deTector: a Topology-Aware Monitoring System for Data Center Networks

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Data Center Network Monitoring

- Failures are the norm rather than exception
 - Typical first year for a new cluster (Jeff Dean, Google)
 - 8 network maintenances
 - 15 router reloads/failures
 - 26 rack failures/moves
 - Dozens of blips of DNS
 - 1000 individual machine failures
- SLA violation (99.999%)
 - Packet losses and latency spikes
 - Difficult to troubleshoot (up to days to fix the issues)

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Challenges

- Clean failures
 - Easy to detect, e.g., server down.
- Gray failures
 - Not reported by the device (SNMP/CLI)
- Low-rate losses
 - Covered up by ECMP
- Transient failures
 - Difficult to play back and pinpoint

Challenges

- Clean failures
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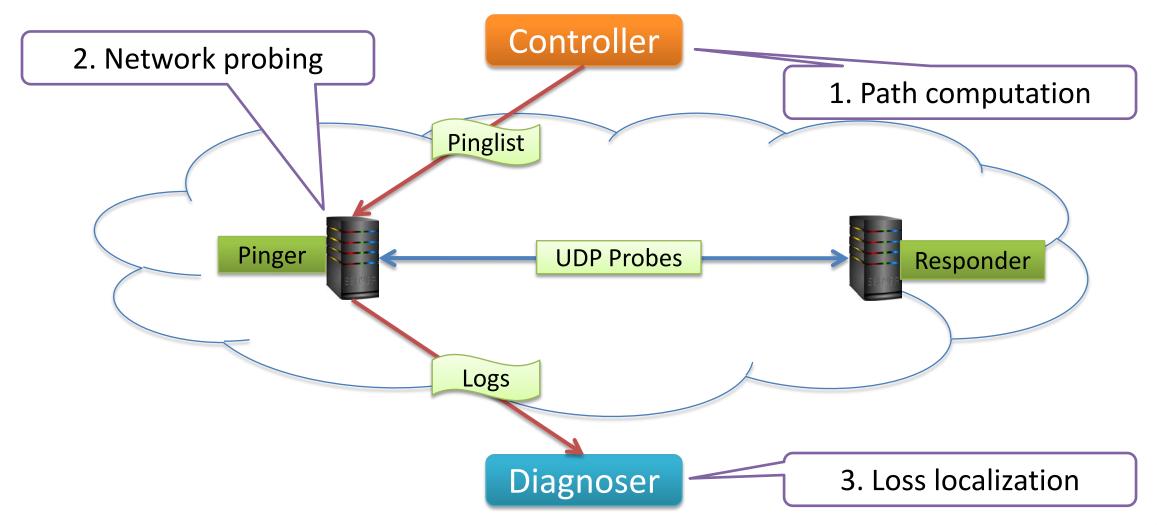
Existing Solutions

- Existing systems
 - Passive: CLI/SNMP
 - Active: Pingmesh, NetNORAD
- Limitations
 - Fail to detect at least one type of losses
 - High overhead
 - Can not pinpoint failures without other tools (e.g., tracert)

Existing Solutions

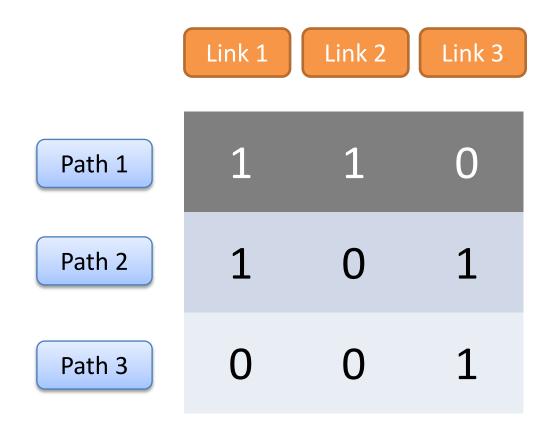
- Existing systems
 - Passive: CLI/SNMP
 - Acti Can we design a better network
- Limitation monitoring system by exploiting
 - Fail o detect at network topology?
 - High overhead
 - Can not pinpoint failures without other tools (e.g., tracert)

