

Improving File System Performance of Mobile Storage Systems Using a Decoupled Defragmenter

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Outline

- **Impact of File Fragmentation/Defragmentation**
- **Key Observations on Flash-based File Fragmentation**
 - ◆ **Decoupled Fragmentation**
 - ◆ **Dominant Impact of Logical Fragmentation**
- Janusd: a Decoupled Defragmenter
- Experimental Results
- Conclusions

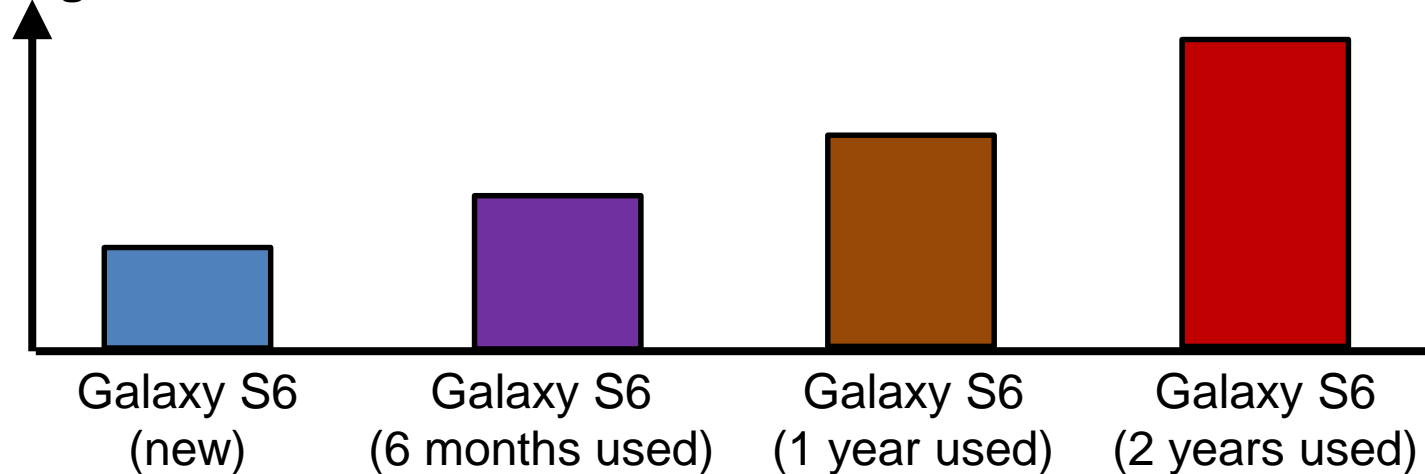
Gradual Performance Degradation on Smartphones



