A DSL Approach to Reconcile Equivalent Divergent Program Executions

<u>Luís</u>	Daniel	Anastasios	Cristian
<u>Pina</u>	Grumberg	Andronidis	Cadar

{I.pina / daniel.grumberg14 / a.andronidis15 / c.cadar}@imperial.ac.uk Imperial College London London, UK

July 13th, 2017

2017 USENIX Annual Technical Conference (ATC)

What are "Equivalent Divergent Program Executions"?

And why should I care about reconciling them?

Equivalent Divergent Program Executions

>./hello1
Hello world

>./hello2
Hello world

Equivalent Divergent Program Executions

>./hello1
Hello world
>ldd hello1
libc.so.6

>./hello2

Hello world >ldd hello2 libc.so.6

jemalloc.so.6

Equivalent Divergent Program Executions

>./hello1
Hello world
>ldd hello1
libc.so.6

>strace hello1
write(1, "Hello world", 11)

>./hello2

Hello world
>ldd hello2
libc.so.6
jemalloc.so.6
>strace hello2
write(1, "Hello ", 6)
write(1, "world", 5)

Run multiple versions as one

Run multiple versions as one

 Improves reliability malloc(WEIRD_NUMBER) SEGFAULT

jemalloc(WEIRD_NUMBER)
OK

Run multiple versions as one

 Improves reliability malloc(WEIRD_NUMBER) SEGFAULT

jemalloc(WEIRD_NUMBER) OK

> Improves security
>strace hello1
write(1,"Hello world")

>strace hello2
write(1,"Hello, ")
fork()
execve("/bin/sh")
write(1,"world!")

Run multiple versions as one

 Improves reliability malloc(WEIRD_NUMBER)
 SEGFAULT

jemalloc(WEIRD_NUMBER) OK

Improves security >strace hello1 write(1,"Hello world")

>strace hello2
write(1,"Hello, ")
fork()
execve("/bin/sh")
write(1,"world!")

Versions should be diverse but equivalent

Versions should be diverse but equivalent

What about **equivalent** executions that issue **divergent** sequences of system calls?

>strace hello1 >strace
write(1,"Hello world", 11) write()

>strace hello2

- write(1,"Hello ", 6)
- write(1,"world", 5)

Versions should be diverse but equivalent

What about **equivalent** executions that issue **divergent** sequences of system calls?

>strace hello1 >strace hello2
write(1,"Hello world", 11) write(1,"Hello ", 6)
write(1,"world" , 5)

Describe the divergences with a **Domain Specific Language (DSL)**

write(1,"Hello ", 6)
write(1,"world", 5)

write(1,"Hello world", 11)



Varan



Varan



Varan



Varan



Varan



Varan



Varan



Varan



Varan





Varan



DSL Architecture



DSL Architecture



DSL Architecture



DSL Rules

- Default rule: read(_,_,_) as r => r
- Actions
 - ► MATCH
 - ► NOP
 - ► SKIP
 - ► EXECUTE
 - STORE

Further examples

- Hello world
- nothing keyword
- C predicates













Recorded	
read	
"foobar"	
write	
"Hello world"	
?	

Replayed



nothing Keyword

nothing => sched_yield()


nothing => sched_yield()



nothing => sched_yield()

Recorded	Replayed
read	read
"foobar"	"foobar"
write	write
"Hello world"	"Hello "
2	write
:	"world"
	sched_yield

nothing => sched_yield()

Recorded
read
"foobar"
write
"Hello world"
sigaction
sig1
sigaction
sig2
?

Replayed



and multiple left-hand side

// extern int sig1, sig2; sigact(sig,_,_) { \$(sig) == sig1; } as s1, sigact(sig,_,_) { \$(sig) == sig2; } as s2 => s2, s1













Deployment scenarios

- Different configurations
- Different releases
- Different dynamic analyses

Deployment scenarios Different configurations

Recorded Redis minimal config Replayed 1 Redis with persistency (3 rules) Replayed 2 Redis with verbose logs (4 rules) Replayed 3 Redis with persistency and verbose logs (7 rules)

Deployment scenarios

Different releases

ID	Redis Versions <i>Recorded – Replayed</i>	Commits Rules		
1	1.3.8 - 1.3.10	40	0	
2	1.3.10 - 1.3.12	105	0	
3	1.3.12 - 2.0.0	92	1	6
4	2.0.0 - 2.0.5	34	1	0
5	2.0.5 - 2.2.0	730	3	
6	2.2.0 - 2.2.15	110	2	