deTector in One Slide



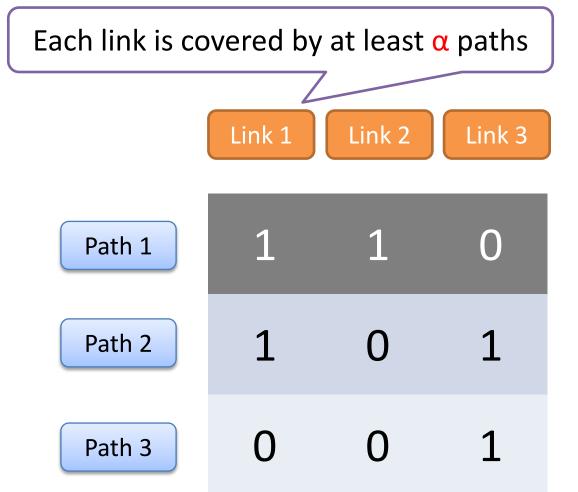
Phase I: Path Computation

Path Selection Problem



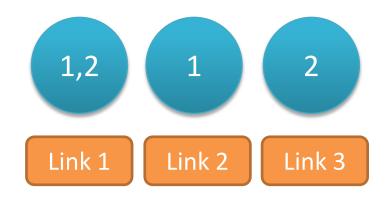
- Given a routing matrix, select probing paths to send probes:
 - path number minimizing
 - α-coverage
 - $-\beta$ -identifiability

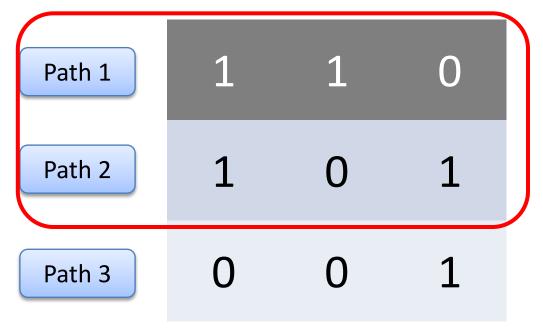
α-coverage



- Ensure even and enough path coverage of each link
- w[link]: track the number of paths through it.
- If w[link] > α, then the link
 has enough paths through it.

β-identifiability

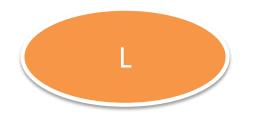




 Any β failed links can be identified correctly

Probe matrix: path1 + path2
 1-identifiability but not 2 identifiability

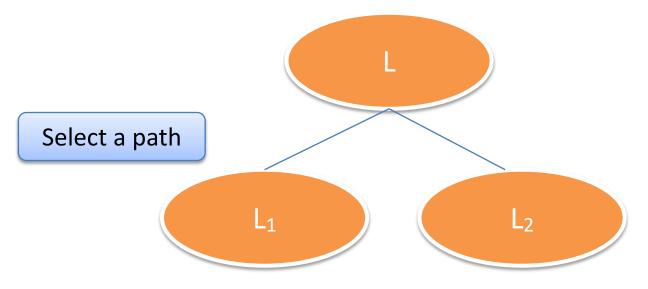
Algorithm for 1-identifiability



 Select the minimum number of paths so that each link has a different set of probe paths.

 Greedily select the path which splits the largest number of link sets in each iteration.

