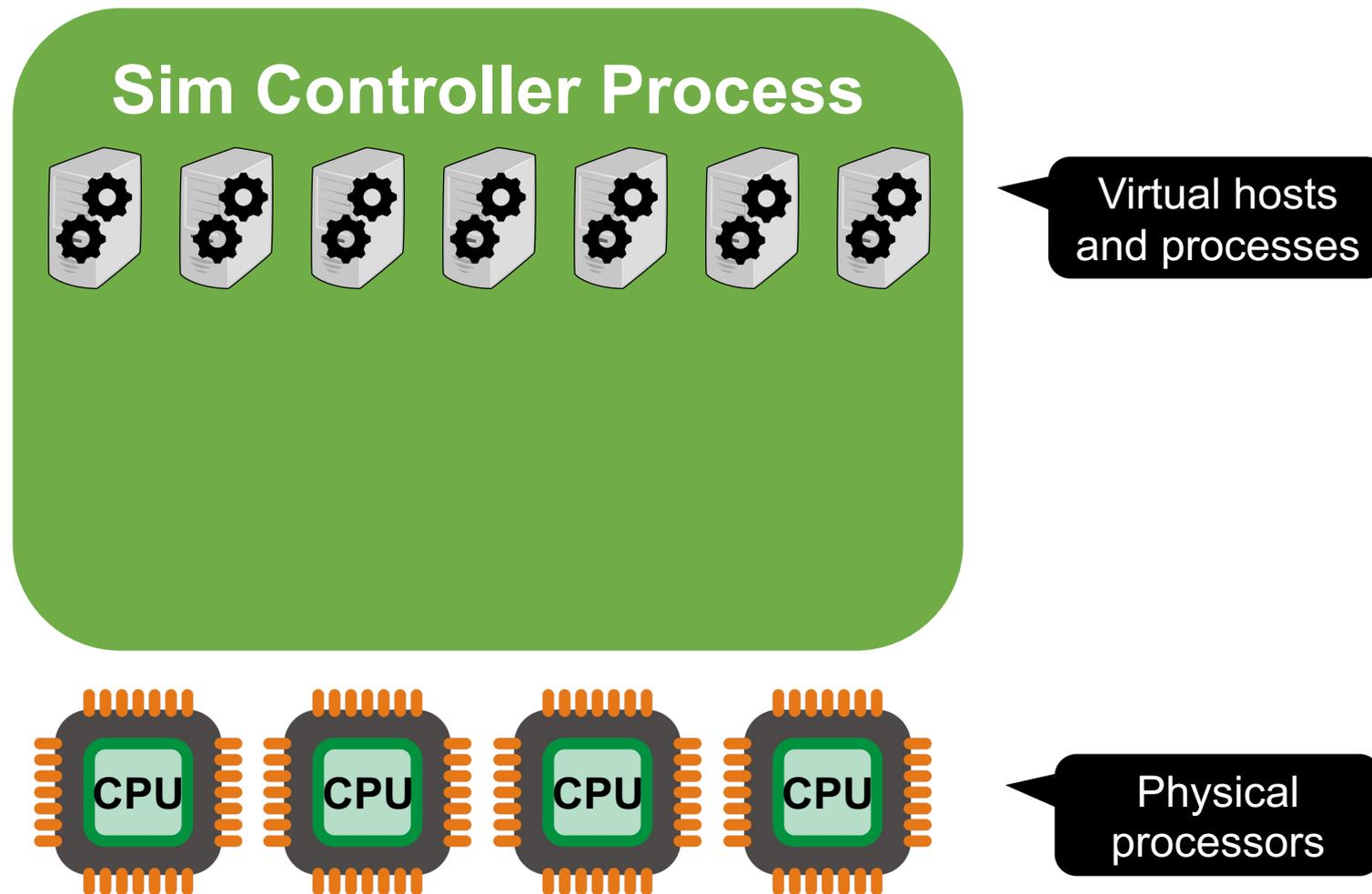


Physical
processors

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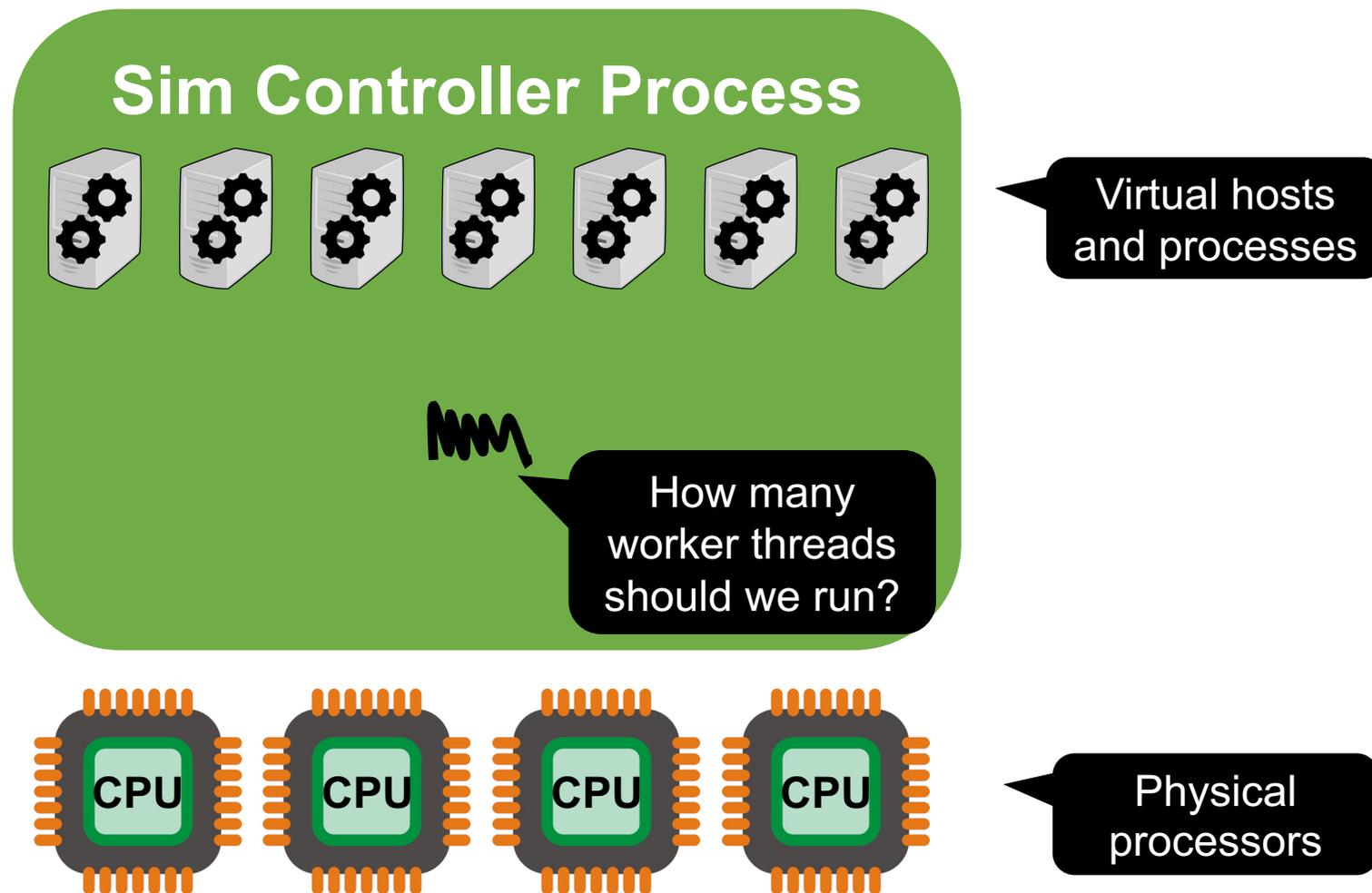
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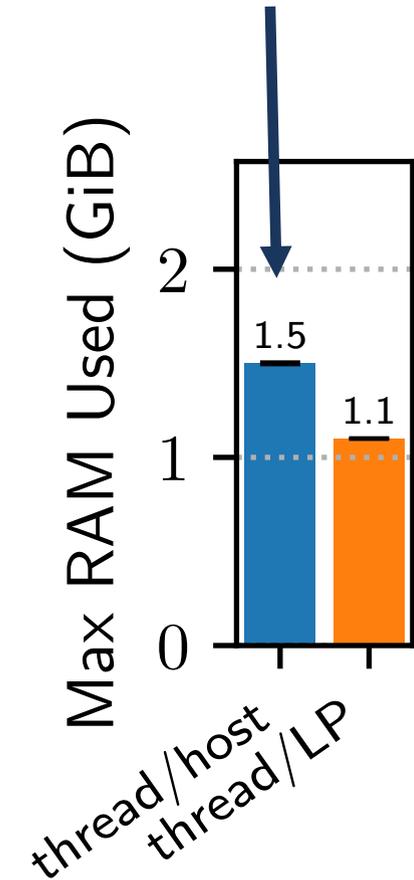
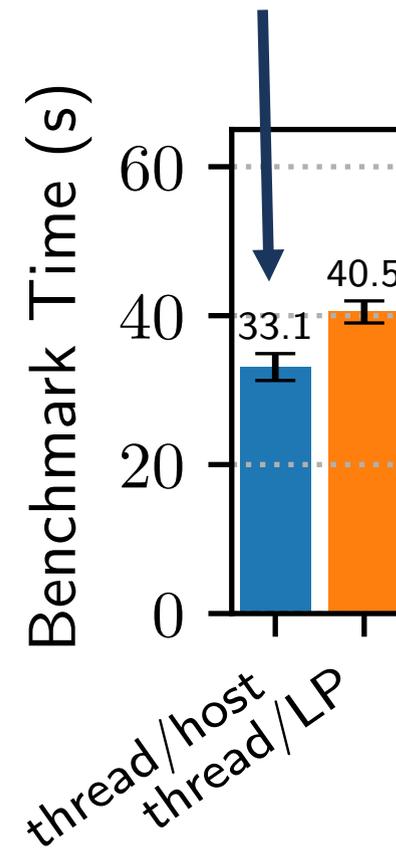
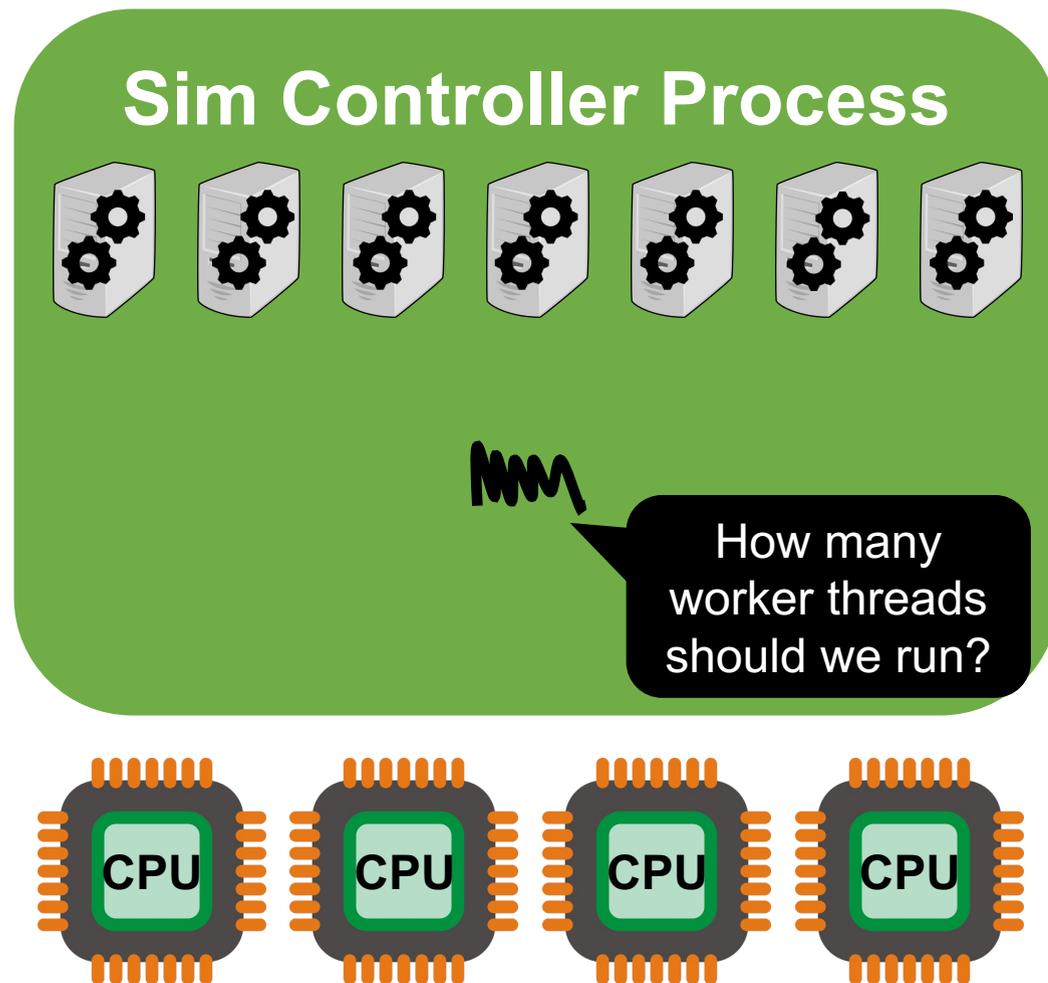
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Parallel Worker Threads

