

iJournaling: Fine-Grained Journaling for Improving the Latency of Fsync System Call

Daejun Park and Dongkun Shin
Sungkyunkwan University, Korea
pdaeju@skku.edu, dongkun@skku.edu

Why fsync() latency is important?

- **Fsync() system call**
 - used by many applications to guarantee durability of a file
 - blocks until the flushing data is completed.

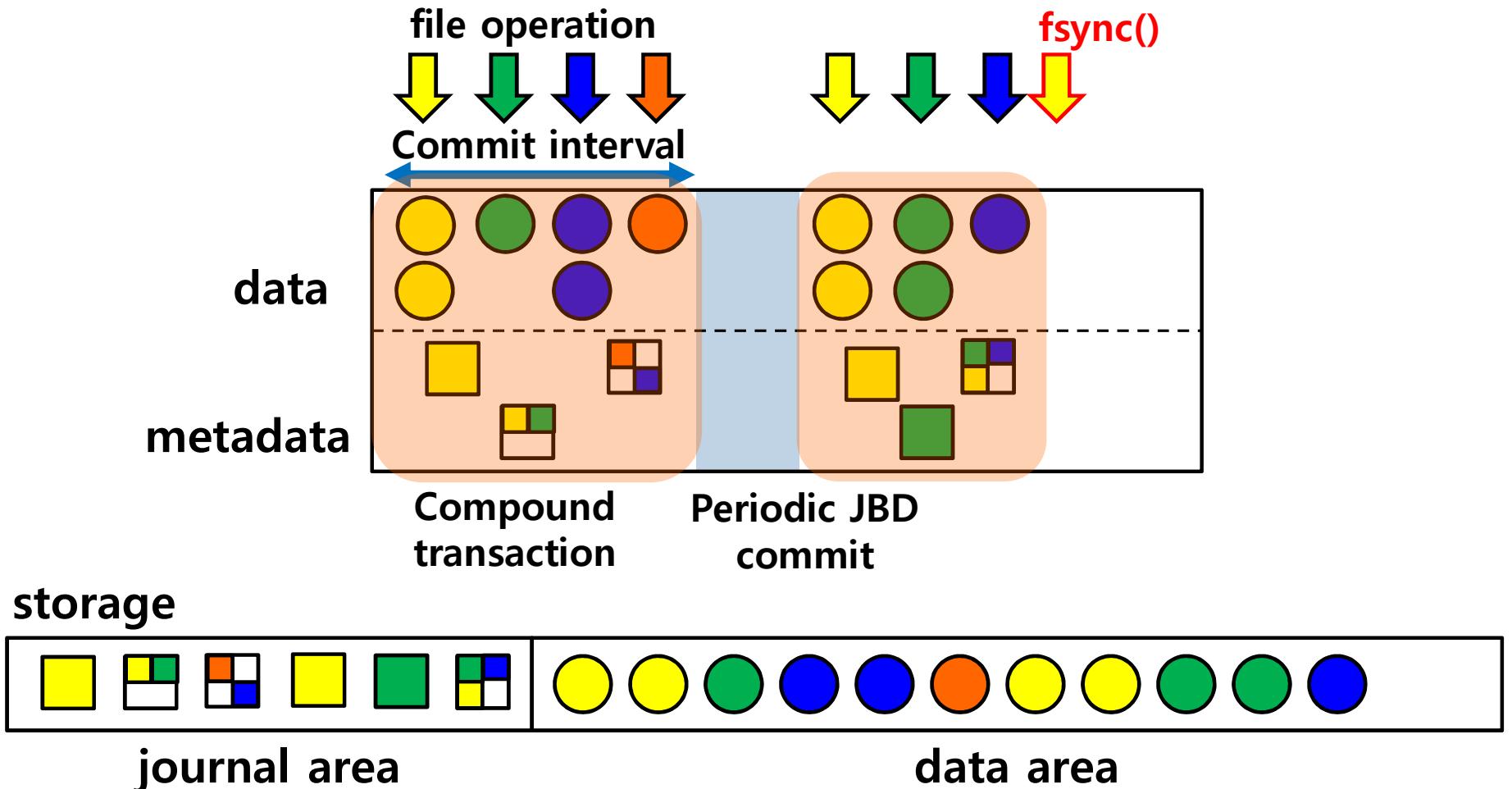
- **Database management system** 
 - MYSQL (tpcc-mysql) calls fsync() about **140 call/s**