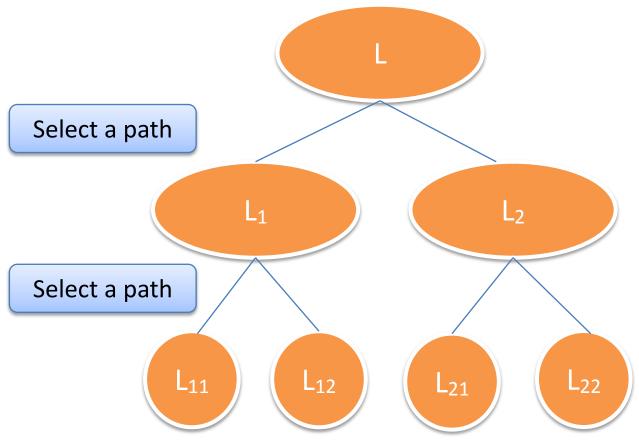
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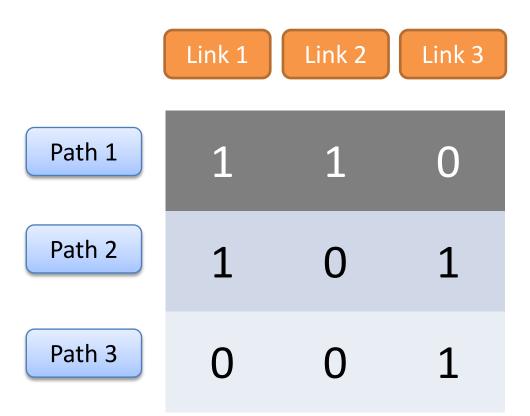
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1-identifiability => β-identifiability

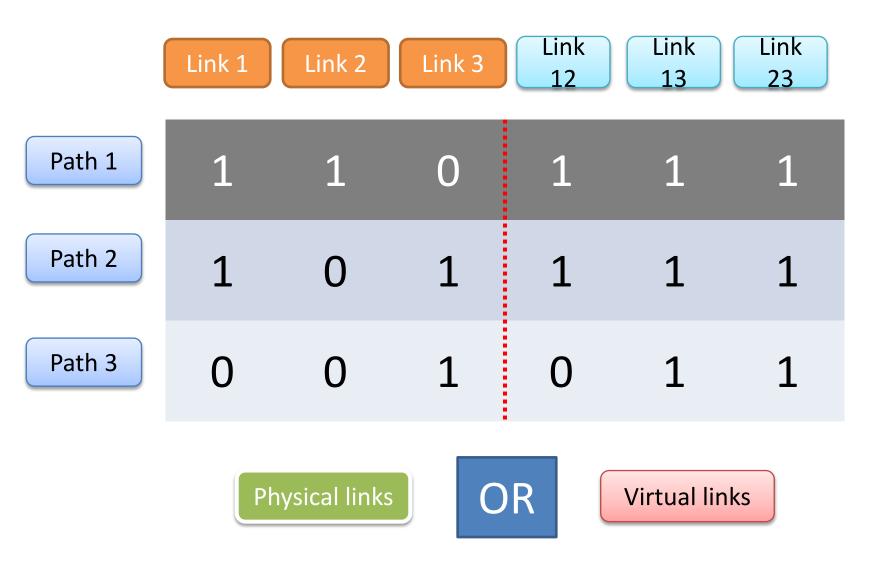


Physical links

 Extend routing matrix with virtual links.

 If a virtual link is down, we say its corresponding physical links have failed.

1-identifiability => β-identifiability



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Probe Matrix Construction (PMC) Algorithm

- Extend the routing matrix with virtual links • Define a score for each path Quantify coverage $score(path) = \sum_{link \in path} w[link] - \# \text{ of link sets on } path$ Quantify identifiability
- Select a path with minimal score in each iteration
- Stop when achieving α -coverage and β -identifiability

PMC Algorithm

- Achieve 63% approximation ratio.
- Time complexity O(n²) where n is the number of paths.
- A Fattree(64) DCN has more than 2³² paths, running time > 24 hours