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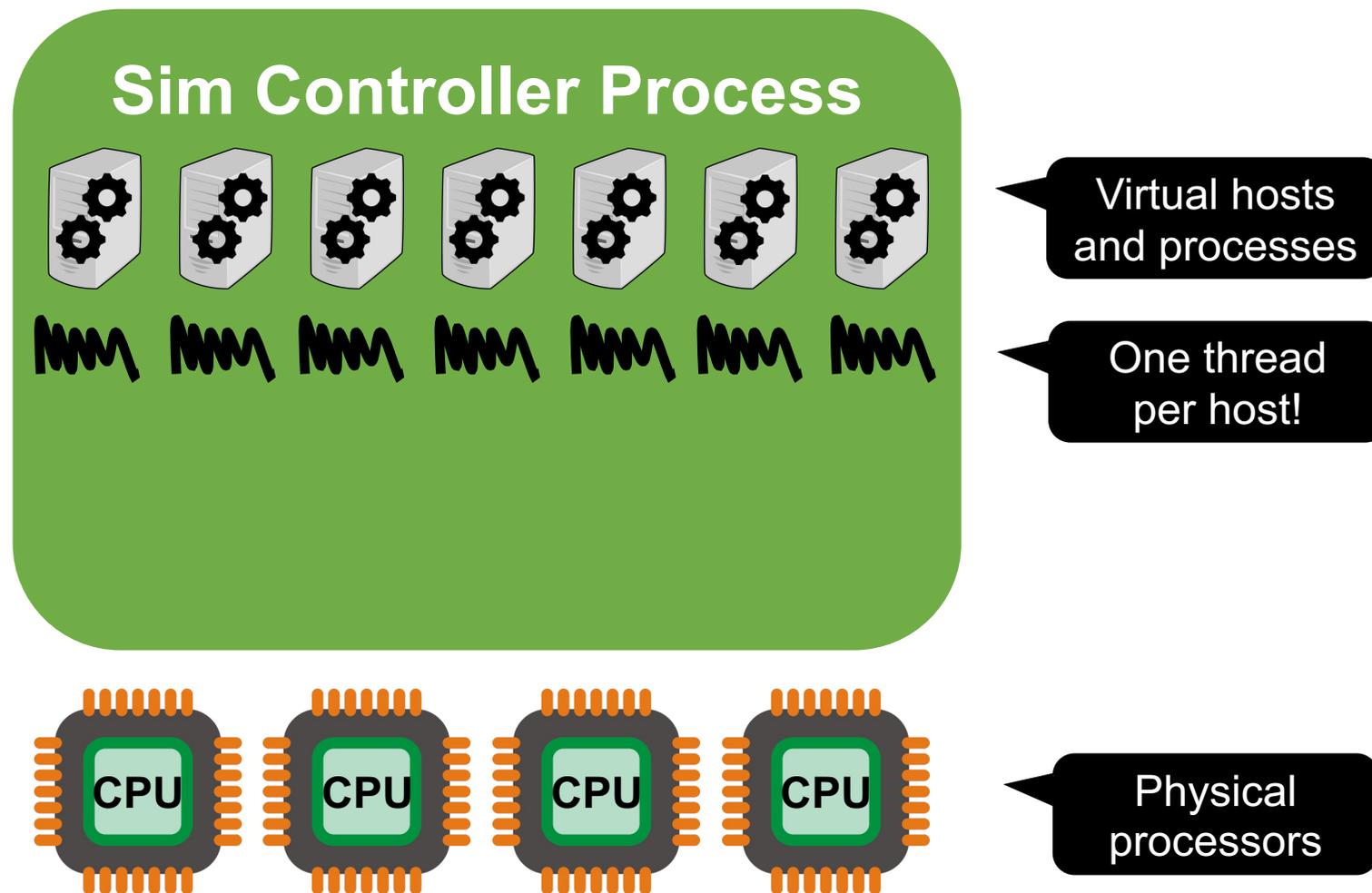
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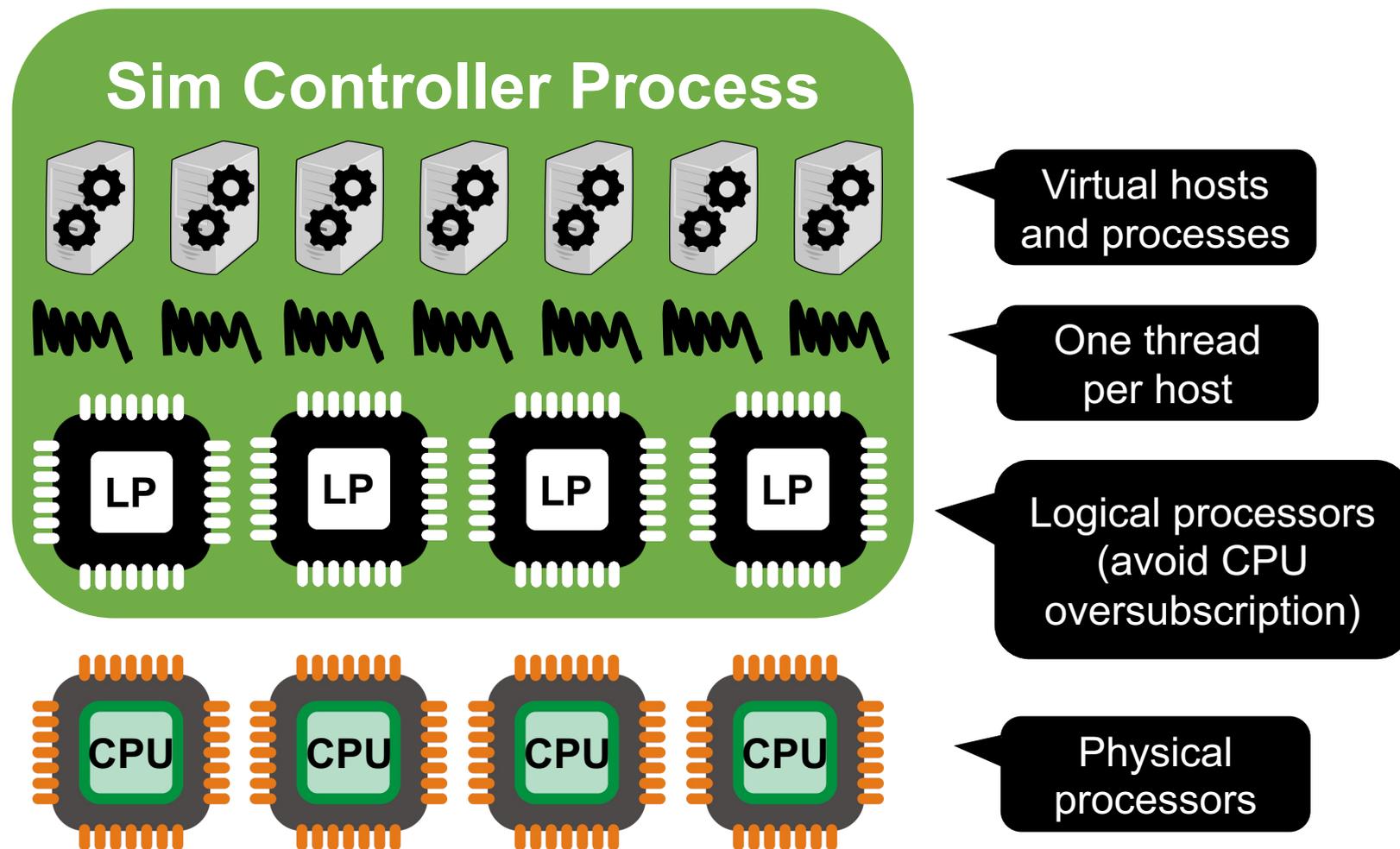
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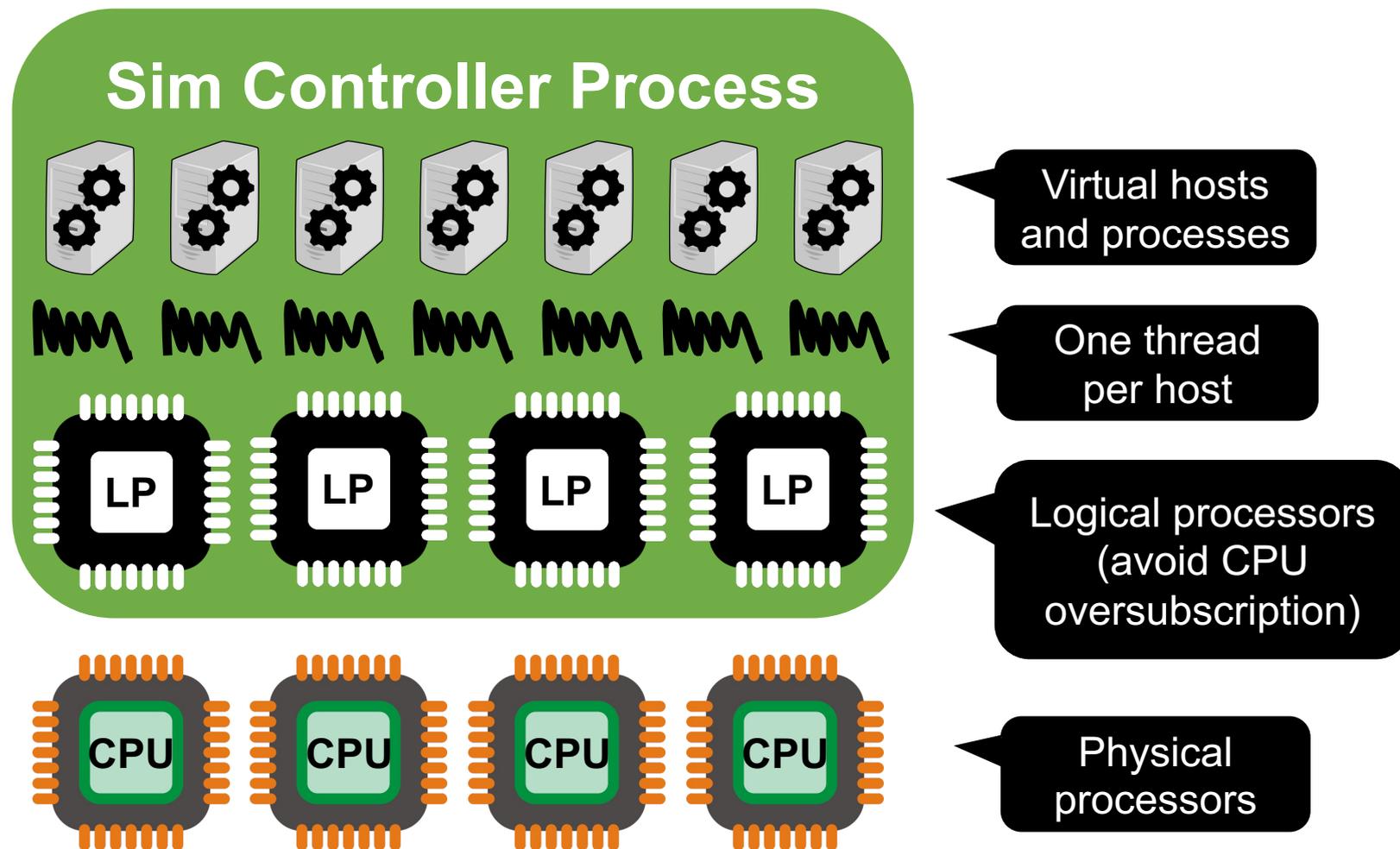
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Goal: efficiently parallelize simulation workload

Thread scheduling:



- Work stealing
- Each LP starts a thread
 1. Runs all assigned events in current round (1 ms)
 2. Set thread to *waiting*
 3. Starts next *waiting* thread (if any)
- When all threads *waiting*
 - Advance round clock
 - Repeat

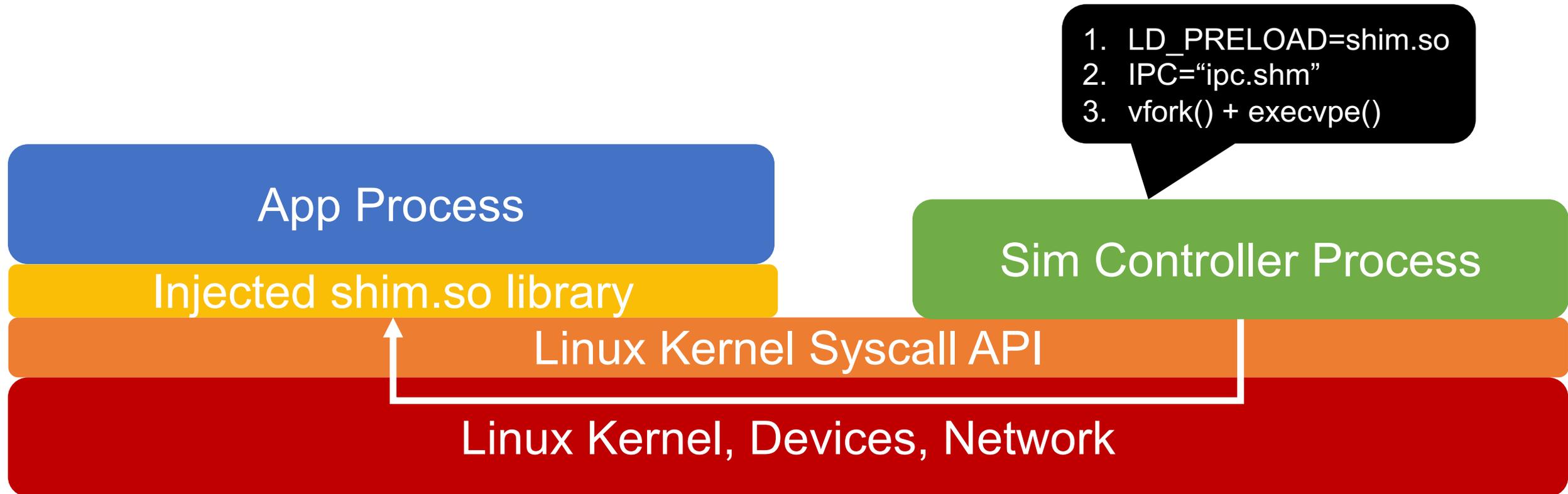
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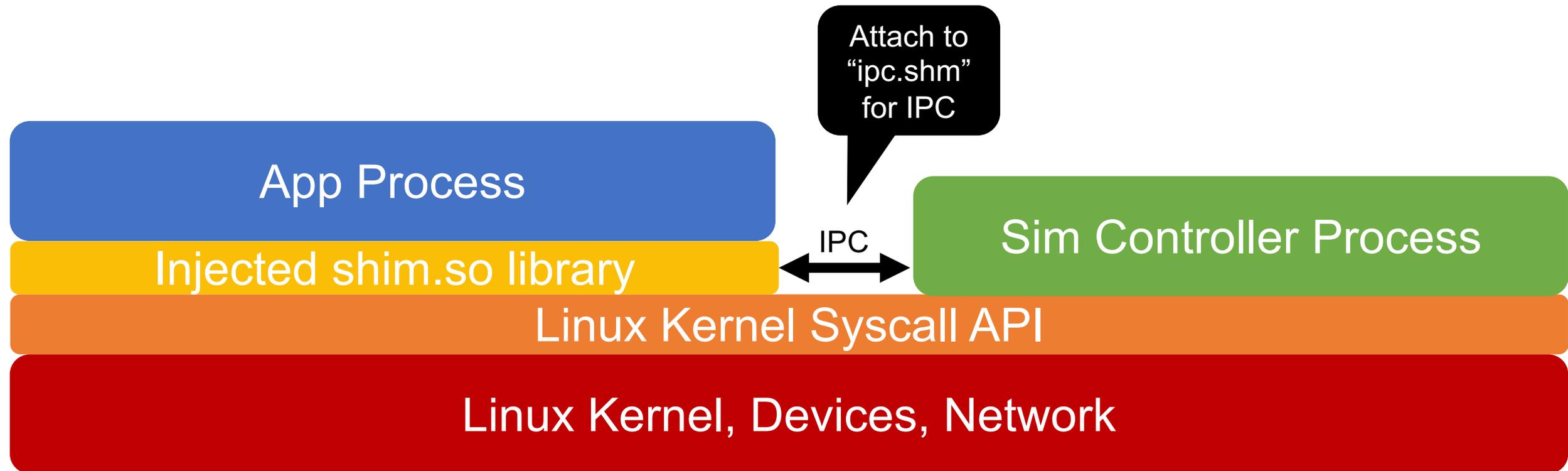
Linux Kernel Syscall API