- **Smartphone application**

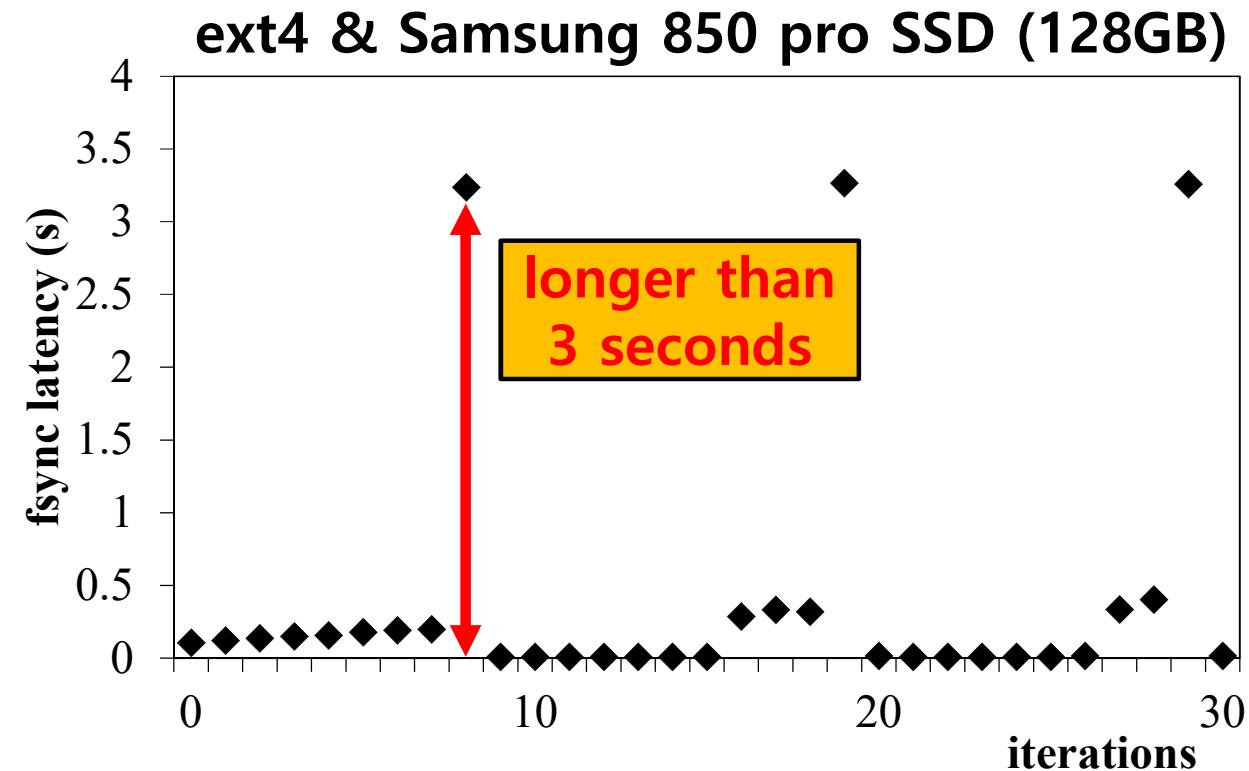
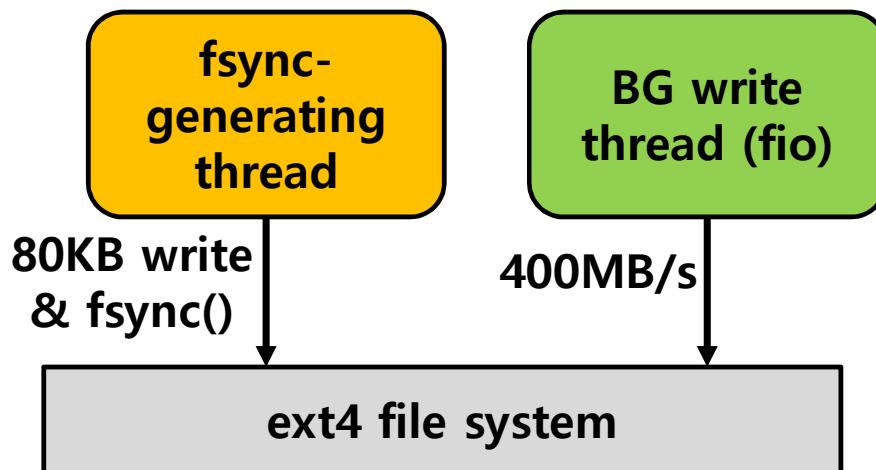
-  – A burstshot (20 pics) at smartphone calls fsync() **100 times**
-  – Adding a phone number on contract app calls fsync() **7 times**