Performance of Android smartphones
gradually degrades as smartphones *age*

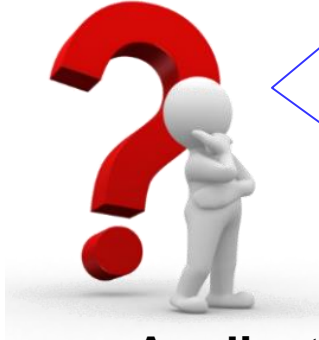


Application
launching time



Application launching times increase up to 3 times on 2-year used smartphones

Root Cause: File Fragmentation

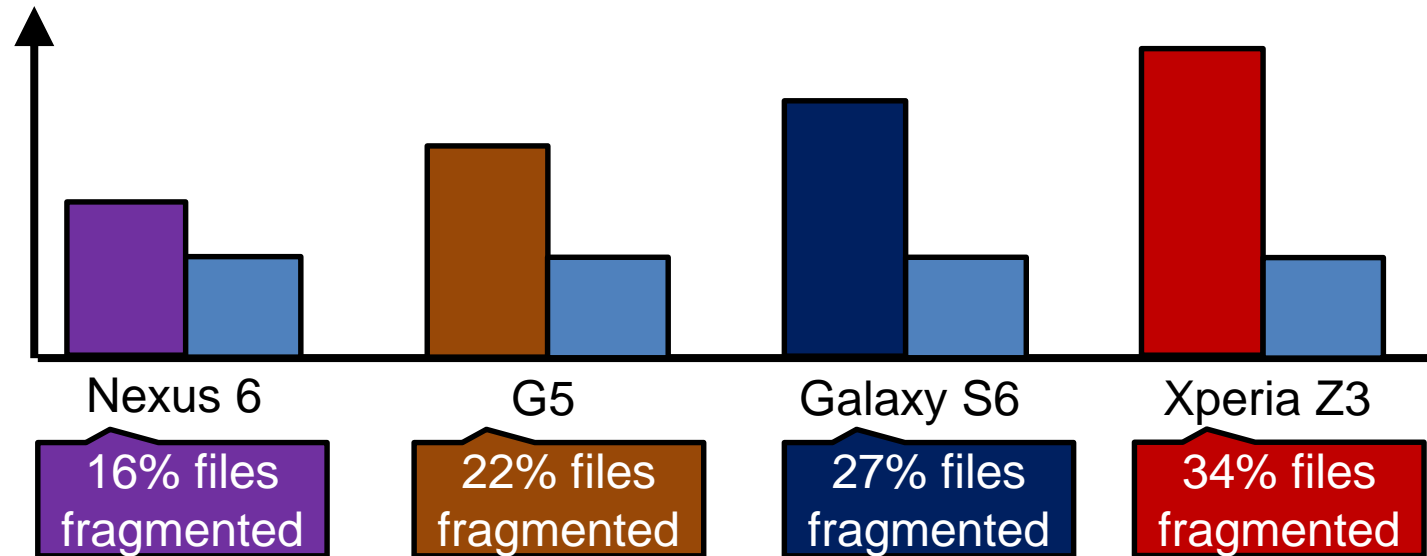


Q: **Why** does performance degrade ?

A: File fragmentation



Application launching time



Defragmentation can improve the degraded performance by fragmentation

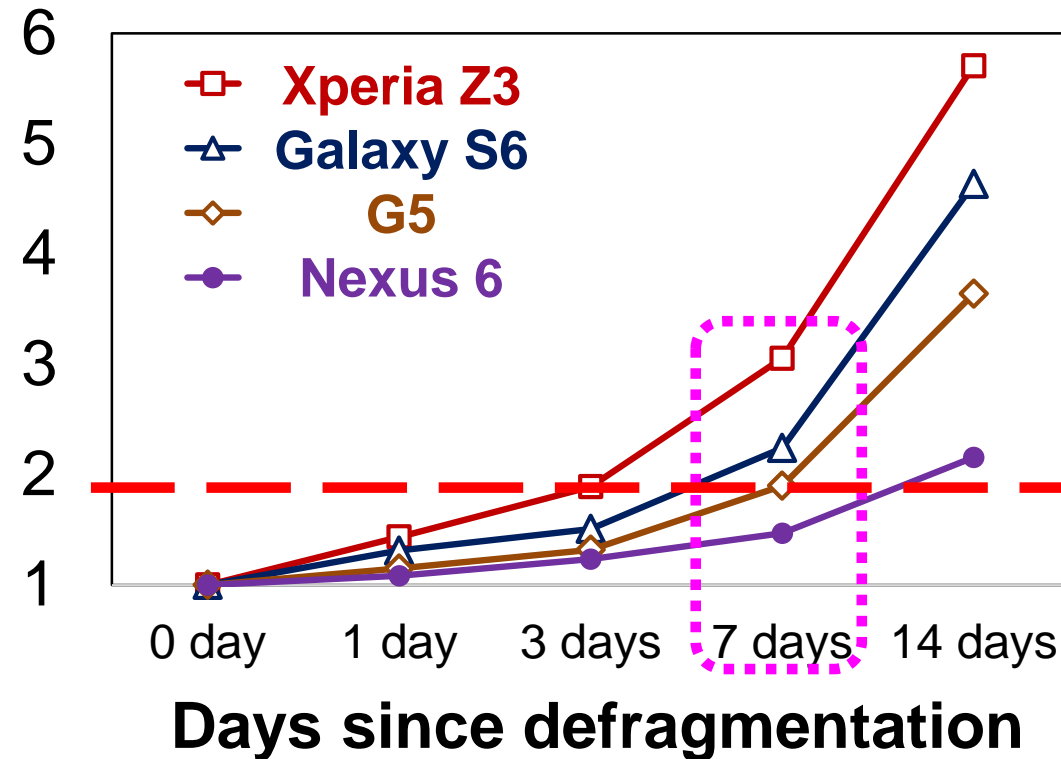
File Fragmentation Recurrence on Smartphones



Q: **How often** should we defrag smartphones ?