Linux Kernel, Devices, Network

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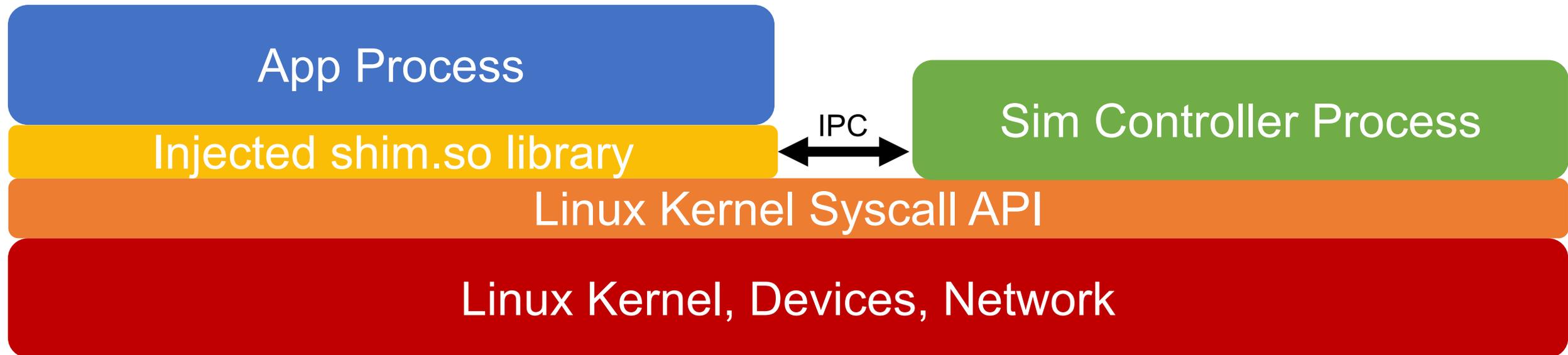


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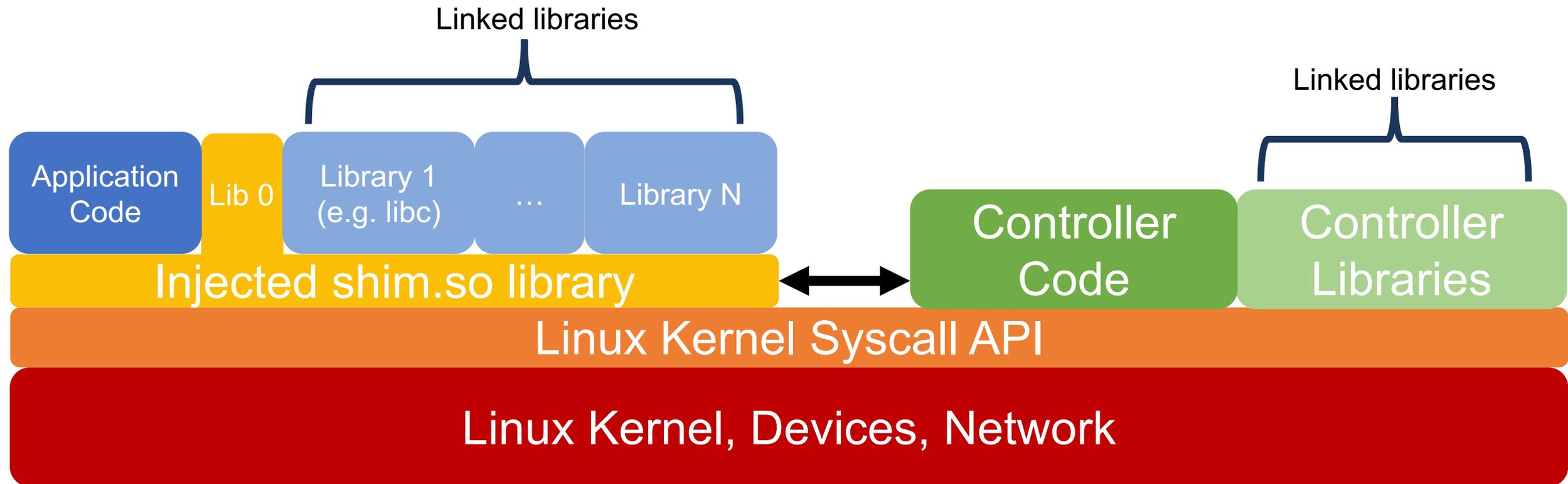
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App is running and can communicate with controller



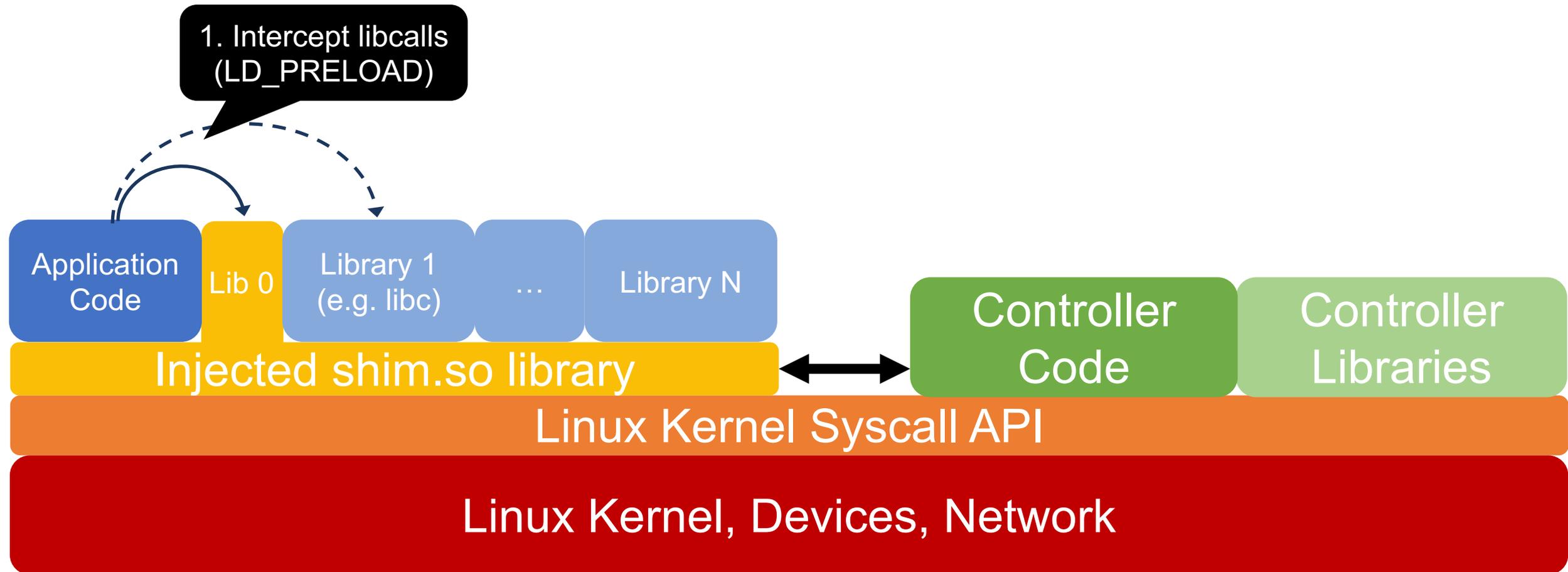
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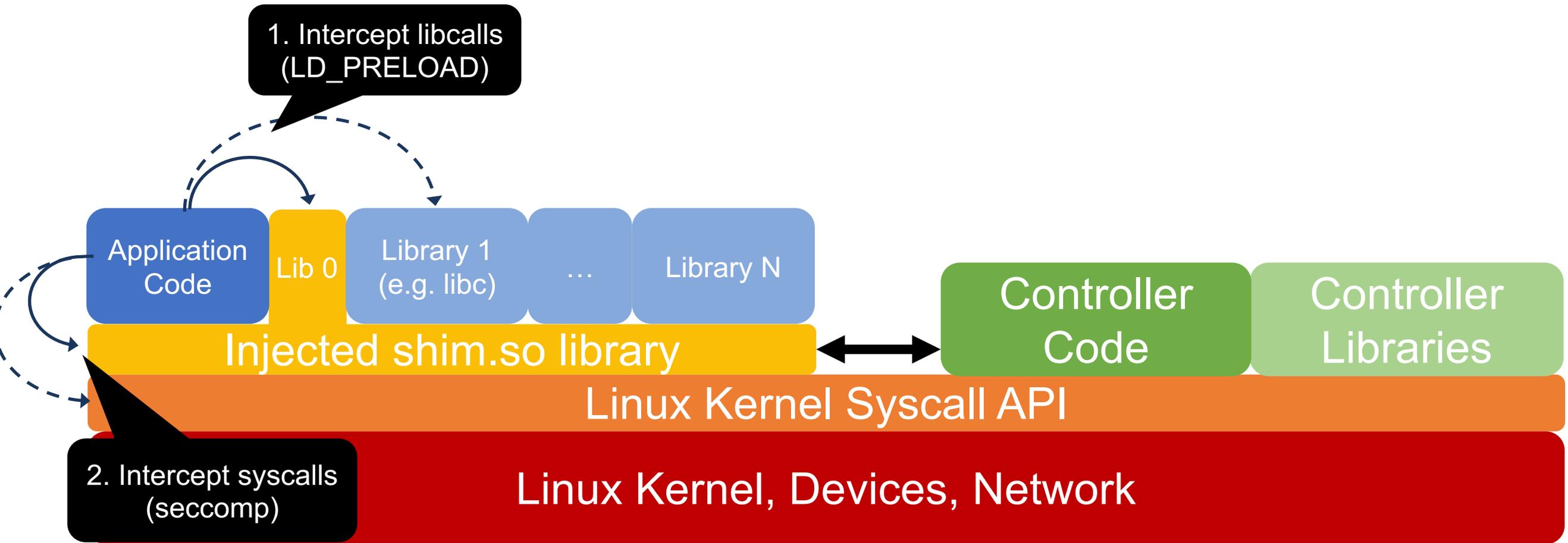
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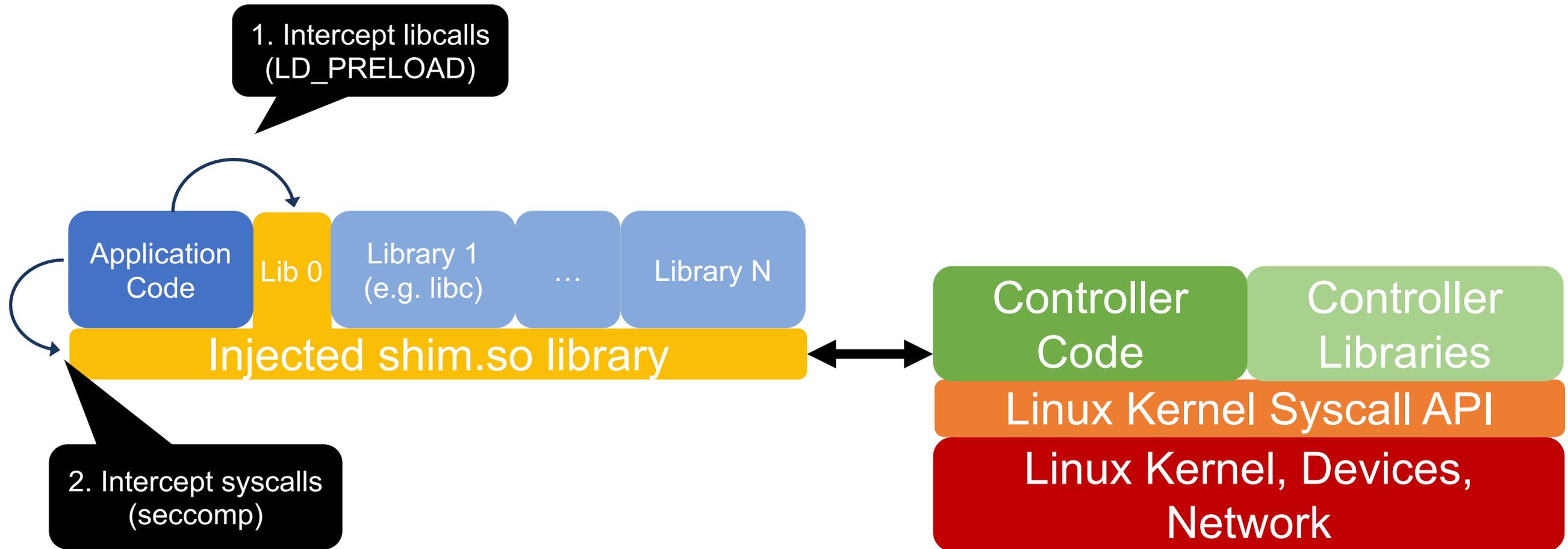
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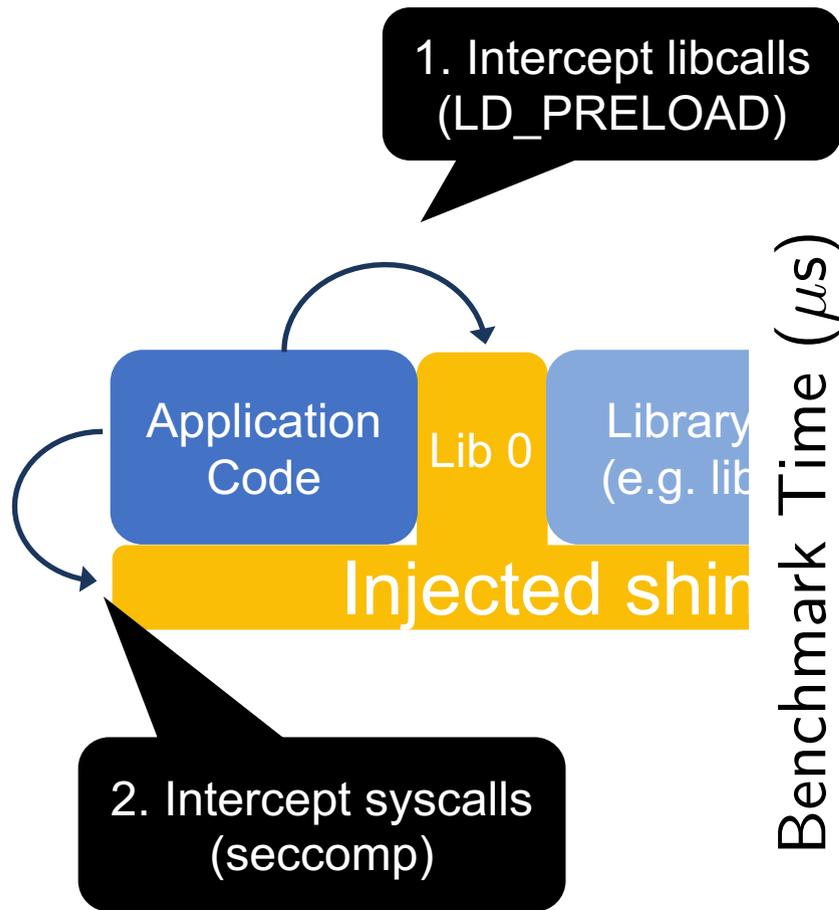
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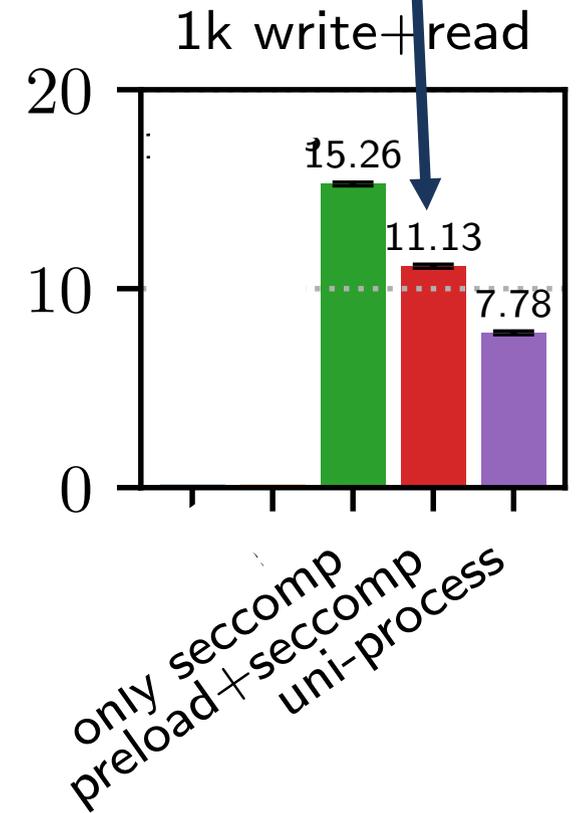
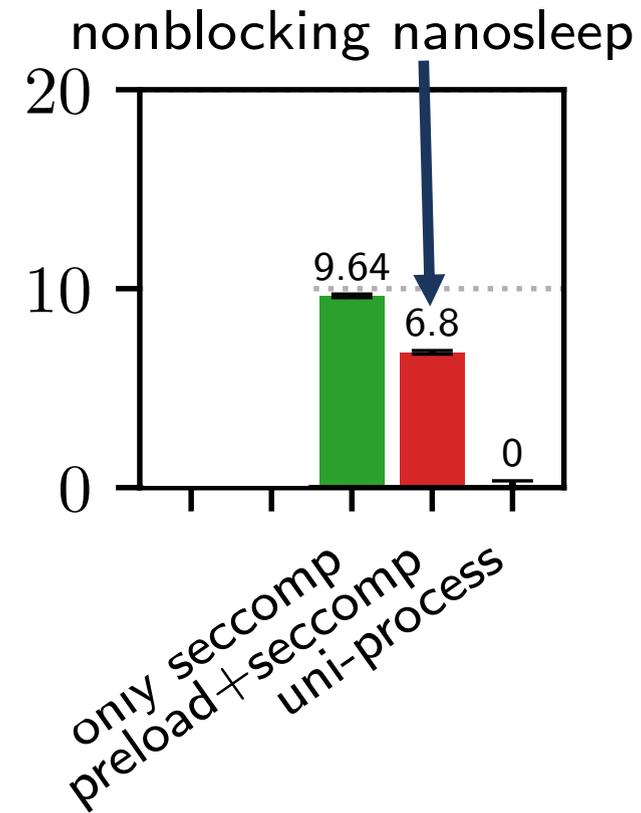
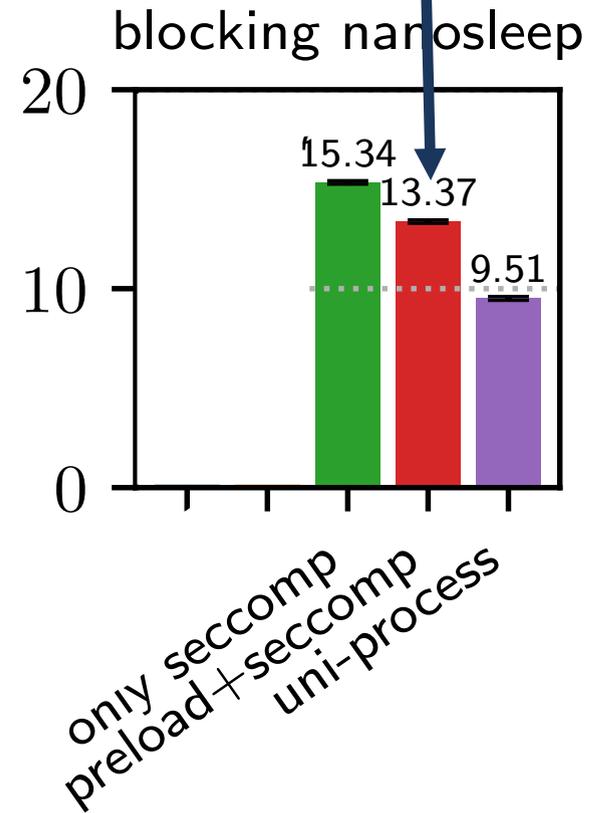