Fsync latency can affect on application performance

Legacy journaling in the ext4

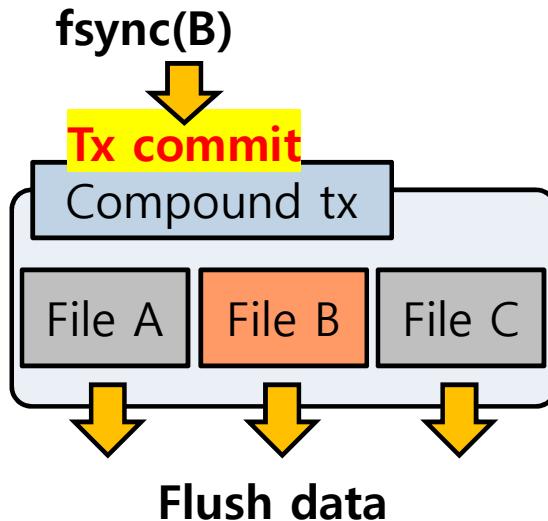


Long fsync() latency with heavy BG write

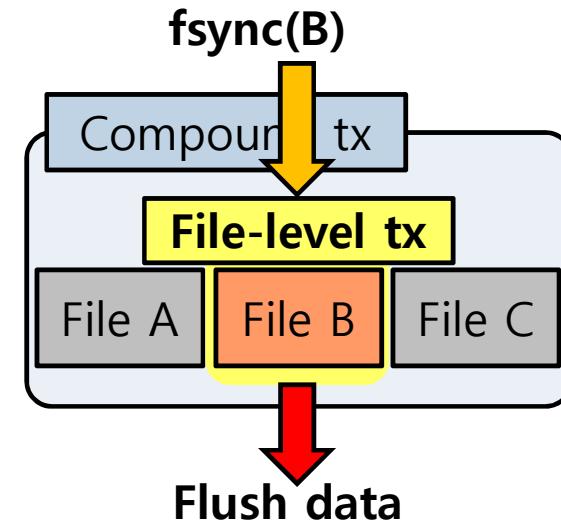


File level journaling

Normal journaling



File-level journaling

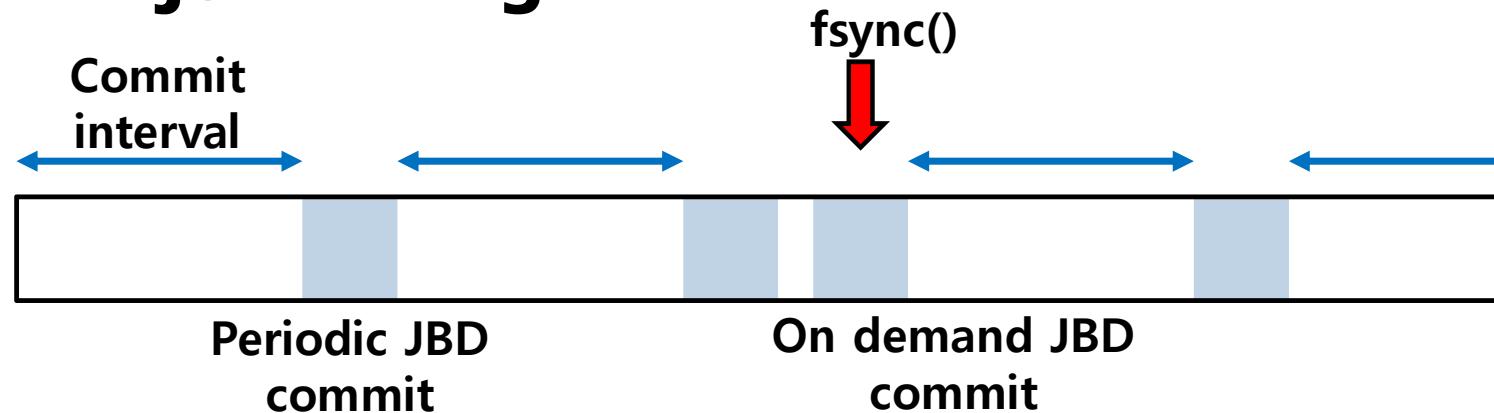


- **File-level journaling is needed for fsync()!**

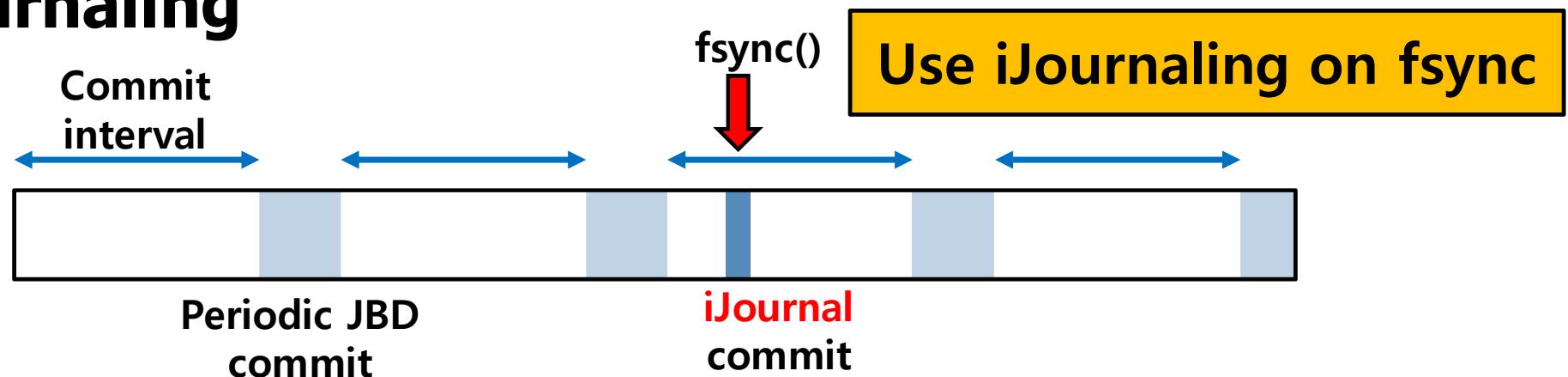
iJournaling

Hybrid journaling

- Normal journaling



- iJournaling



Related Work

- **CCFS[FAST '17], PBG[ISCE '14], Eager synching[NVMSA '14], IceFS[OSDI '14], Spanfs[ATC '05]**
 - Minimizes the compound transaction problem by isolation
- **ScaleFS[*], ZFS**
 - uses a logical logging technique.
 - performance overhead occurs because each file-system operation must record its own log
- **Xsyncfs [TOC '08], NoFS [FAST '12], and OptFS [SOSP '13]**
 - improved the fsync latency by delaying sync operations or changing the implementation of ordering constraint.

[*] Rasha Eqbal. ScaleFS: A multicore-scalable file system. Master's thesis, Massachusetts Institute of Technology, August 2014.

Challenges for iJournaling

1. Physical logging

- One metadata block includes info. of multiple file operations

2. Inter-file dependency

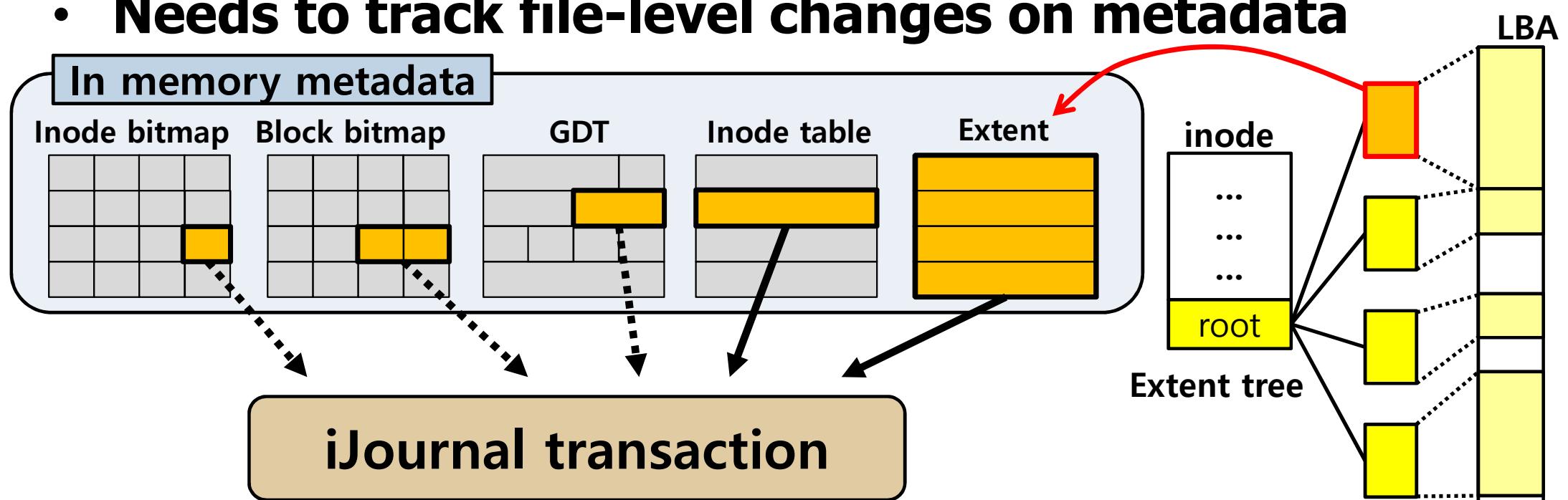
- Needs to flush files related to fsynced file
- e.g. parent DE, hard link