Average
of fragments
per file

Degree of
file fragmentation



User begins
to experience
performance
degradation



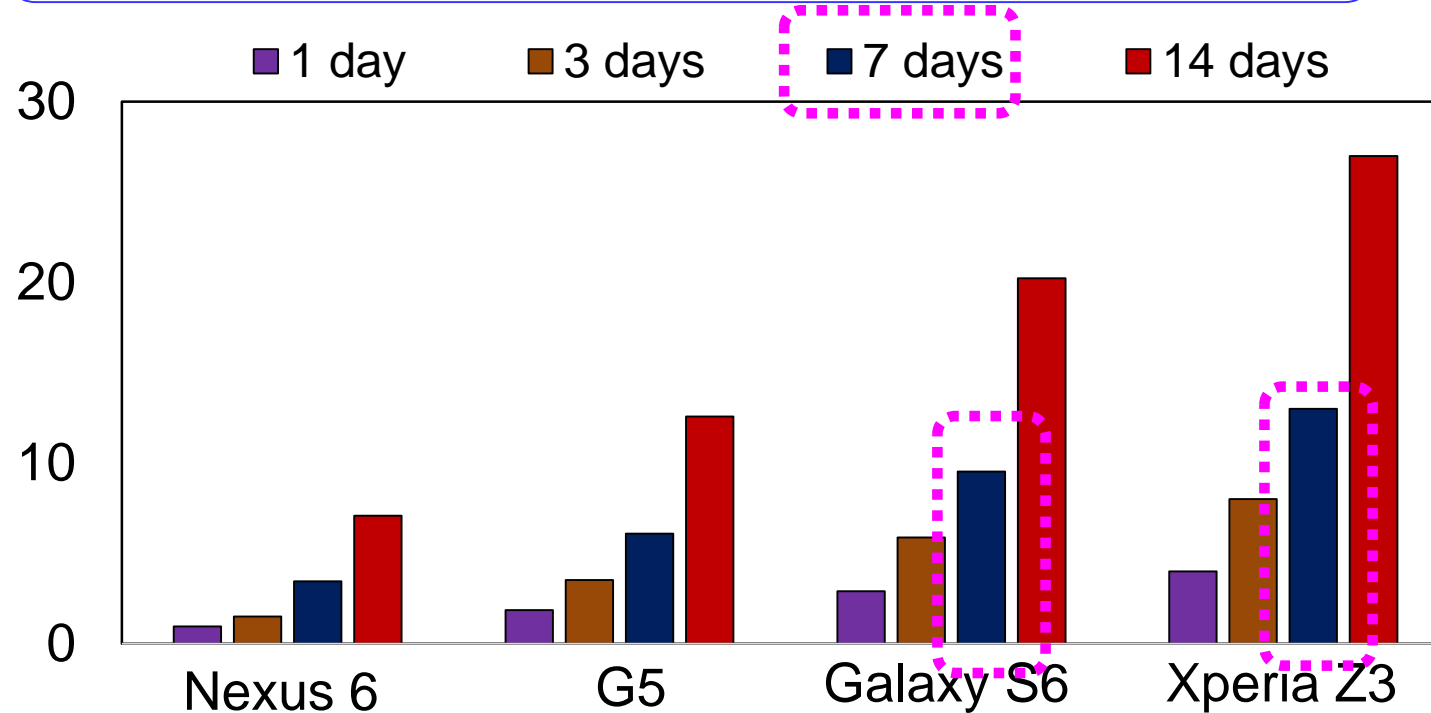
File fragmentation recurs even in a week