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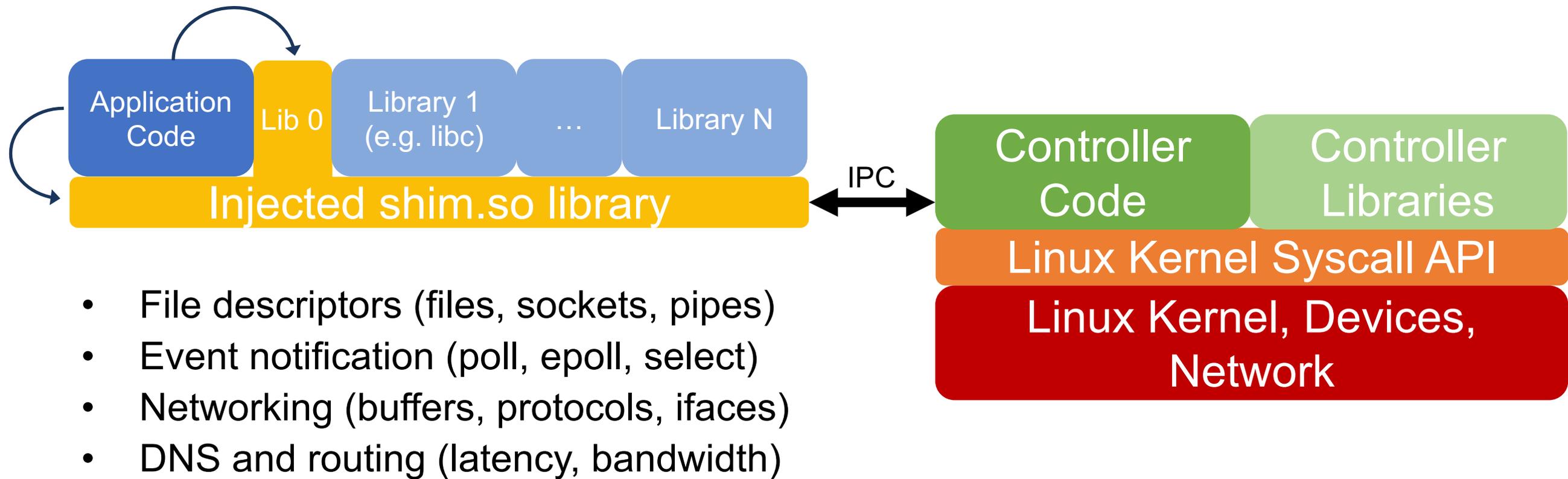
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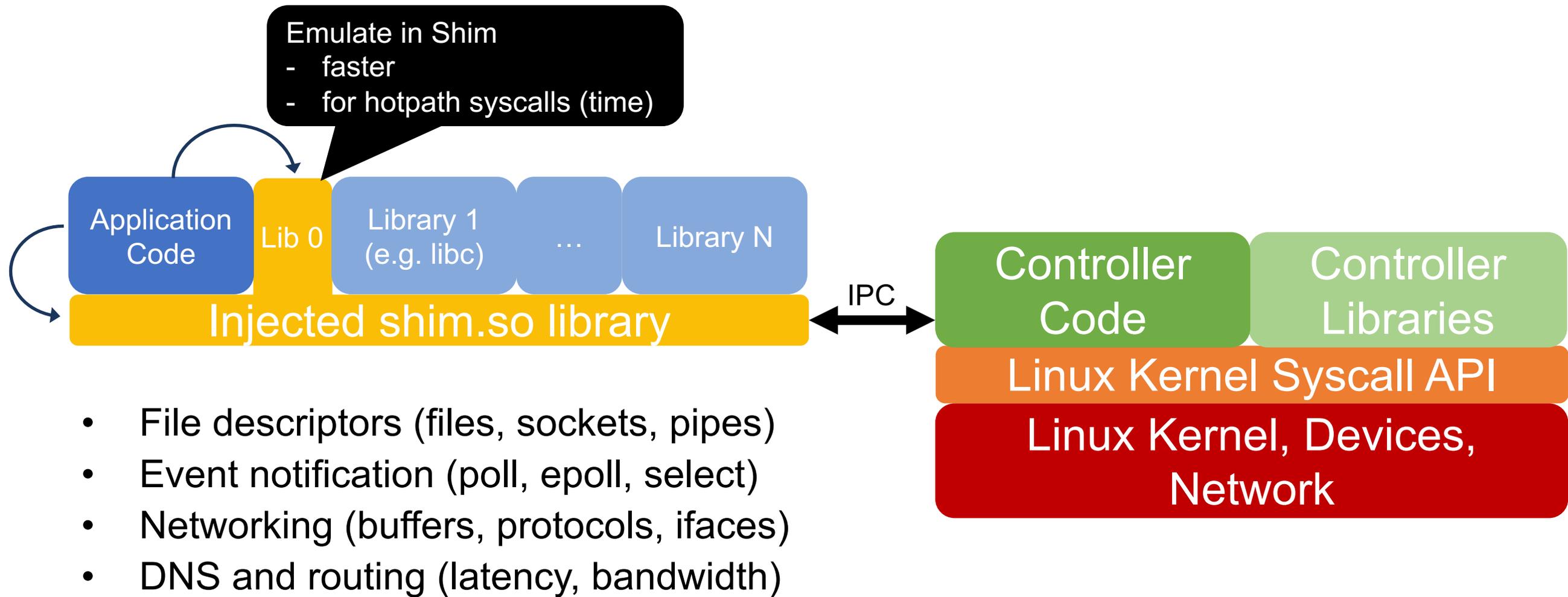


Benchmark Time (μ s)



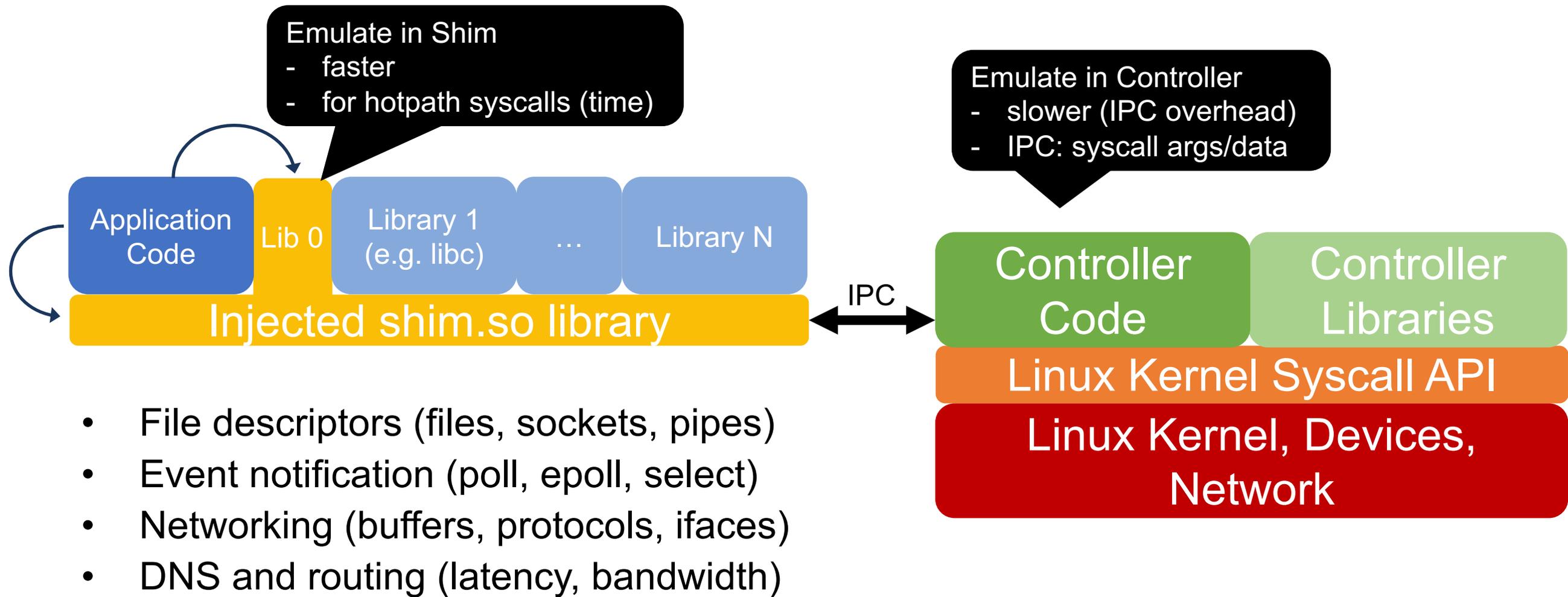
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Syscall Emulation

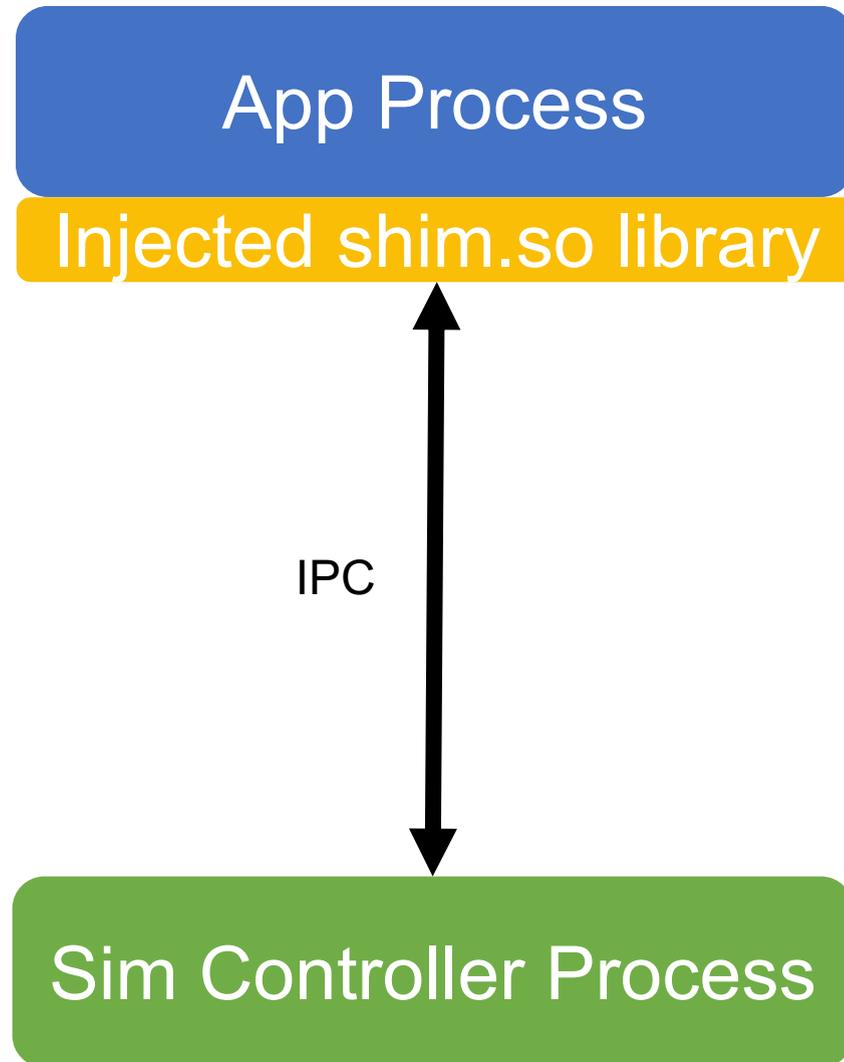
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Inter-Process Communication