3. Managing iJournal area

- Sharing normal journal area or separated iJournal area

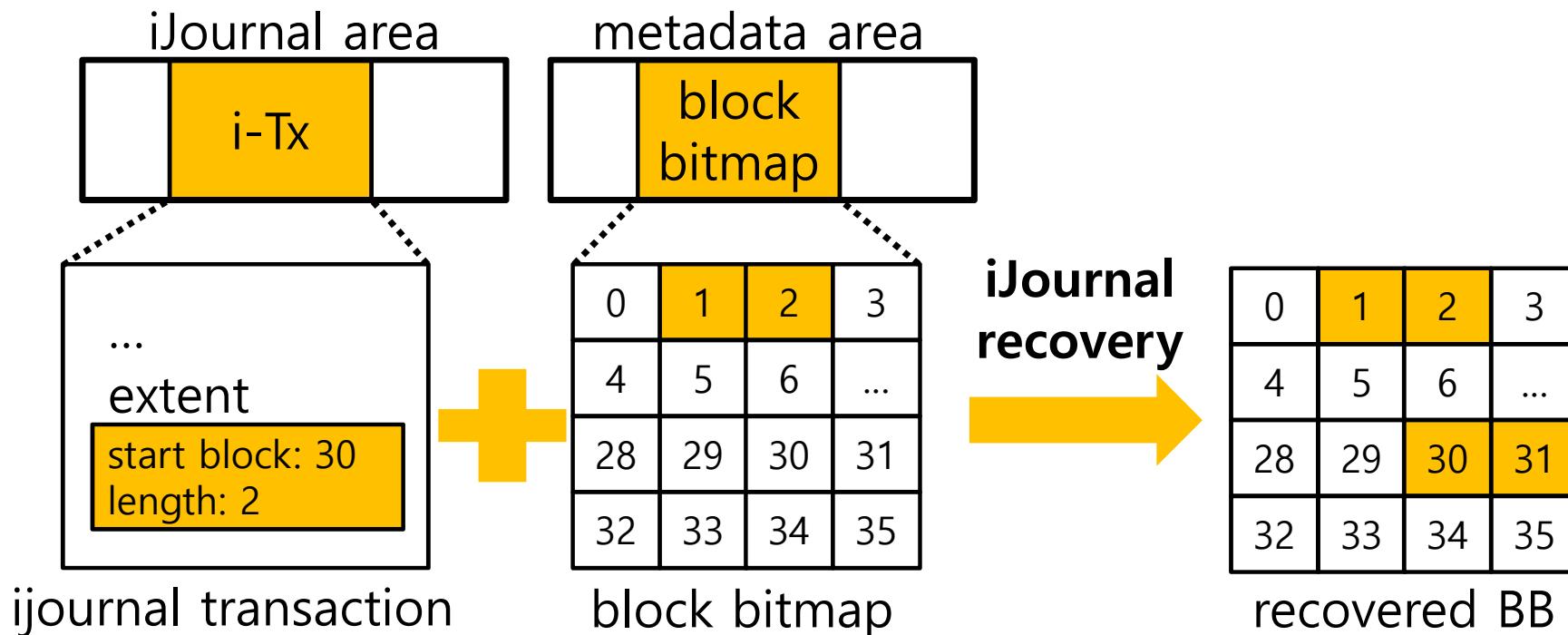
Legacy journaling – physical logging

- Physical logging with 4KB block granularity
 - multiple file operations share metadata
- Needs to track file-level changes on metadata