Lifetime Impact of Frequent Defragmentation



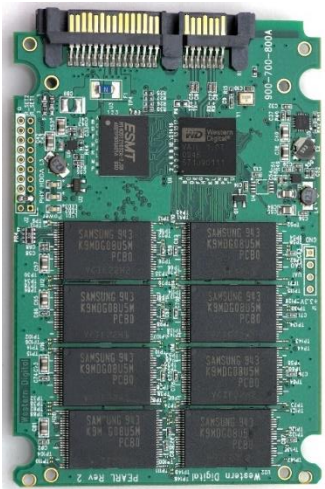
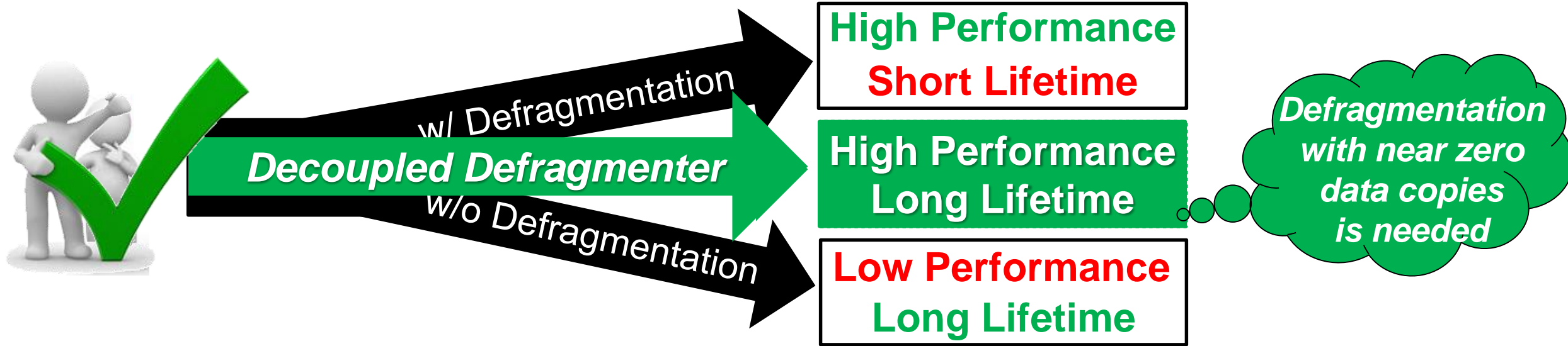
Q: Are there any side effects from frequent defragmentation?

The amount of data copies by defragmentation (GB)



Weekly defragmentation can reduce the storage lifetime by more than 10%

Summary: Impact of File Fragmentation/Defragmentation

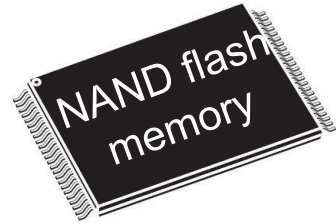


File fragmentation in NAND flash-based storage is **quite different** from conventional one in HDD

1. **Decoupled fragmentation**
2. **High overhead of logical fragmentation**

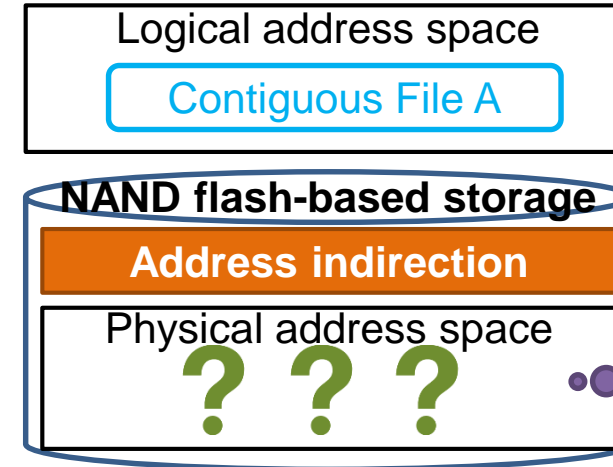
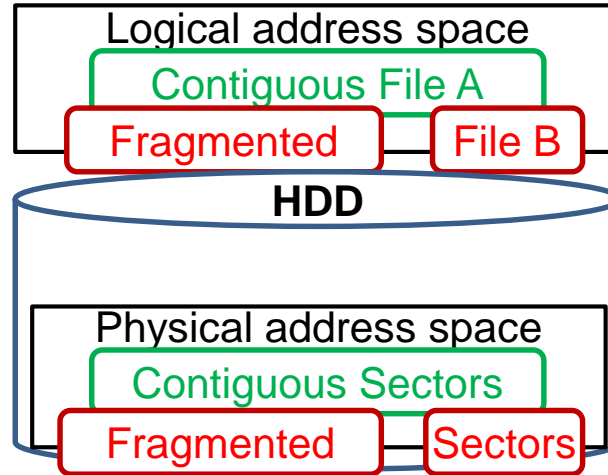
NAND Flash-based
Storage