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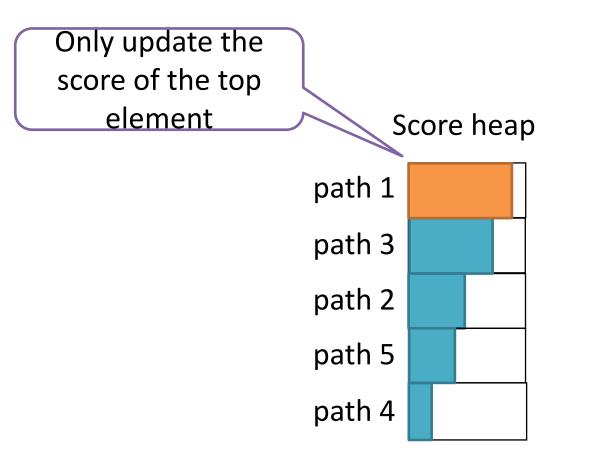
PMC Algorithm

- Achieve 63% approximation ratio.
- Time complexity O(n²) where n is the number of paths. Optimizations for speedup
- A Fattree(64) DCN has more than 2³² paths, running time > 24 hours

Optimization I: Routing Matrix Decomposition

1	1	0	0	0		1	1	0
1	0	1	0	0	Share no links	1	0	1
0	1	1	0	0	and paths	0	1	1
0	0	0	1	0			1 (0
0	0	0	1	1			1 :	1

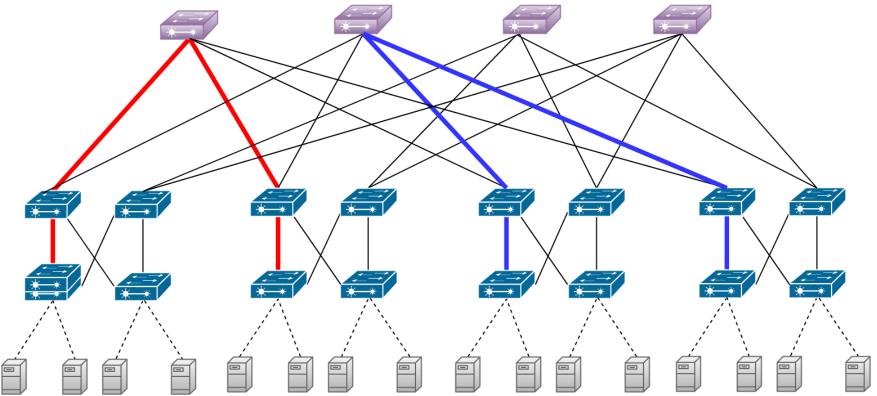
Optimization II: Lazy Update



• Defer the score update of a path as much as possible until we have to.

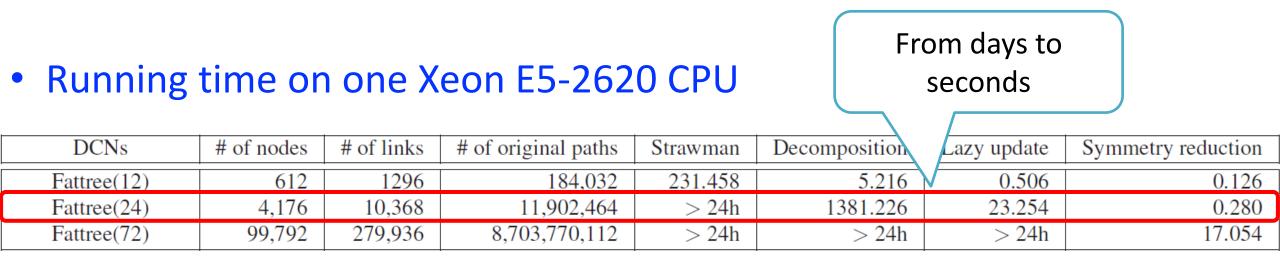
 Correctness guaranteed by the submodularity of the objective function.