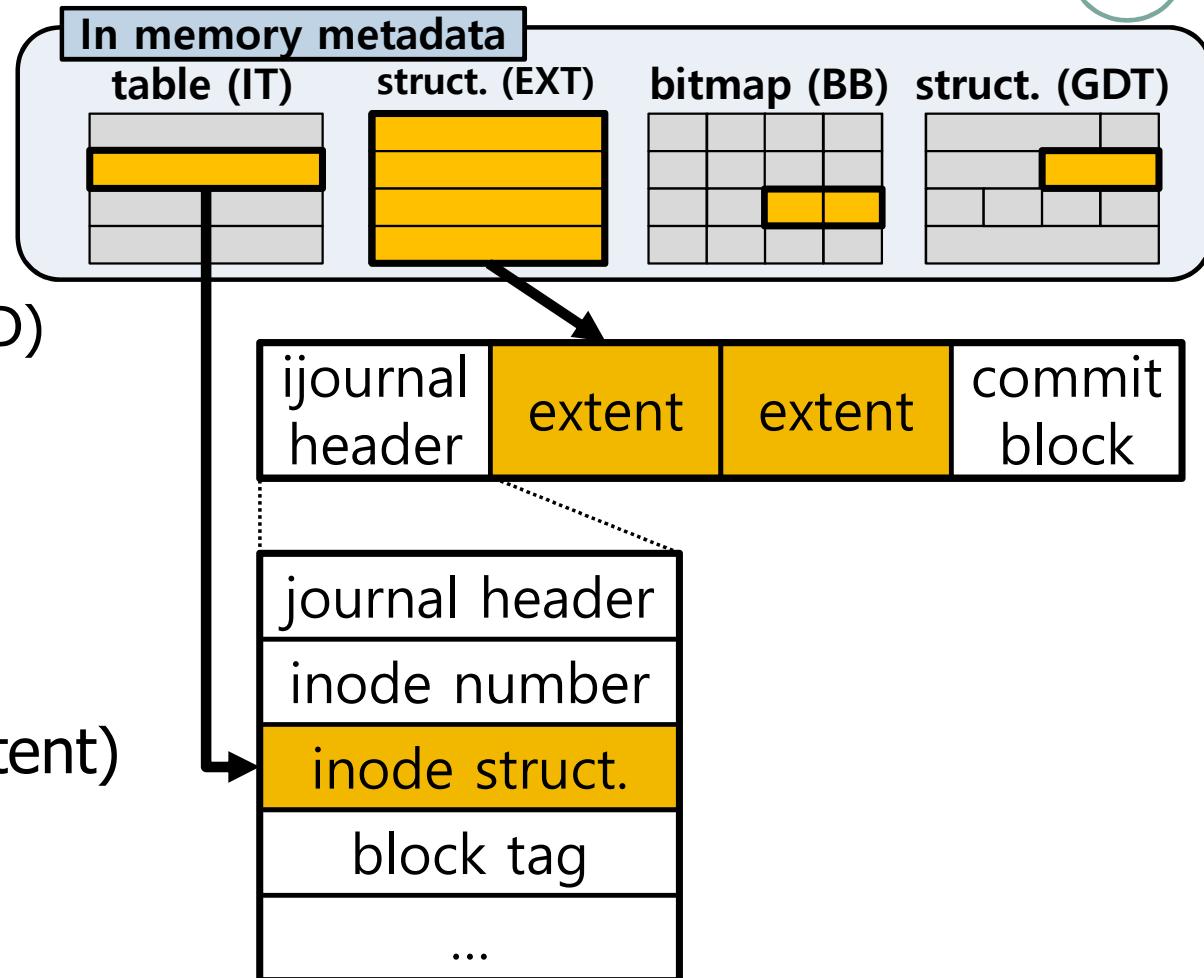
iJournaling – Logging and Recovery

- Compare & update metadata



iJournal transaction

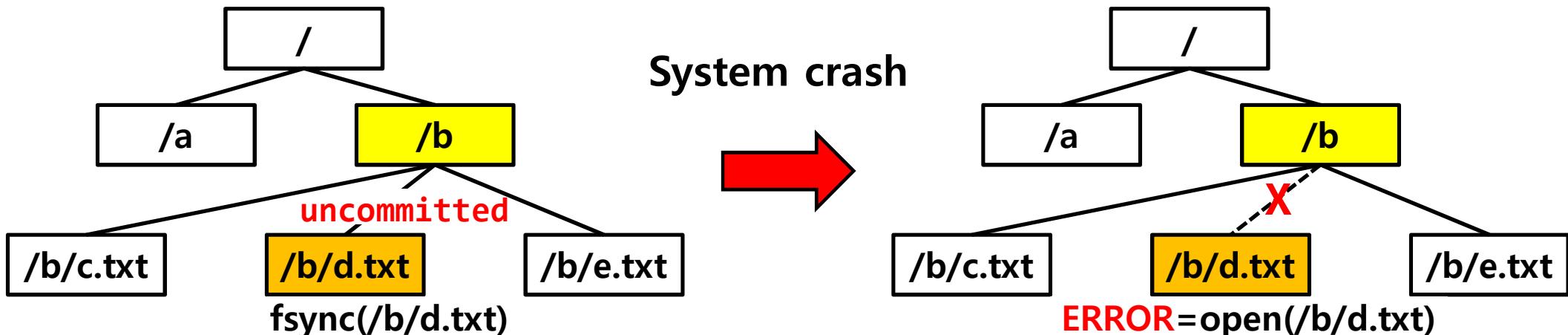
- **iJournal header block**
 - journal header
 - magic number
 - transaction ID (+ sub-txID)
 - inode number
 - inode structure
 - block tag
- **Journal**
 - metadata contents (ex. extent)
- **Commit block**



Inter-file dependencies

- **Parent DE**

- If a file is fsynced but its parent directory entry is not committed
→ **unreachable**
- Log DE of uncommitted parent directory

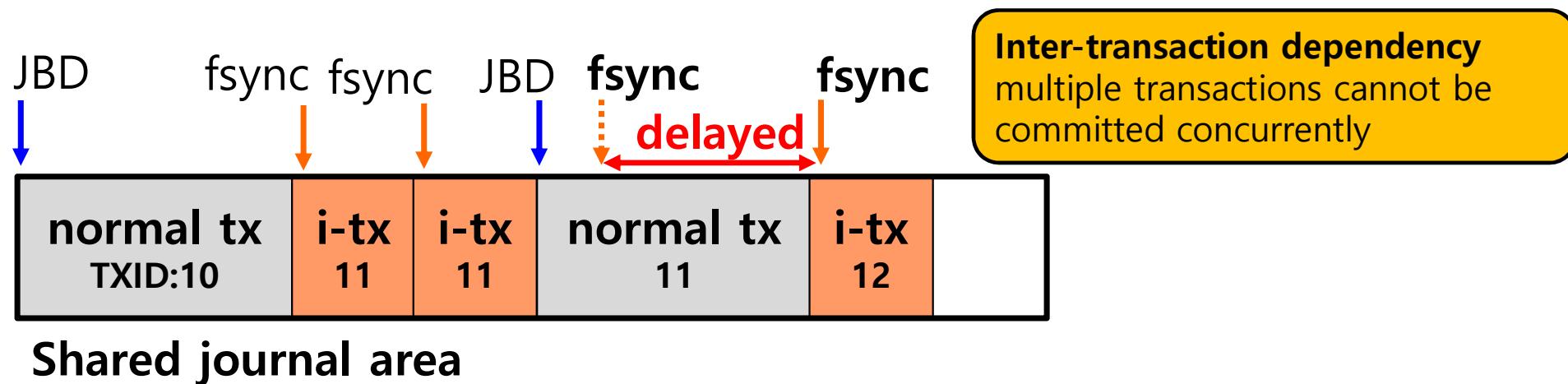


- **Uncommitted hard link modification of fsynced file**

Managing iJournal area

- **Shared journal area**

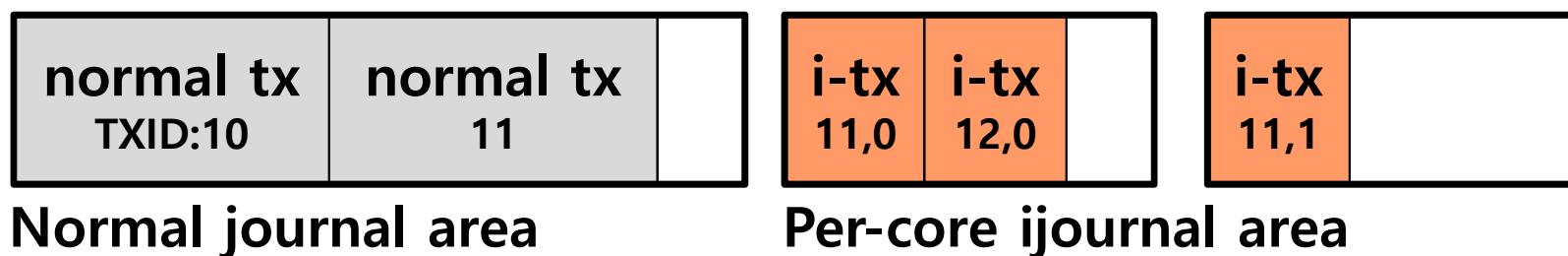
- No required additional space
- iJournaling performed at the service routine of fsync()
- However, ijournaling must wait until the block allocation of JBD is completed