Observation 1: Decoupled Fragmentation



All data are stored using *address indirection*

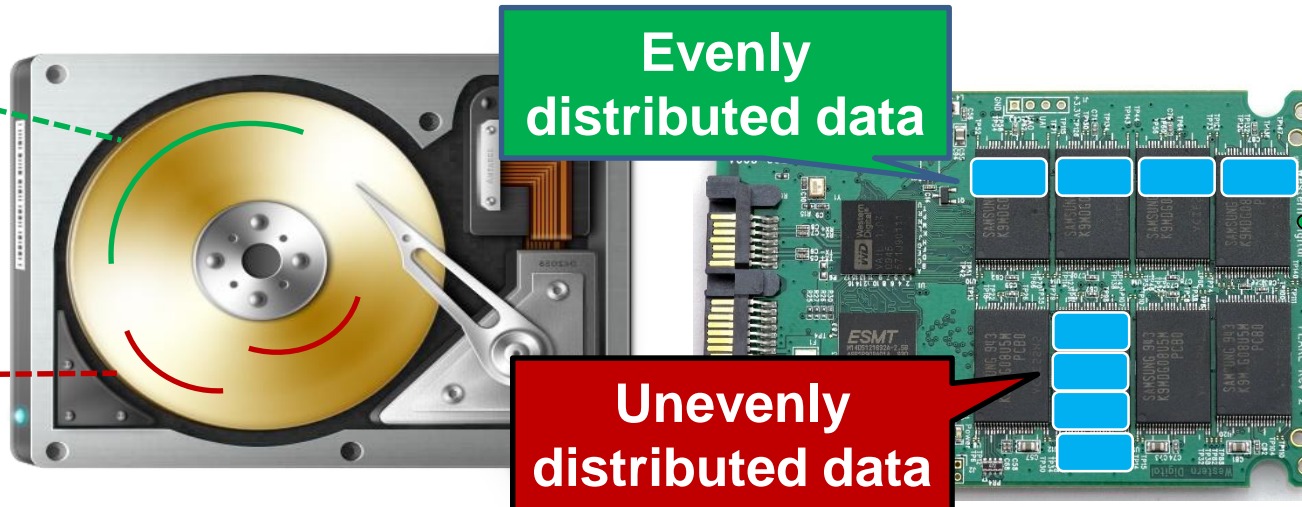
Fragmentation at *logical space* and *physical space*