Optimization III: Symmetry Reduction

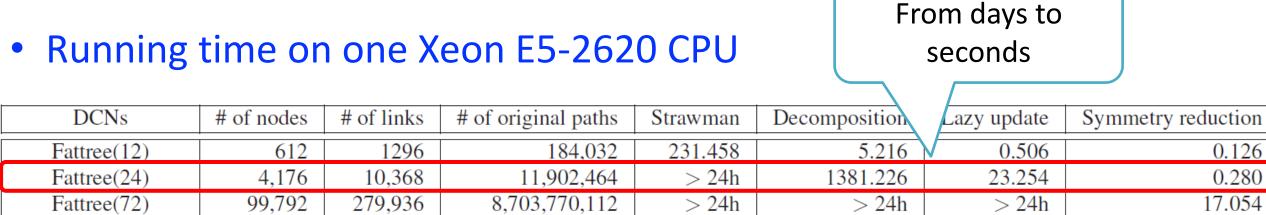


Most DCN topologies are symmetric!

PMC Algorithm Results of Fattree



PMC Algorithm Results of Fattree





DCNs	# of original paths	# of sele	(α,β)	
DCINS		(1, 0)	(1, 1)	(3, 2)
Fattree(32)	66,977,792	4,096	7,680	12,288
Fattree(64)	4,292,870,144	32768	61,440	98,304

The optimal needs

at least 52428

paths

Phase II: Network Probing

Network Probing

- UDP probes: varying packet length, DSCP, source port
- Source routing: IP-in-IP encapsulation and decapsulation
- Responders: simply echo probes back

Phase III: Loss Localization

Loss Localization Problem



 Given the probe matrix and loss measurements, select the least number of links to explain the observation.

NP-hard

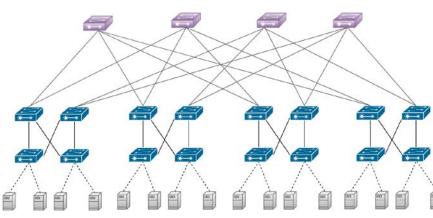
Packet Loss Localization (PLL) Algorithm

- In each iteration we select a link that can explain the largest number of probe losses until all are explained
 - If a link lies in the packet path, then the link can explain the loss.

- Two simple improvements
 - Matrix decomposition to speedup computation
 - Use threshold to filter false positives
 - The ratio of # of lossy paths over # of all probe paths through the link

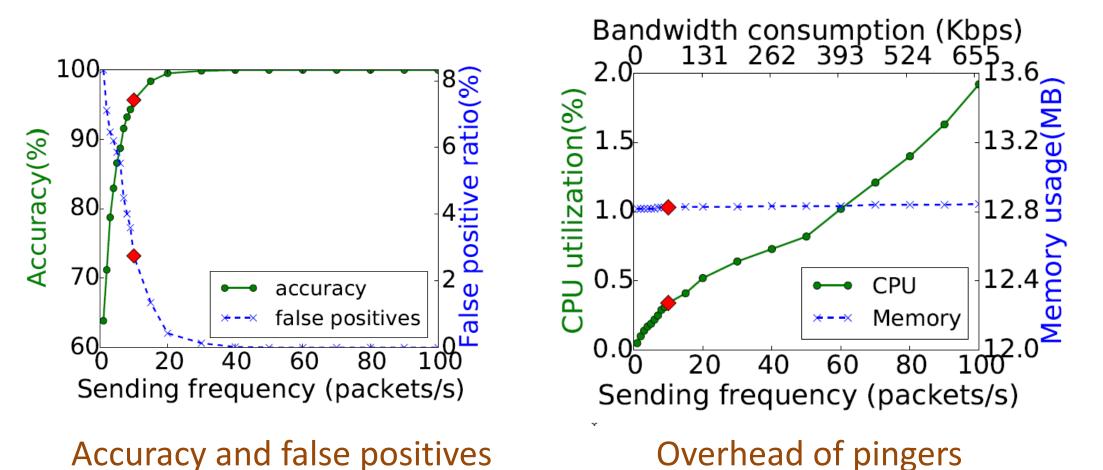
- A 4-ary Fattree testbed with 20 SDN switches
- Install OpenFlow rules to emulate losses caused by various failures
 - Full packet loss: link down etc.
 - Deterministic partial loss: packet blackhole etc.
 - Random partial loss: bit flips etc.
- Performance metric
 - Accuracy
 - False positive ratio