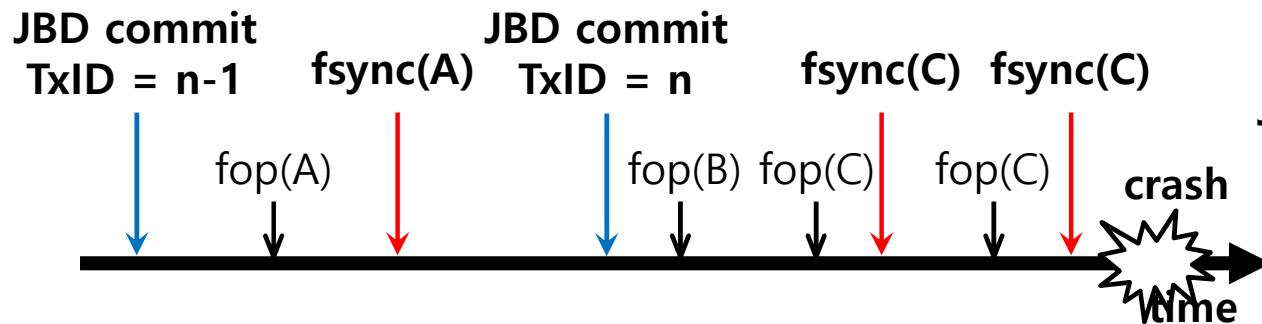
Managing iJournal area

- **Per-core iJournal area**

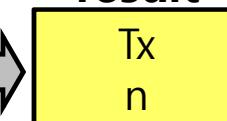
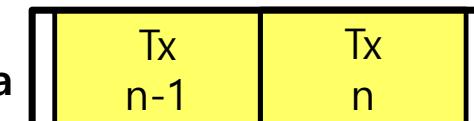
- Required to additional space for per-core iJournal area
- separating journal areas can improve the concurrency of journaling operations
- sub-TxID is needed for sorting i-txs in order



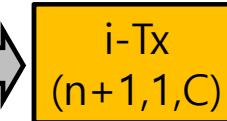
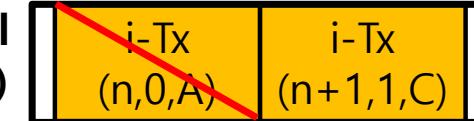
iJournal recovery



Normal
Journal area



iJournal
(core 0)



iJournal
(core 1)



- iJournaling recovery process**

1. normal journal recovery
2. file i-tx recovery
 - Recover block allocation information using extent
3. dir i-tx recovery
 - Recover DE to make accessible for fsynced file

Implementation

- **iJournaling**
 - Based on ext4 file system (ordered mode)
 - iJournaling on fsync() system call applied
 - Delayed allocation enabled
 - boosting technique[FAST '15] optionally applied
- **Desktop**
 - Linux kernel 4.7.3
 - Samsung 850 Pro SSD
- **Smartphone**
 - Linux kernel version: 3.4.5
 - Android OS version: 4.2.2 (Jelly Bean)
 - eMMC 32 GB

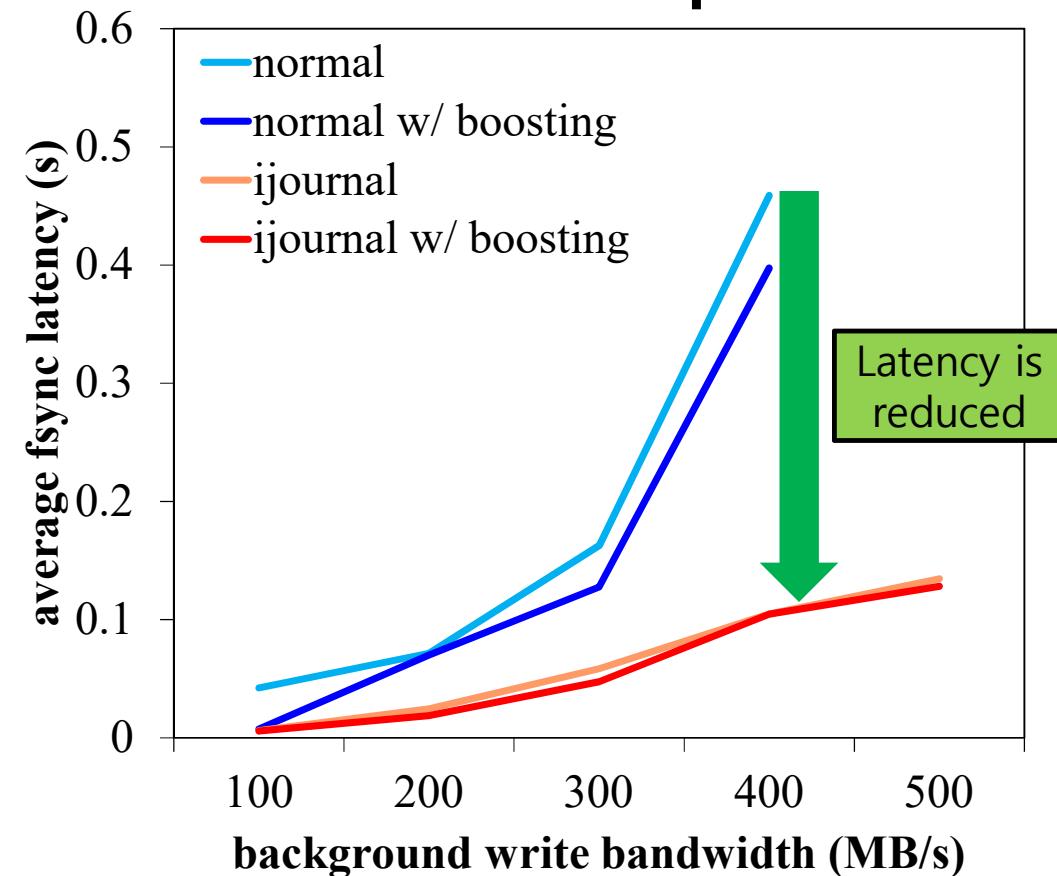
[1] Jeong et al. Boosting quasi-asynchronous I/O for better responsiveness in mobile devices. FAST '15