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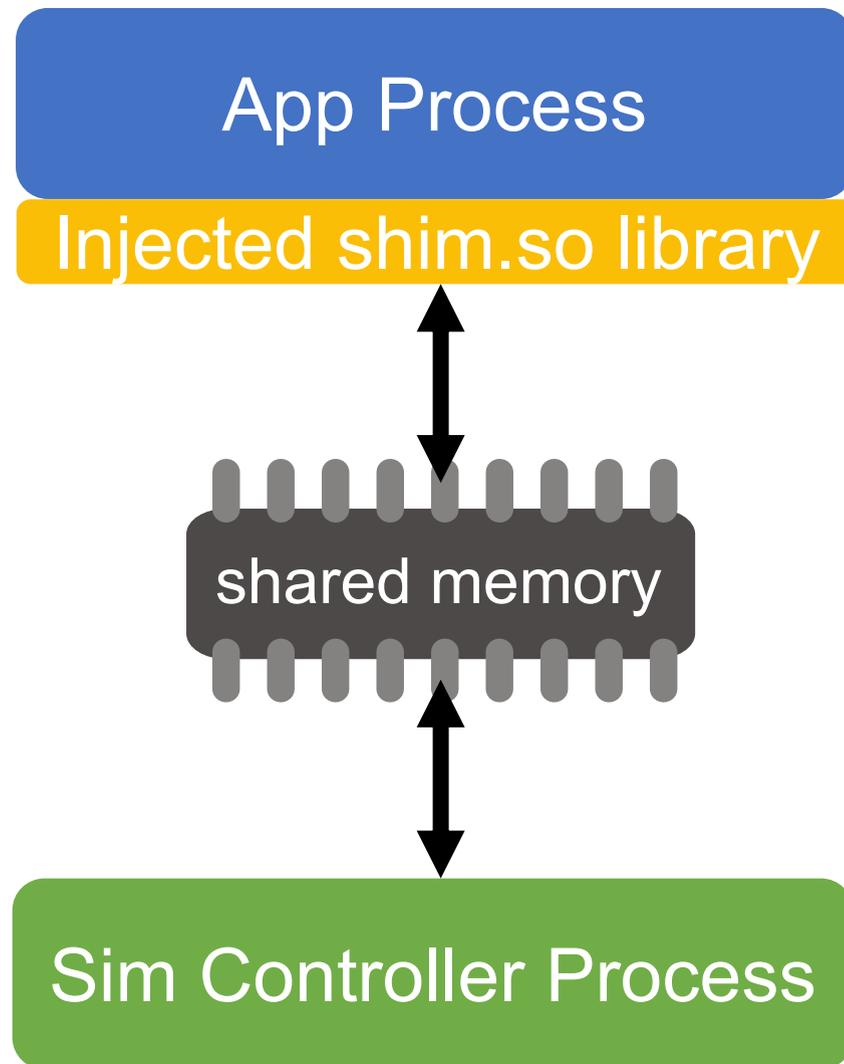
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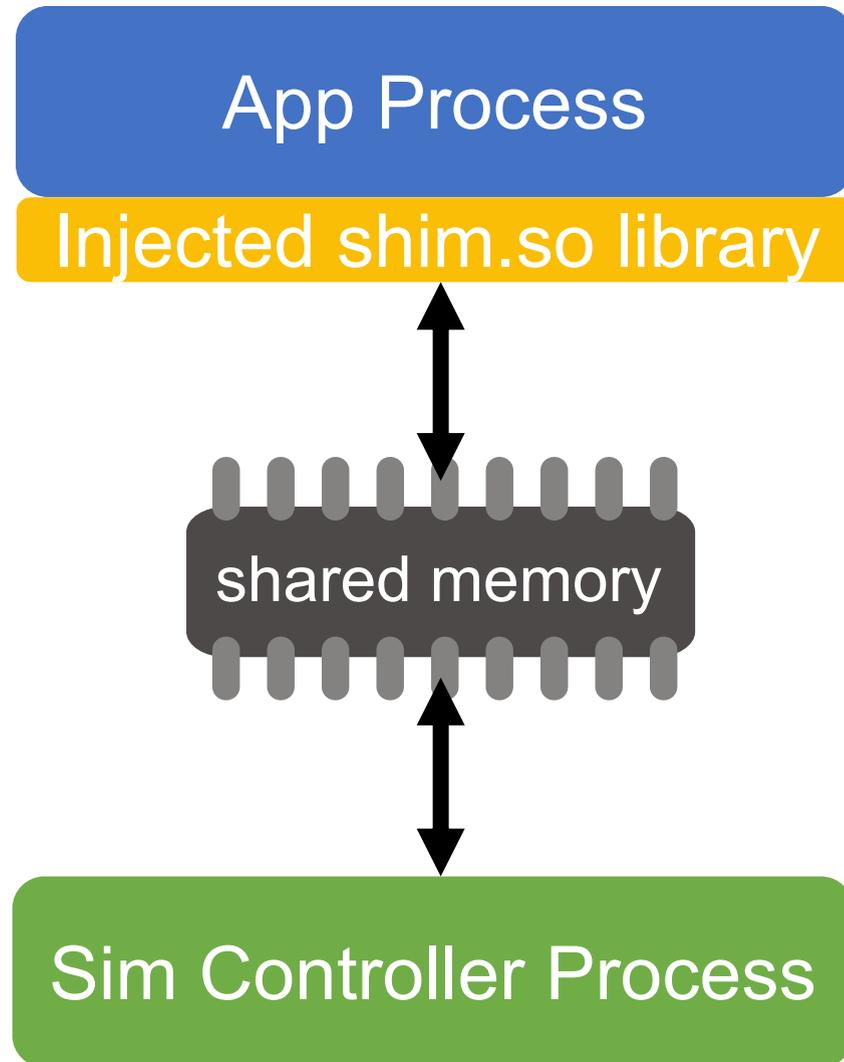
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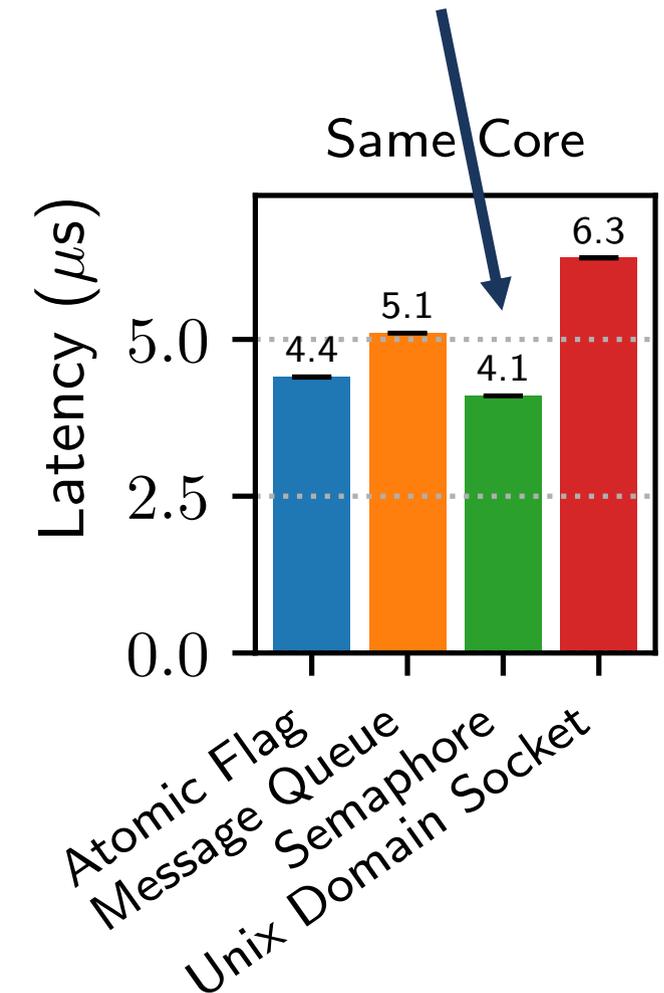
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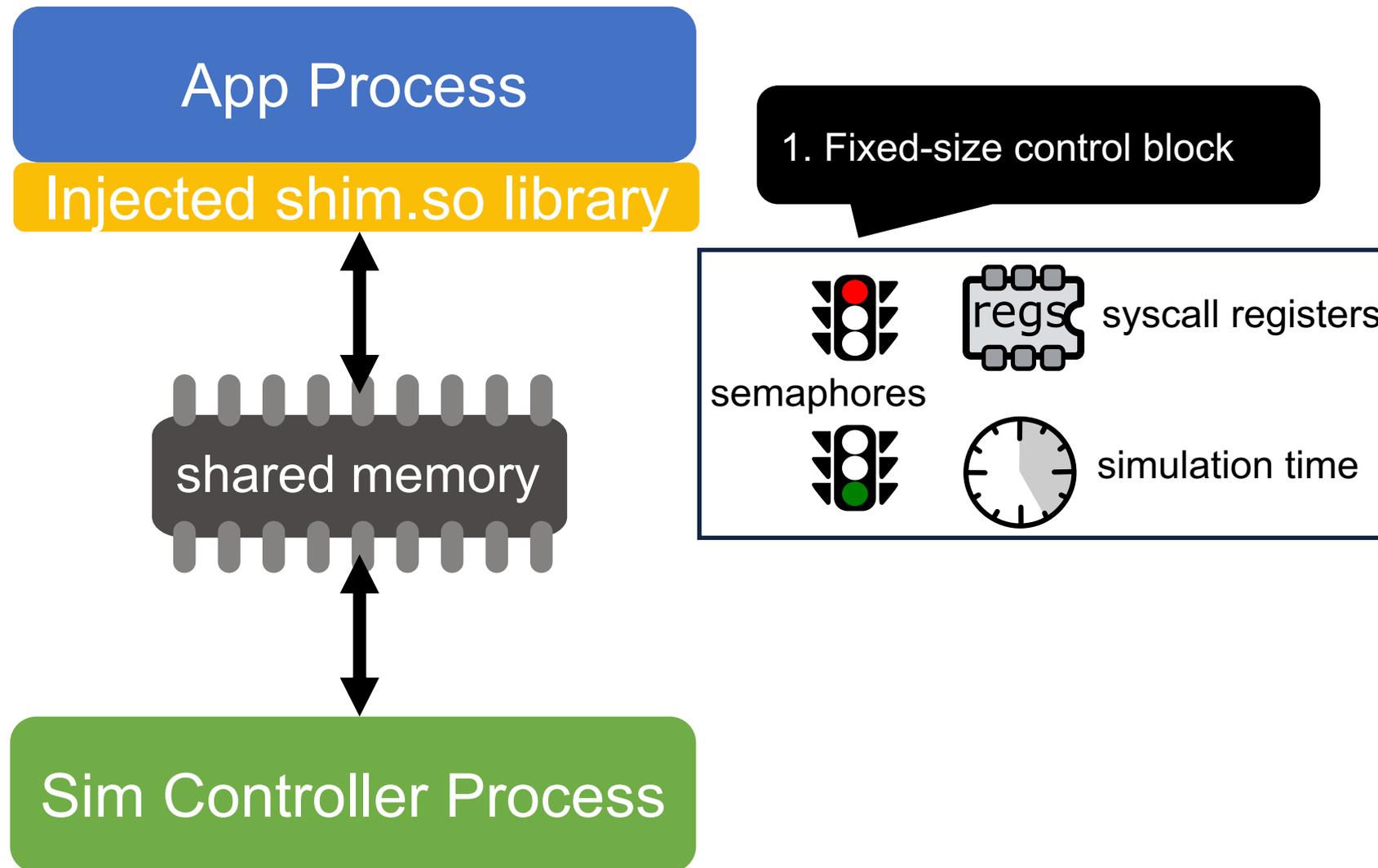
Shared memory +
semaphores is the fastest
IPC method for two
processes running on
the same core