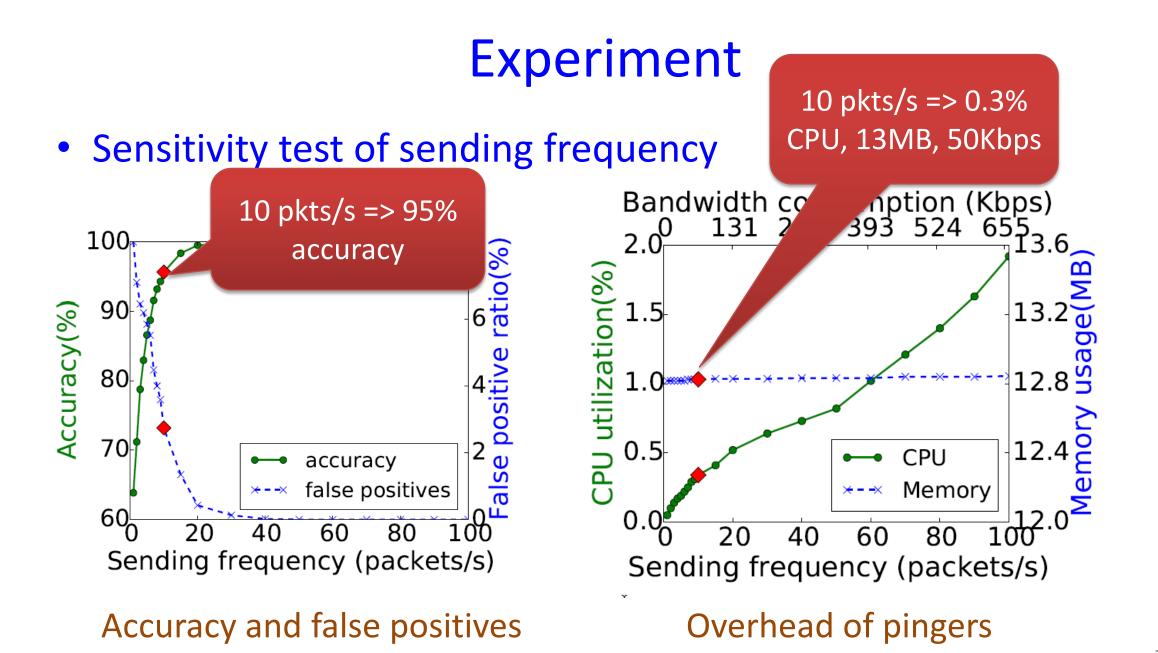
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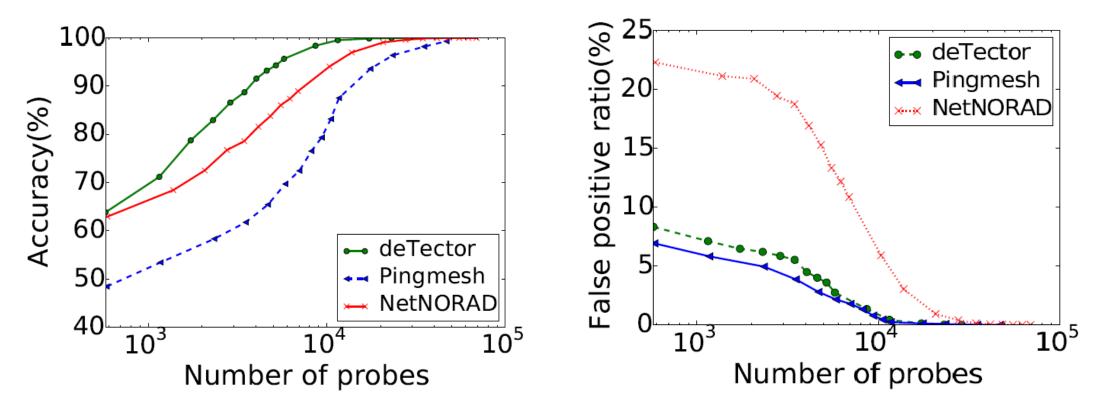