Fsync latency on various BG write traffic

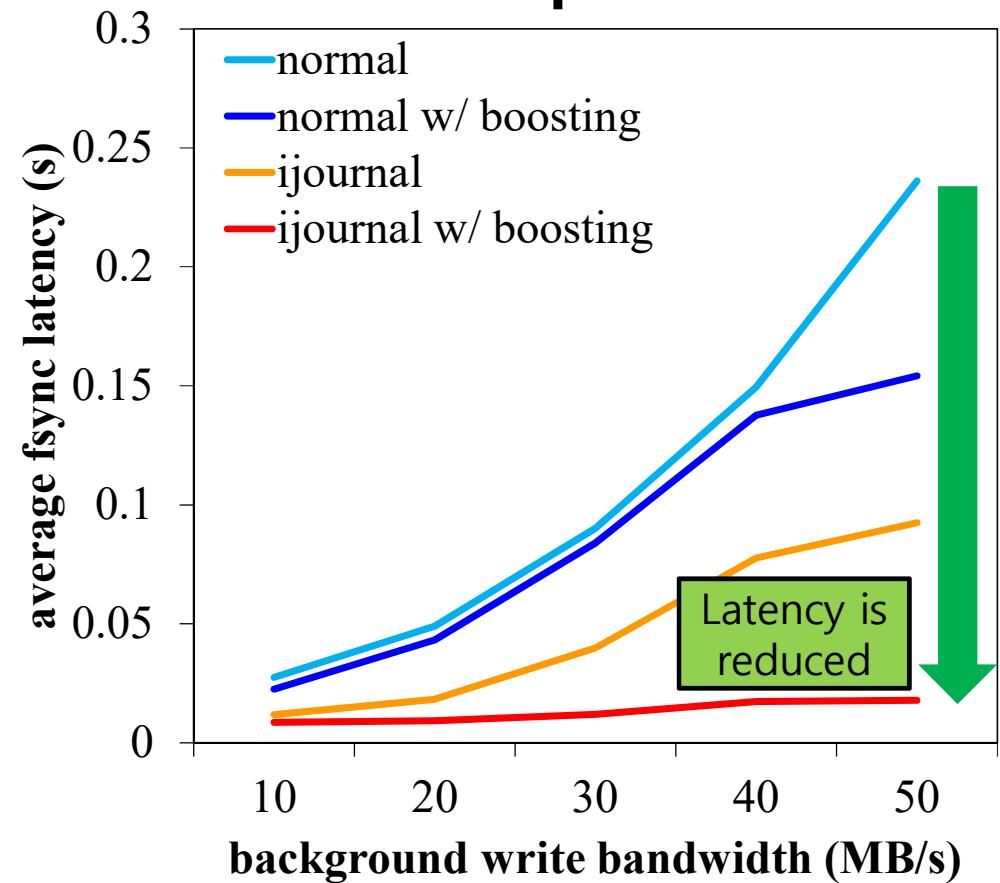
17

21

desktop



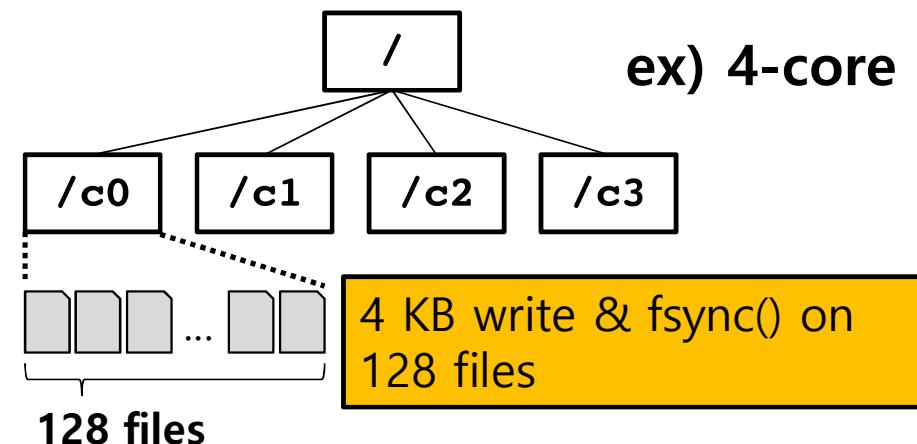
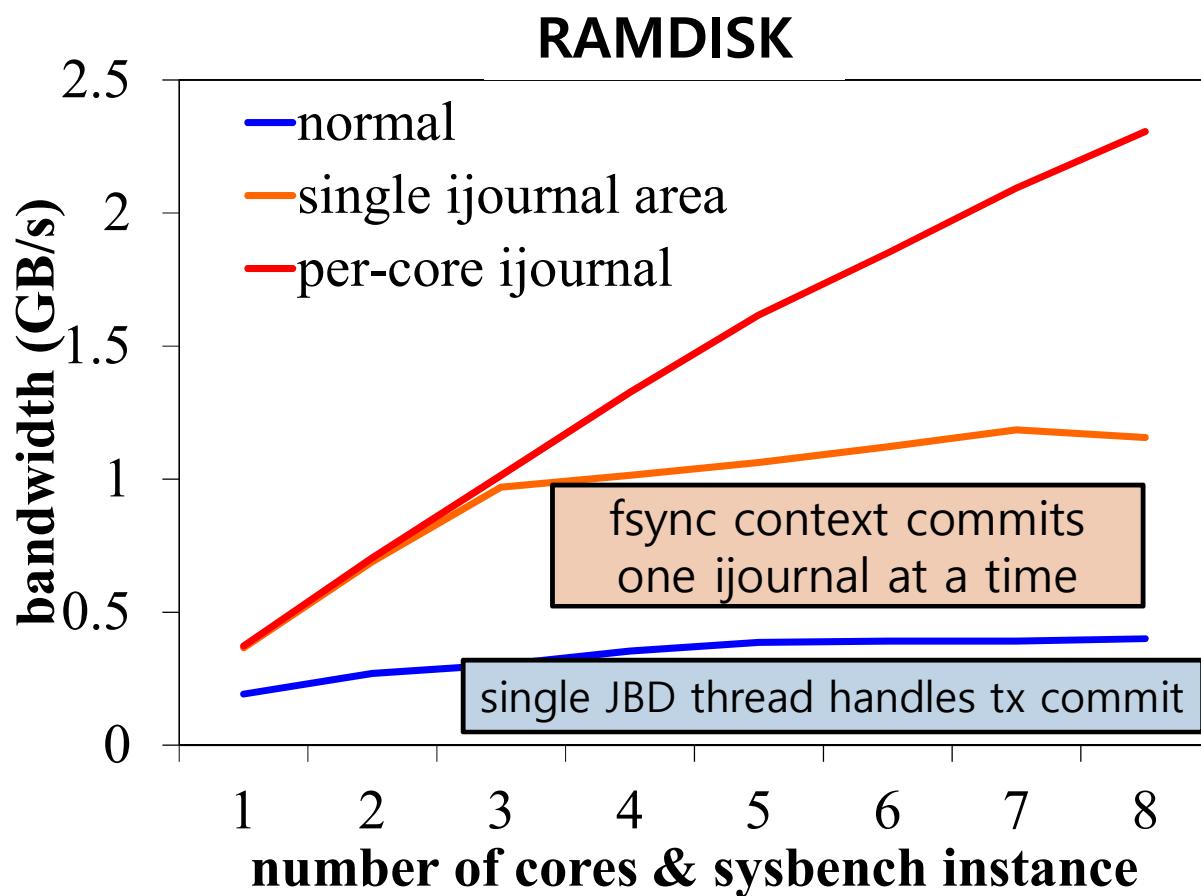
smartphone