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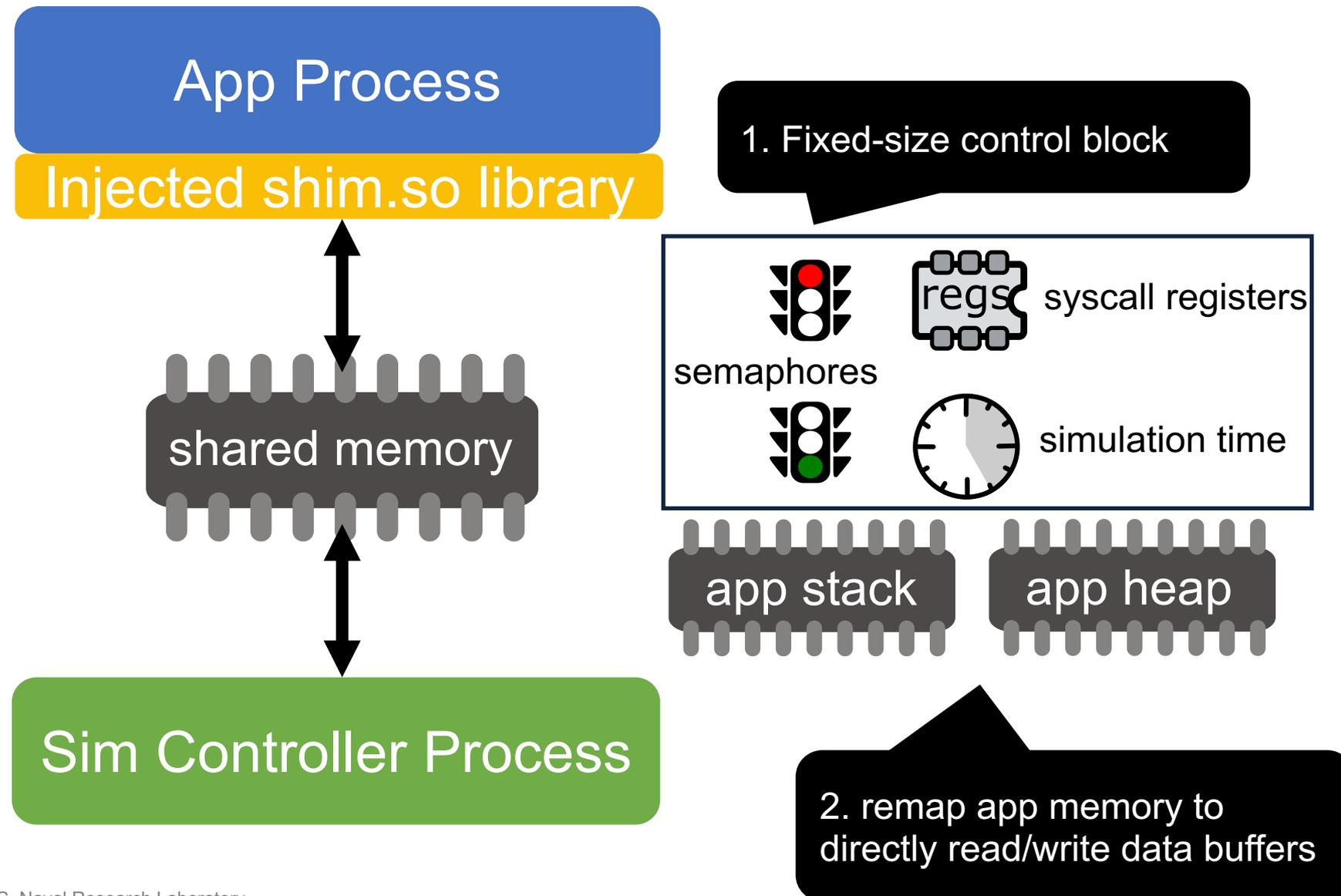
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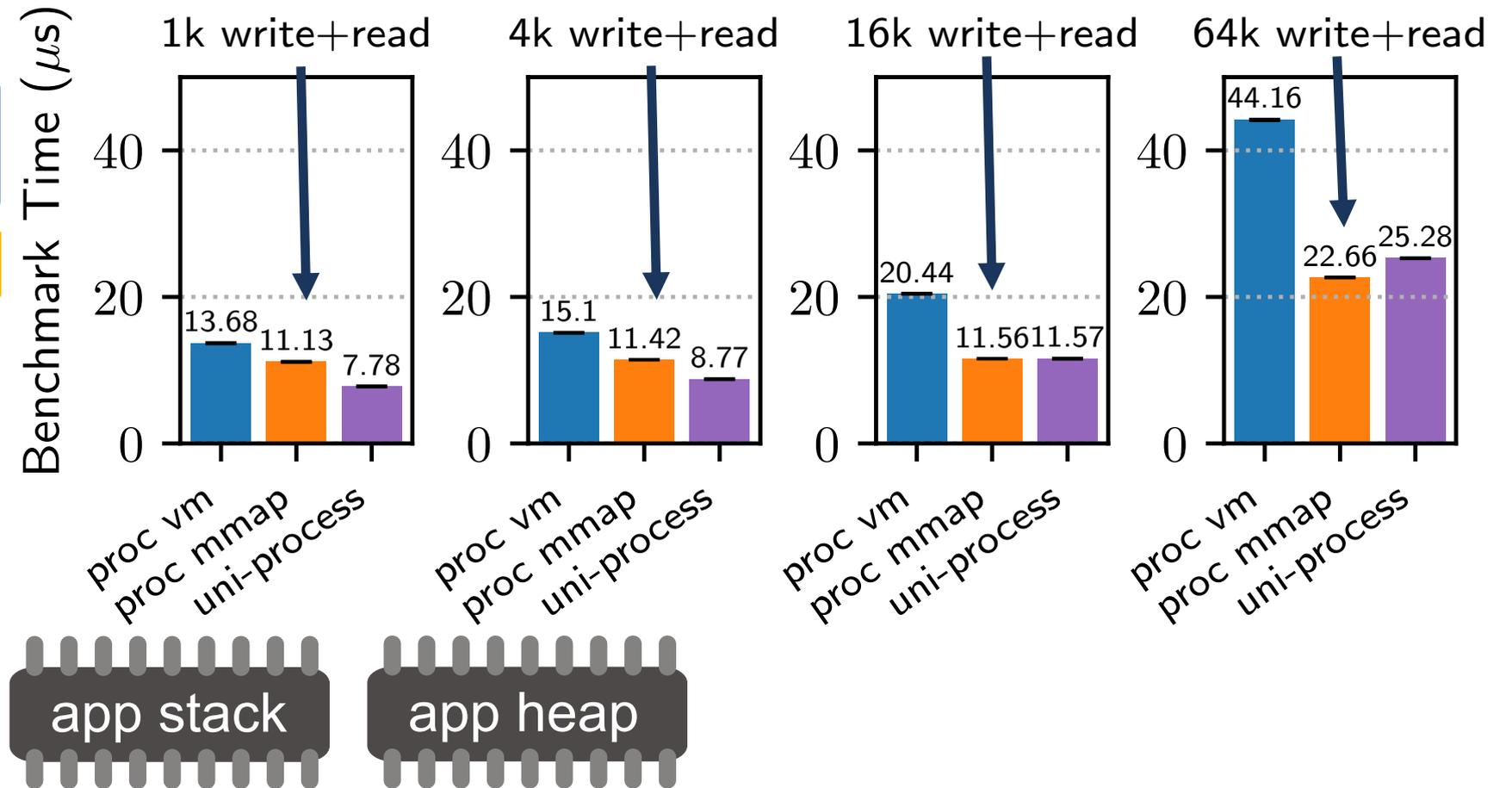
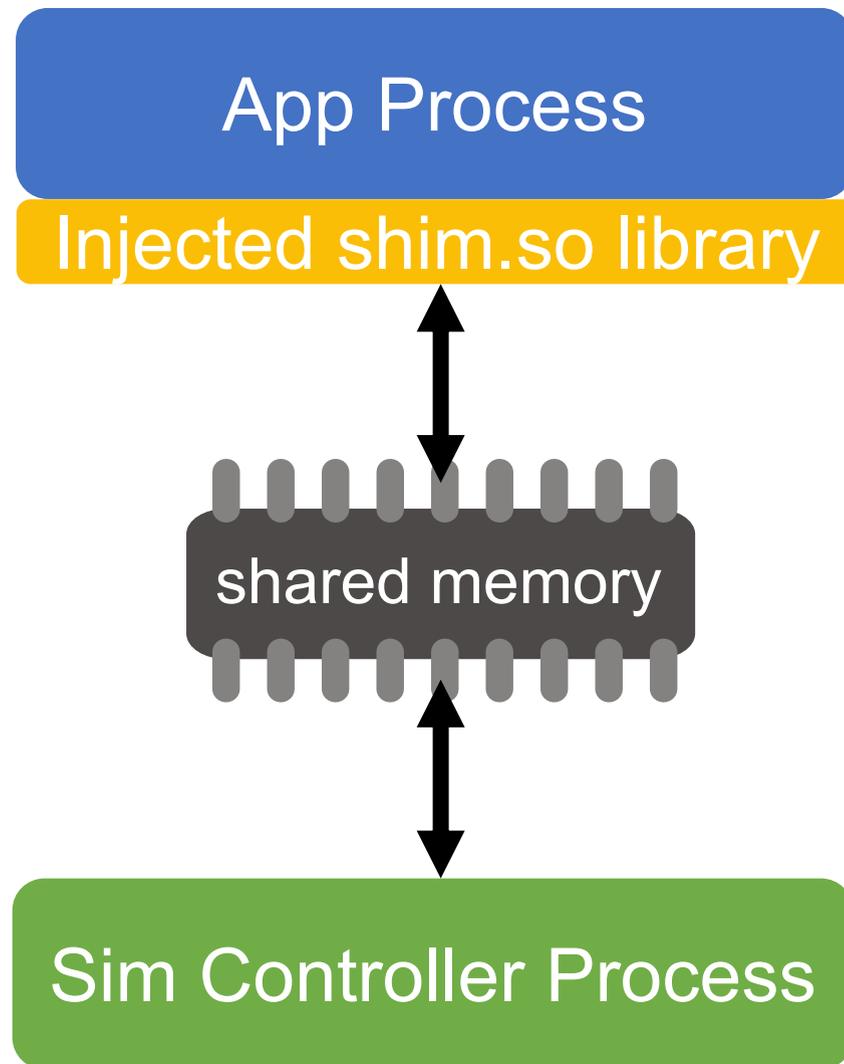
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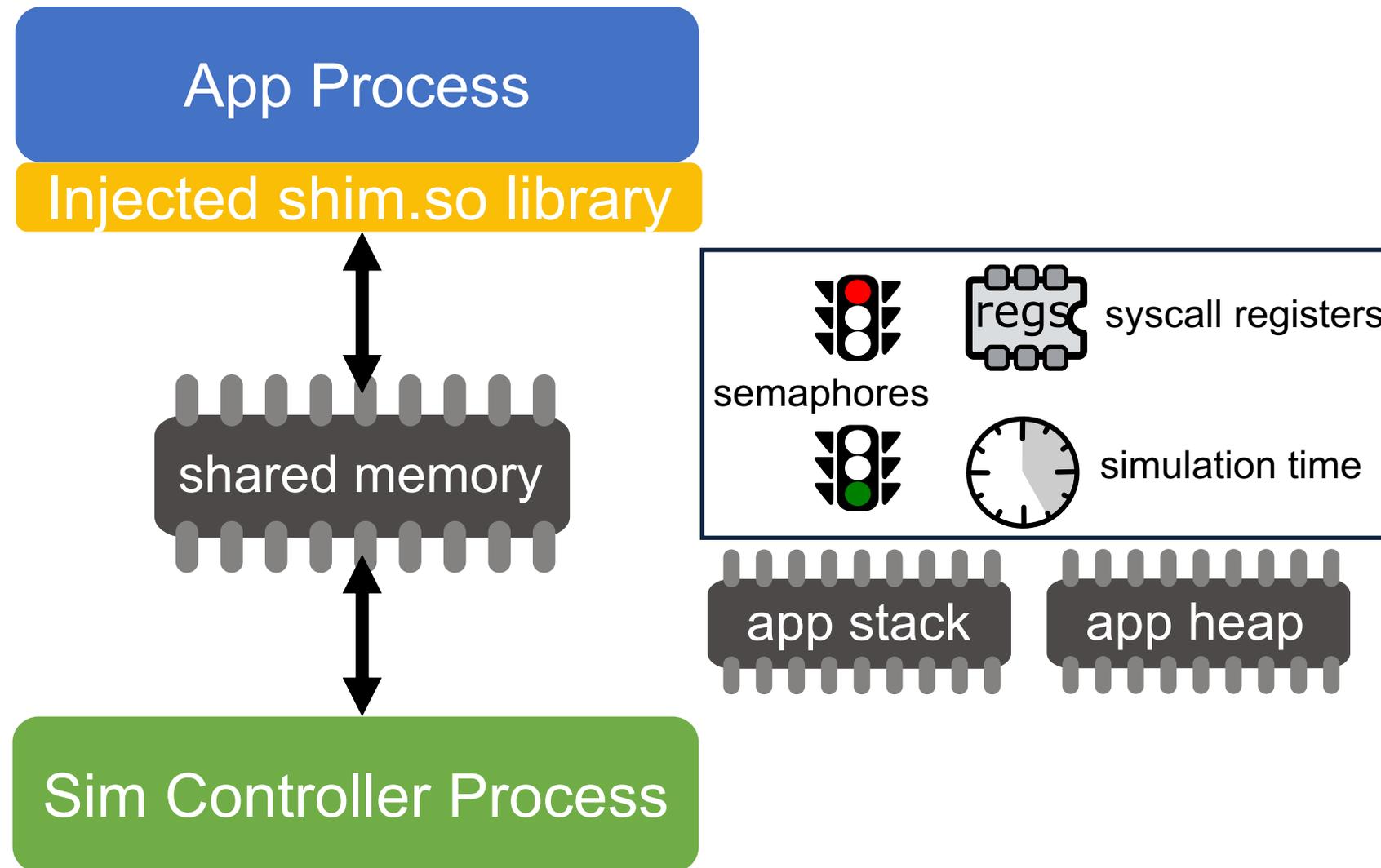
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2. remap app memory to directly read/write data buffers