Sensitivity test of sending frequency

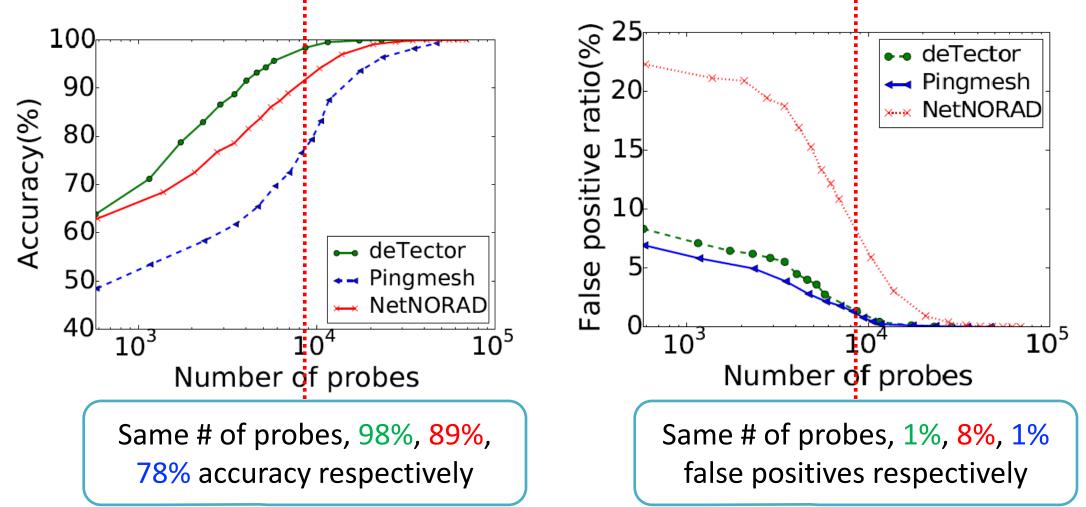




• Accuracy and false positives of three monitoring systems



Accuracy and false positives of three monitoring systems



Simulation

• Accuracy in a 18-radix Fattree, with probe matrices of different levels of α -coverage and β -identifiability

(lpha,eta)	# of paths	Accuracy (%) with $\#$ of failed links						
		1	5	10	20	50		
(1, 0)	729	30.56	30.87	30.30	30.26	29.19		
(2,0)	1485	58.43	57.43	57.08	56.81	57.11		
(3, 0)	2187	68.22	70.61	69.89	70.40	70.14		
(1, 1)	1269	94.74	93.37	94.21	93.43	90.29		
(1, 2)	1512	99.26	99.06	99.02	98.77	95.92		
(1, 3)	2349	99.63	99.63	99.67	99.62	98.07		

Simulation

Identifiability is more effective than coverage for failure localization.

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Summary

- The core of deTector is a carefully designed probe matrix, enabling fast and accurate loss detection and localization.
- deTector is practically deployable.
- Discussions
 - Packet entropy: limited destination IP addresses
 - Loss diagnosis: do not know why packets are dropped
 - Beyond deTector: apply probe matrix to optimize in-band techniques

Thanks