Multicore scalability on iJournaling

18

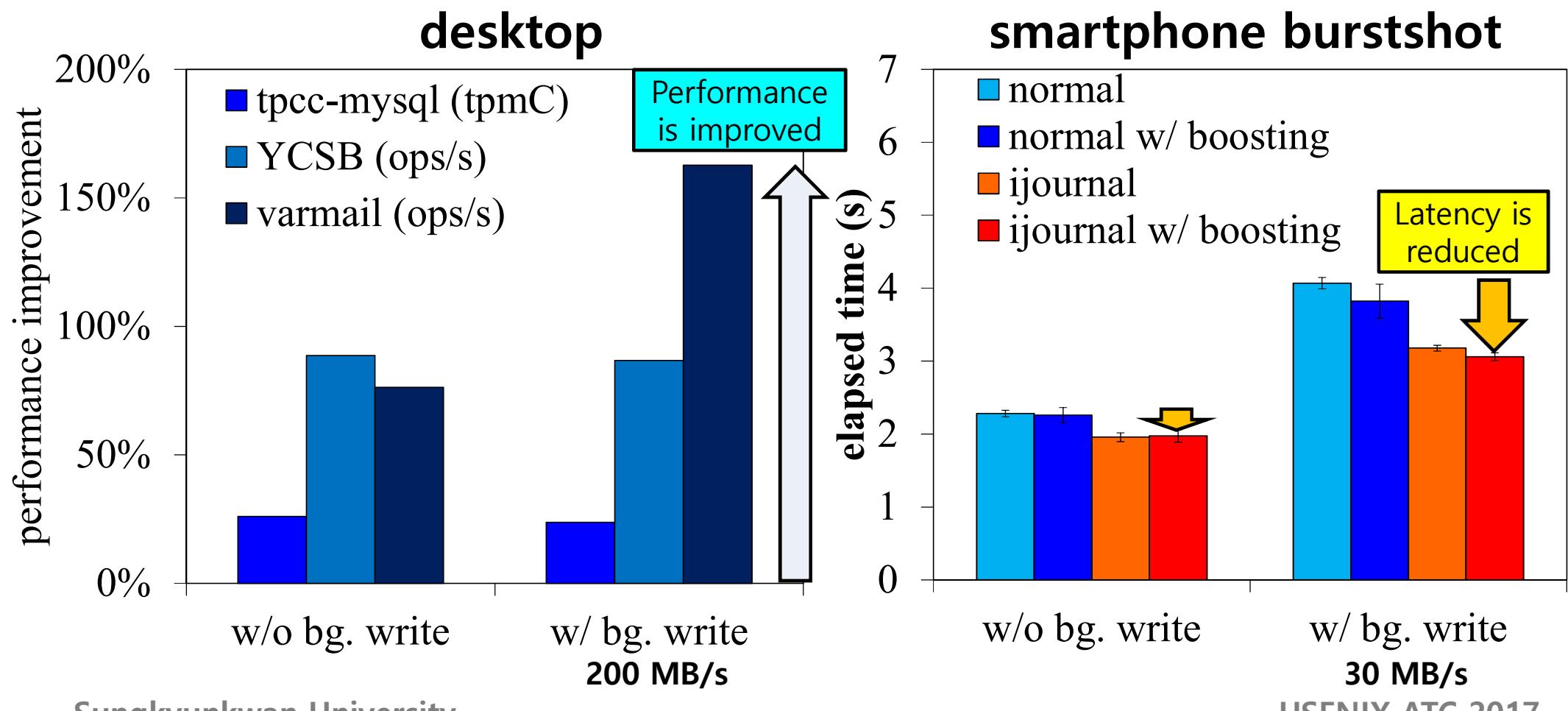
21



Real workload on iJournaling

19

21



Conclusion

20

21

- **fsync() system call can be delayed under the compound transaction scheme**
- **a hybrid journaling technique, called ijournaling**
 - journals only the related file-level transactions of an fsync call
 - recovers the file-system consistency through file-level journals upon a crash recovery
- **We implemented ijournaling and showed significant improvements to the fsync latencies**

Thank you!

21

21

Q&A