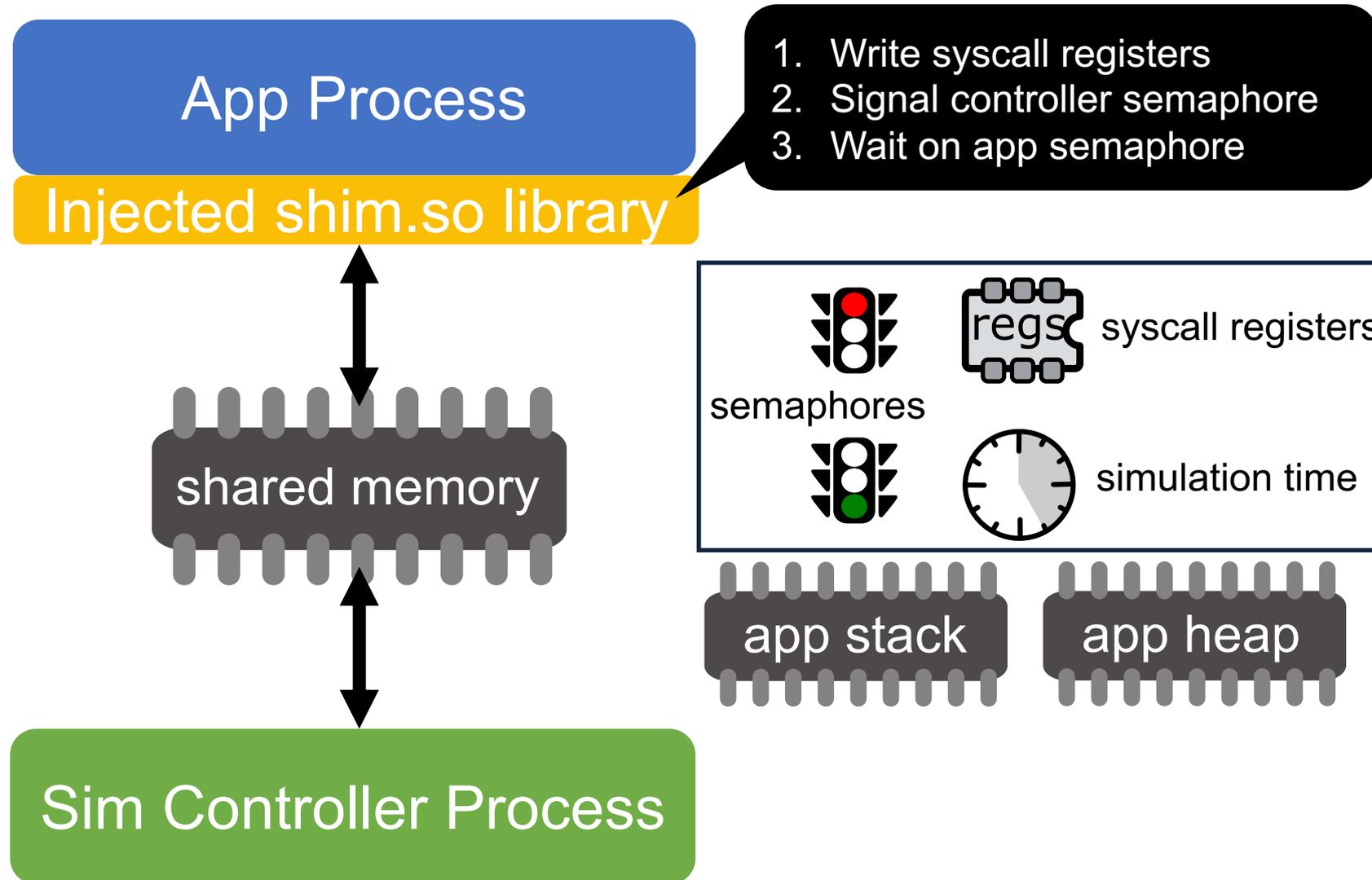
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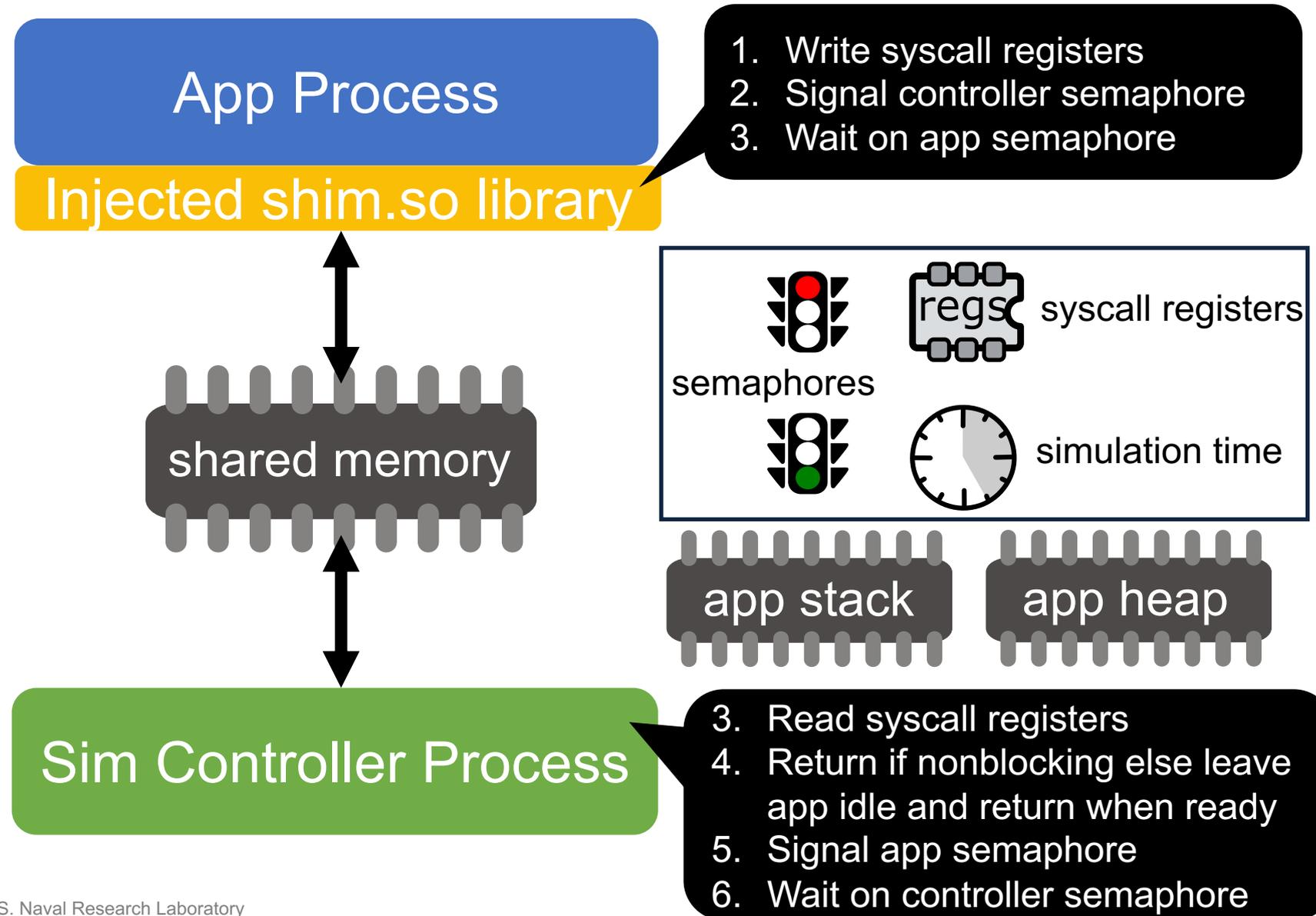
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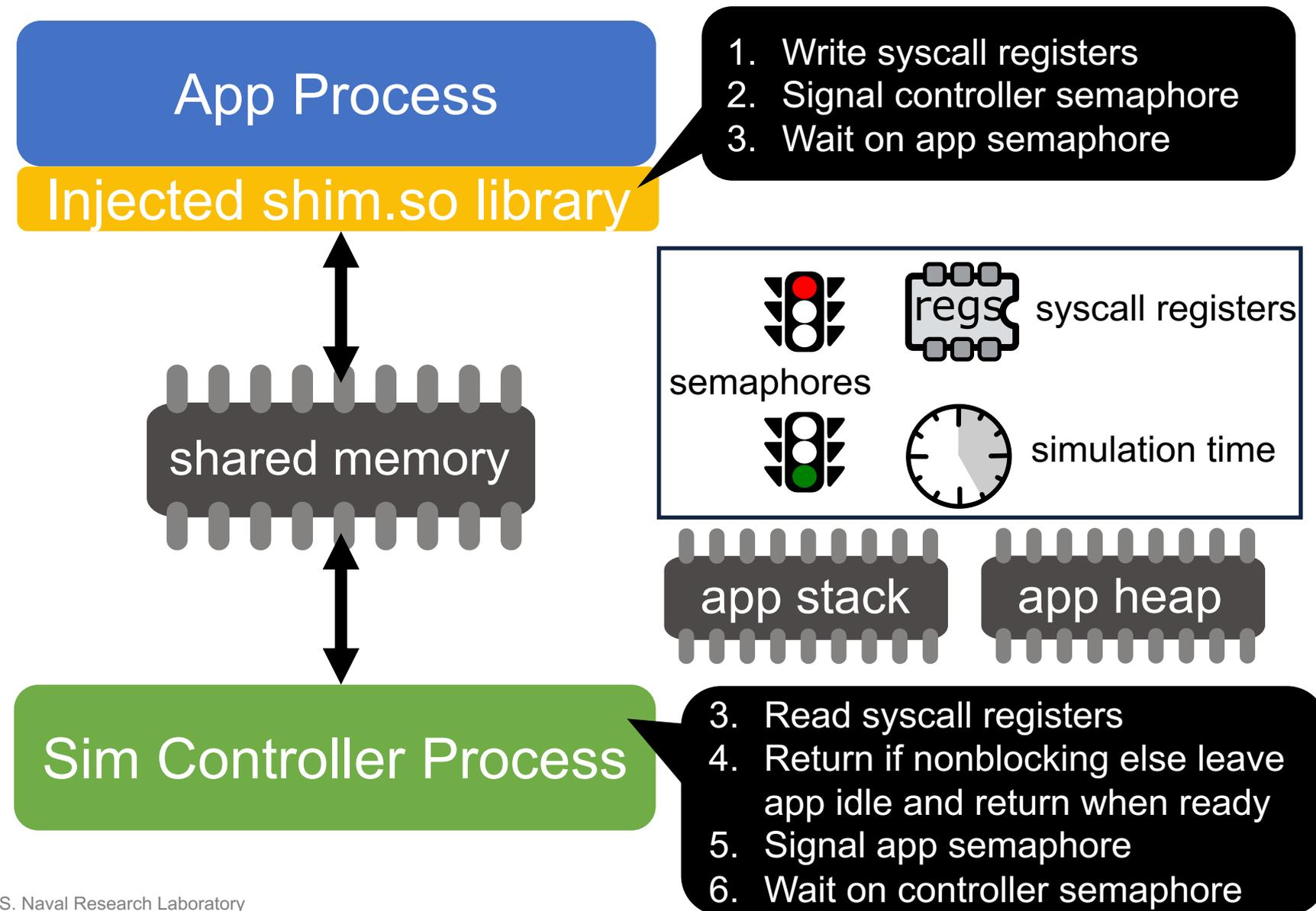
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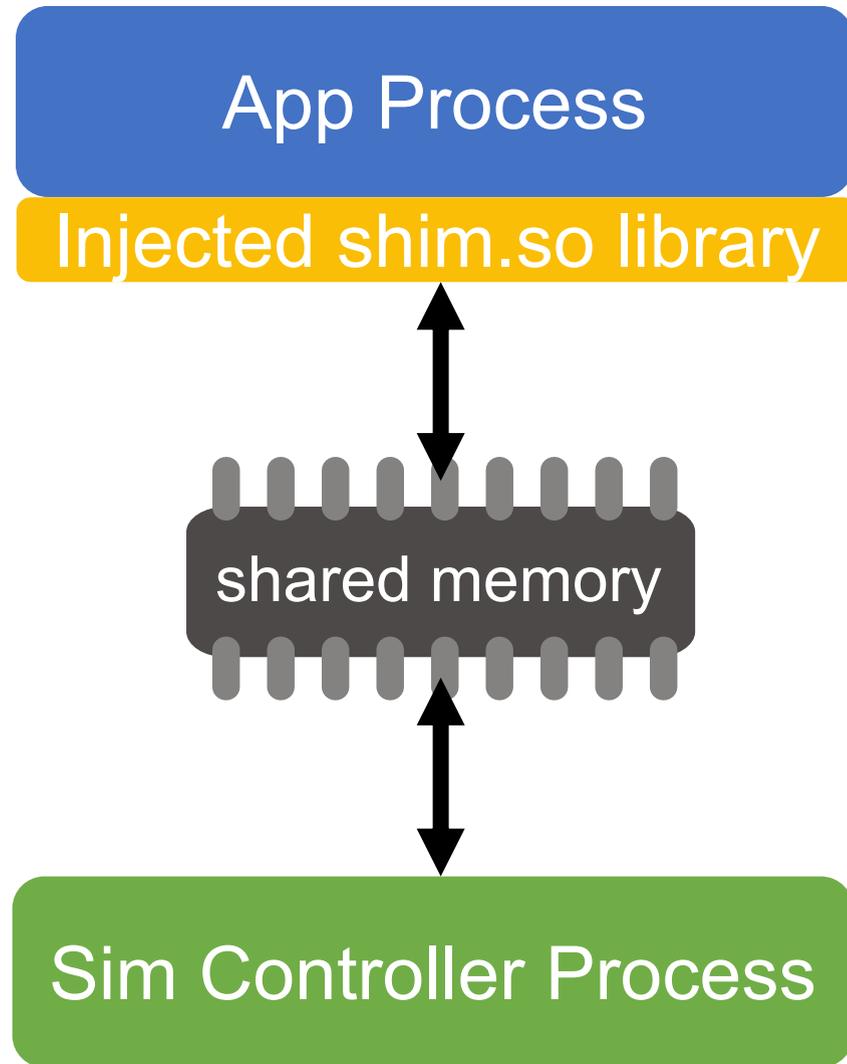
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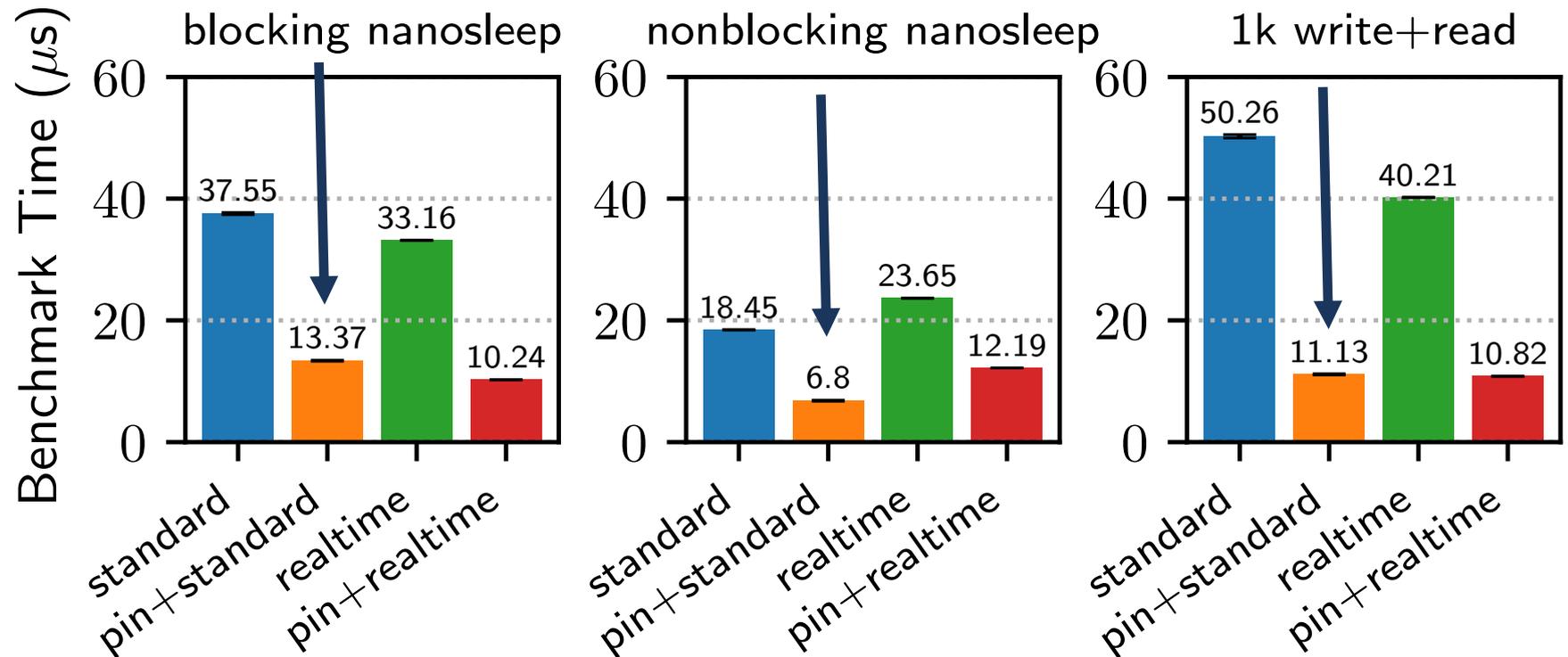
Properties:

- Controller worker thread and its app process run synchronously
- Ensures nonconcurrent access to app stack and heap memory

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Use CPU pinning to pin each worker thread and all of its managed processes to the same core



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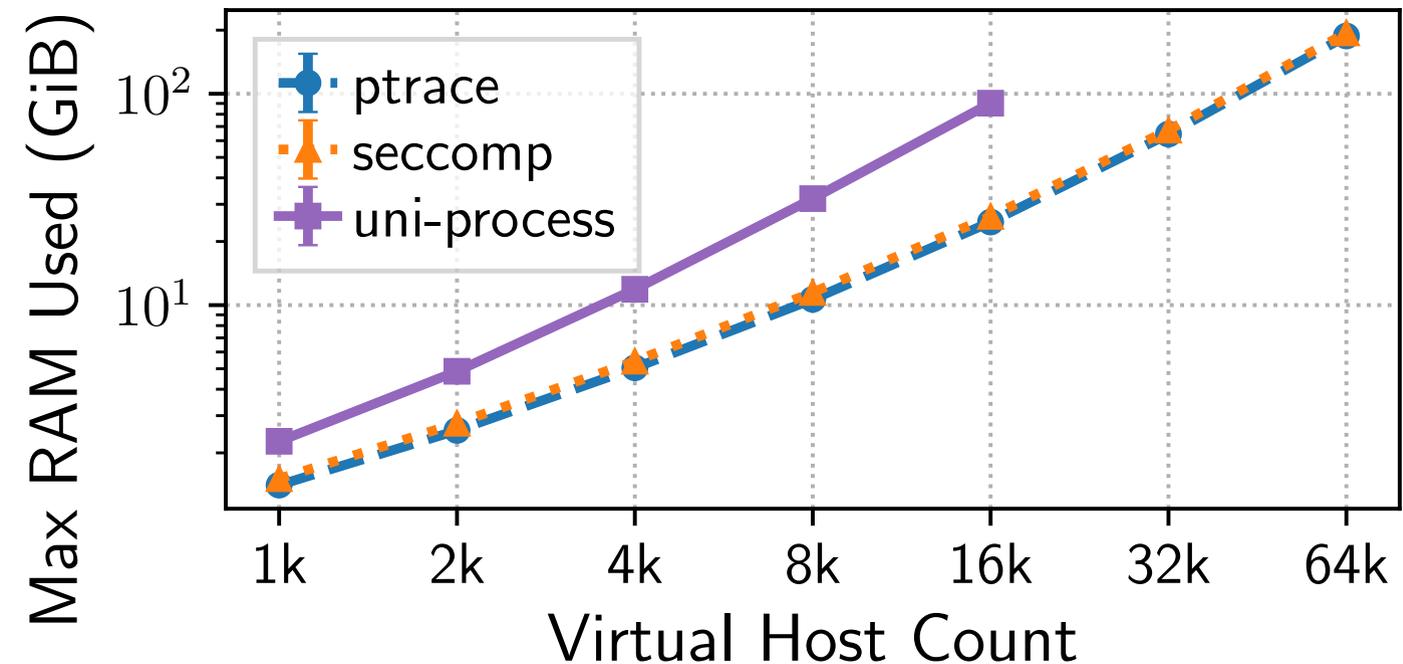
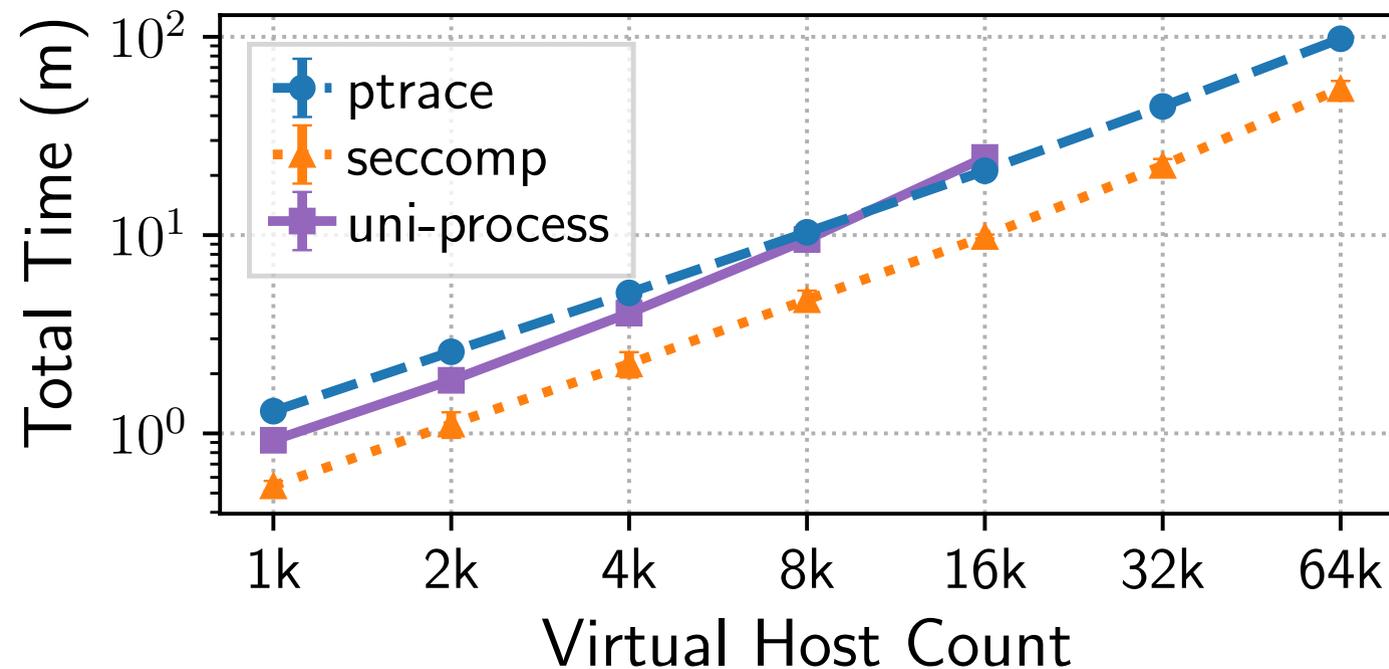
design

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Evaluation: Large P2P Benchmarks

Faster and more scalable than the uni-process, plugin architecture!

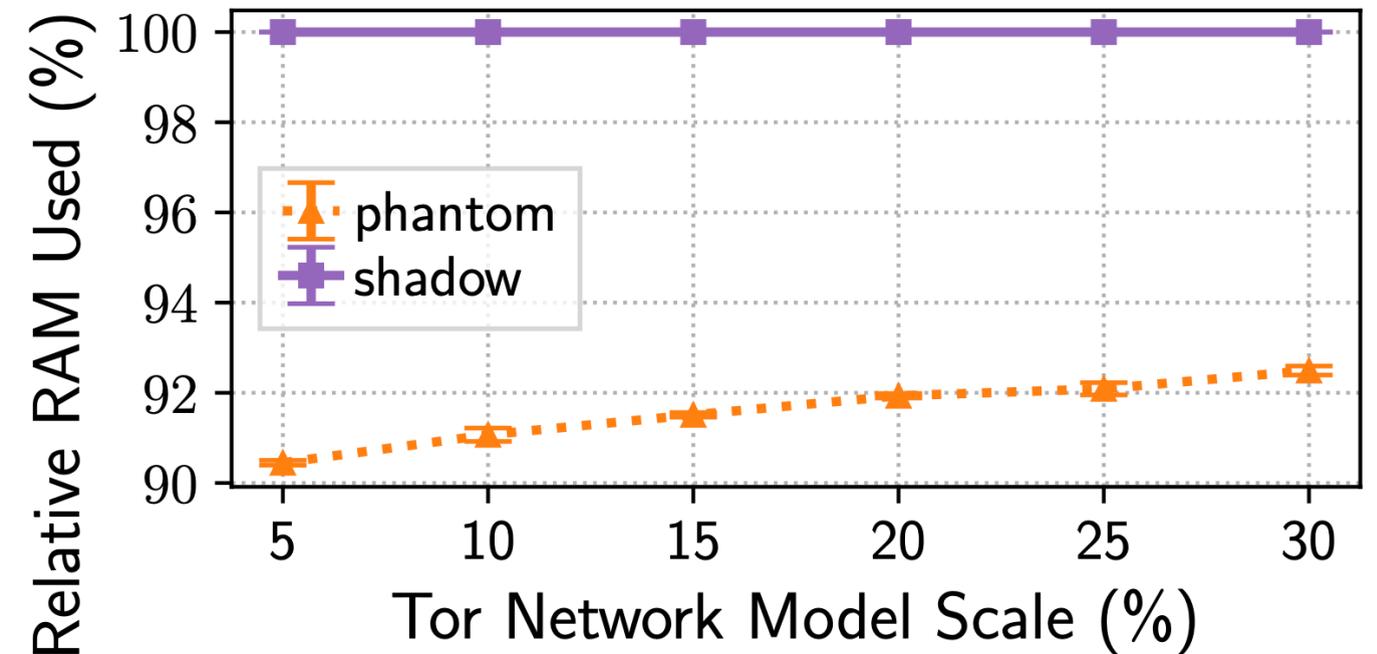
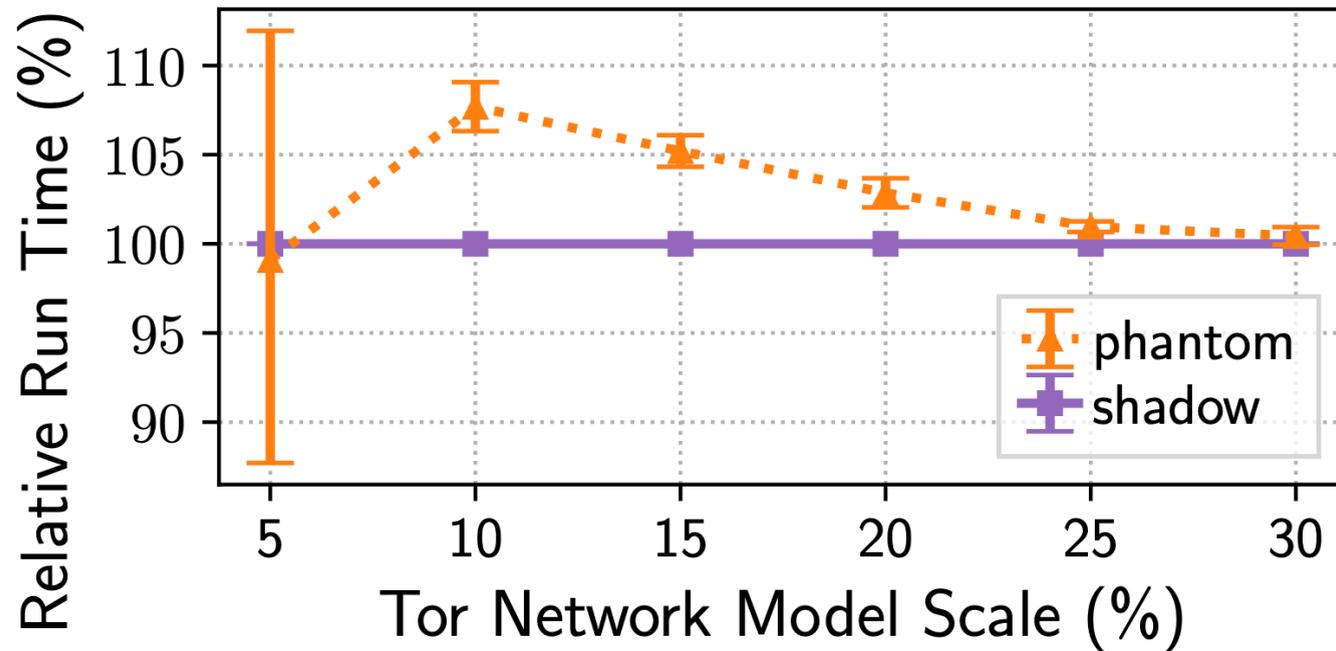
Uses significantly less memory than the uni-process, plugin architecture!



Evaluation: Large Tor Networks

Performance comparable to the state of the art for large Tor networks

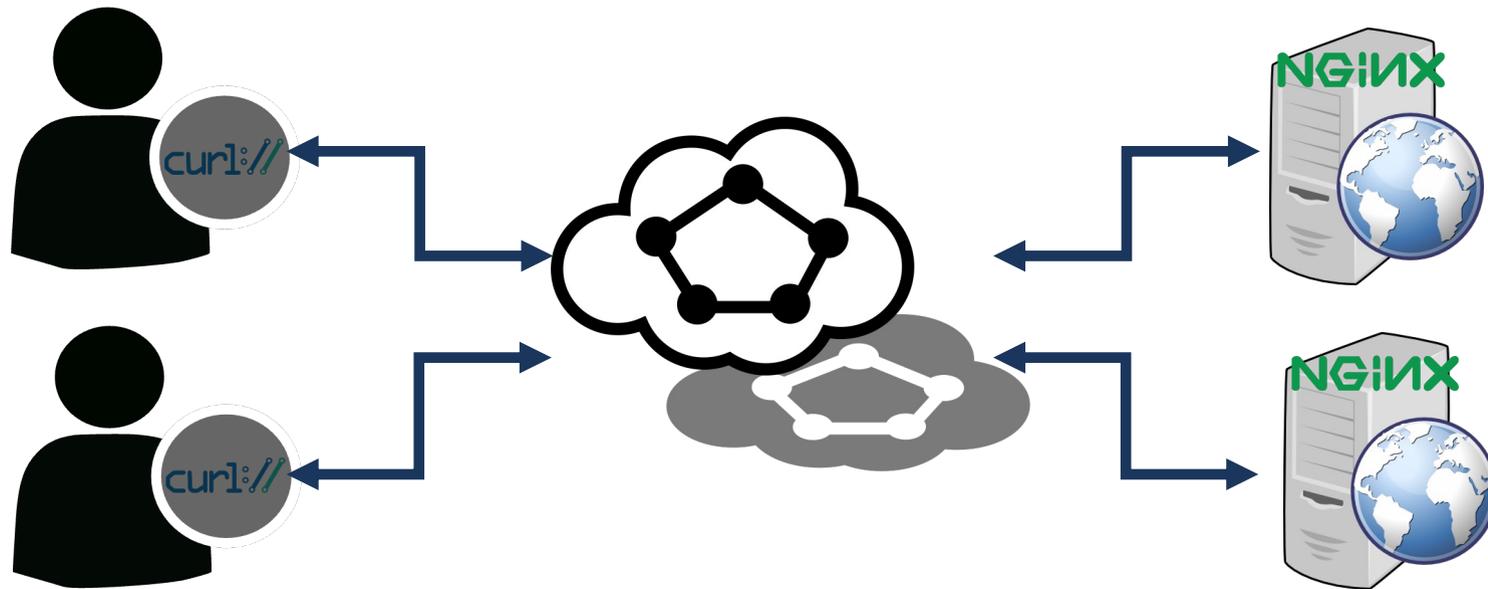
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