Deployment scenarios Analyses

Recorded Native Replayed 1 Asan (3 rules) Replayed 2 Msan (1 rule) Replayed 3 Tsan (5 rules) Replayed 4 Valgrind (14 rules)¹

¹Expands to 31 rules through group syntatic-sugar

Deployment scenarios Analyses

Recorded Native Replayed 1 Asan (3 rules) Replayed 2 Msan (1 rule) Replayed 3 Tsan (5 rules) Replayed 4 Valgrind (14 rules)¹

- ▶ git (log, blame, diff, tag)
- openssh (ssh, ssh-keygen)
- ► htop
- ► vim

¹Expands to 31 rules through group syntatic-sugar

Deployment scenarios Analyses

Recorded Native Replayed 1 Asan (3 rules) Replayed 2 Msan (1 rule) Replayed 3 Tsan (5 rules) Replayed 4 Valgrind (14 rules)¹



- ▶ git (log, blame, diff, tag)
- openssh (ssh, ssh-keygen)
- ► htop
- ▶ vim

¹Expands to 31 rules through group syntatic-sugar

Finding these rules must be hard...

Finding these rules must be hard...

lt isn't

Finding Rules

- 1. strace -o native.log native
- 2. strace -o valgrind.log valgrind
- 3. vimdiff native.log valgrind.log

Finding	g Rules	
vimdiff	native.log	
	read(3,, 4096)	gettid(write(1 sigprod read(3,
	lseek(3, -2347, SEEK_CUR)	sigprod gettid(read(10 lseek(3
	read(3,, 4096)	gettid(write(1 sigproc read(3,
	close(3)	sigprod gettid(read(10 close(3

valgrind.log

```
()
1029, "D", 1)
cmask([], ~[...])
. .... 4096)
cmask(~[...], NULL)
()
028, "D", 1)
3, -2347, SEEK_CUR)
()
1029, "E", 1)
cmask([], ~[...])
. .... 4096)
cmask( [...], NULL)
()
028, "E", 1)
3)
```

Finding Rules $_{vimdiff}$		native.log	
	read(3,	,	4096)
	read(3,	,	4096)

valgrind.log

```
gettid()
write(1029, "D", 1)
sigprocmask([], ~[...])
read(3, ..., 4096)
sigprocmask(~[...], NULL)
gettid()
read(1028, "D", 1)
```

```
gettid()
write(1029, "E", 1)
sigprocmask([], ~[...])
read(3, ..., 4096)
sigprocmask( [...], NULL)
gettid()
read(1028, "E", 1)
```



valgrind.log gettid(), write(1029, _, 1), sigprocmask(_, _), r, sigprocmask(_, _), gettid(), read(1028, _, 1) gettid()

```
gettld()
write(1029, "E", 1)
sigprocmask([], ~[...])
read(3, ..., 4096)
sigprocmask( [...], NULL)
gettid()
read(1028, "E", 1)
```

Rule synthesis algorithm

- Rules with the shape: syscall as s => ..., s, ...
- Input: recorded and replayed traces
- Output: set of candidate rules
- Was able to find 16 out of 19 applicable rules
 - Non-determinism and infrequent syscalls impact quality of rules
- ► Details in the paper

Conclusion

- Increases the applicability of multi-version execution
 - For reliability and security
 - State-of-the-art MVE struggles with divergences
- ► Simple expressive language for reconciling system call sequences
 - Recorded and replayed
 - DSL provides the required action to tolerate divergences
- Necessary rules are easy to identify
 - vimdiff of strace logs
- Automatic algorithm to synthethize rules
 - From equivalent strace logs

We're hiring!



- Post-doc position in Software Systems and Program Analysis
- Starting in November 2017, apply until August 2017
- ▶ Up to 17 months, possibly extendable to 24
- Details: https://srg.ic.ac.uk/vacancies

A DSL Approach to Reconcile Equivalent Divergent Program Executions

- Increases the applicability of multi-version execution
 - For reliability and security
 - State-of-the-art MVE struggles with divergences
- ► Simple expressive language for reconciling system call sequences
 - Recorded and replayed
 - DSL provides the required action to tolerate divergences
- Necessary rules are easy to identify
 - vimdiff of strace logs
- Automatic algorithm to synthethize rules
 - From equivalent strace logs

Luís Pina, Daniel Grumberg, Anastasios Andronidis, Cristian Cadar Imperial College London