Physical fragmentation

Contiguous File A

Fragmented File B



Evenly distributed data

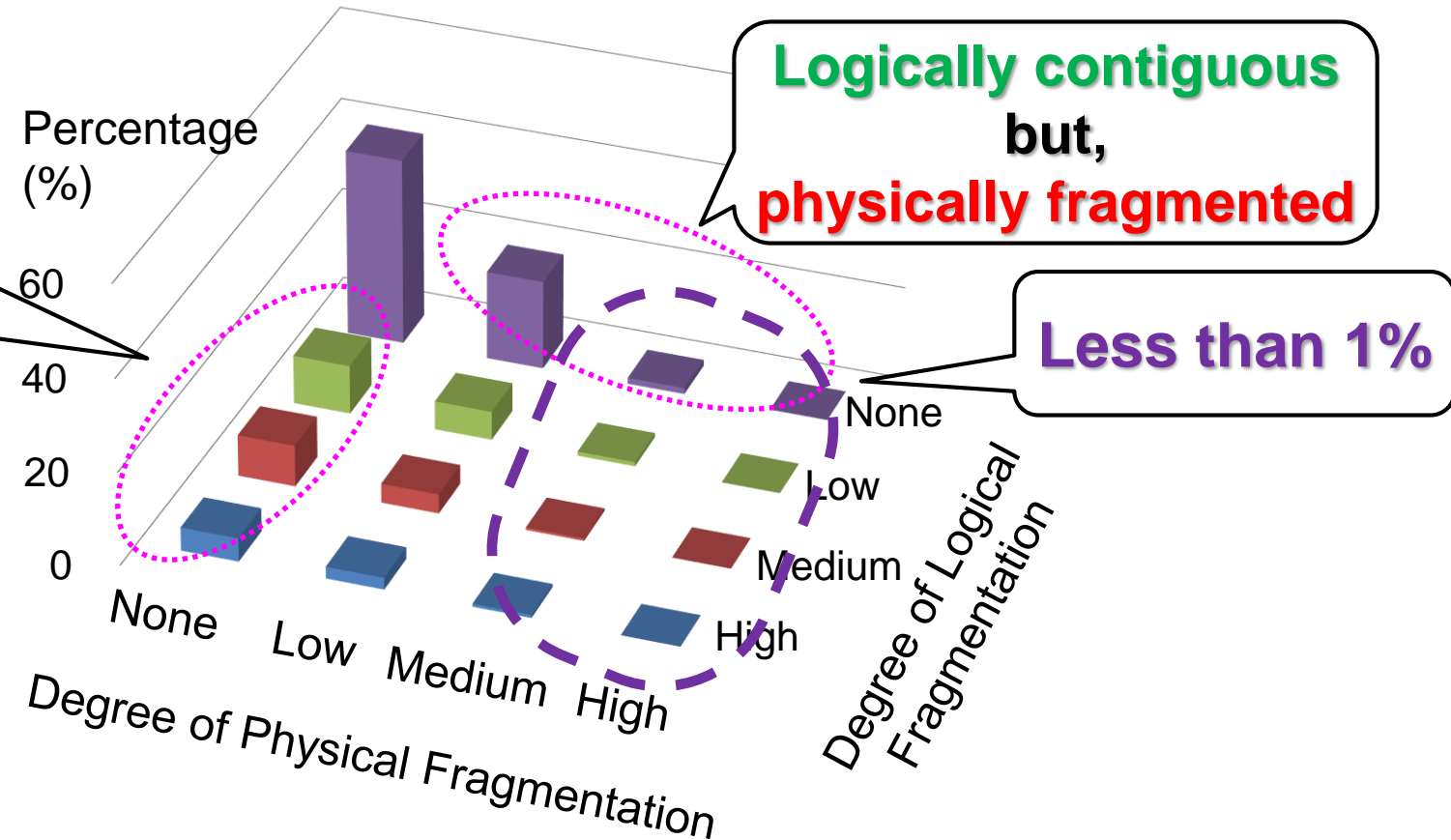
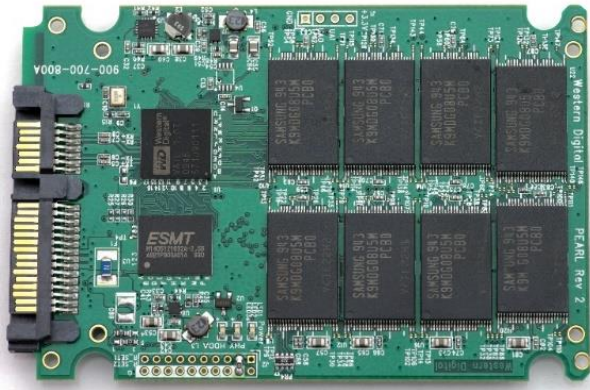
Unevenly distributed data

High degree of I/O parallelism

Low degree of I/O parallelism

Observation 1: Decoupled Fragmentation

Logically fragmented
but,
physically contiguous

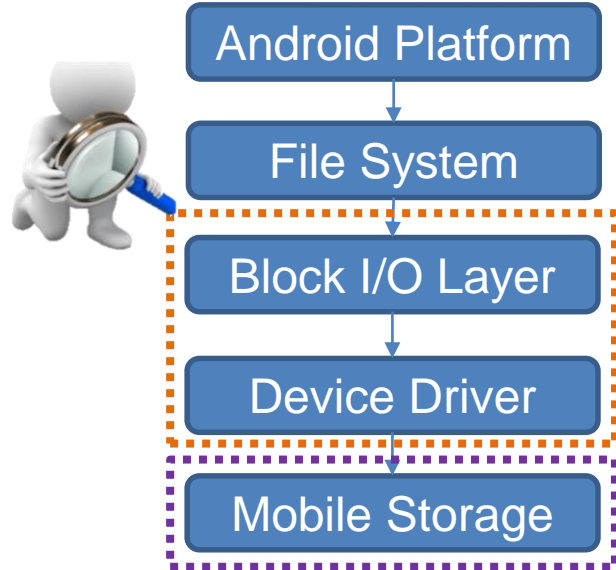


1. There is **no correlation** between **logical/physical** fragmentation
2. **Physical fragmentation rarely** occurs

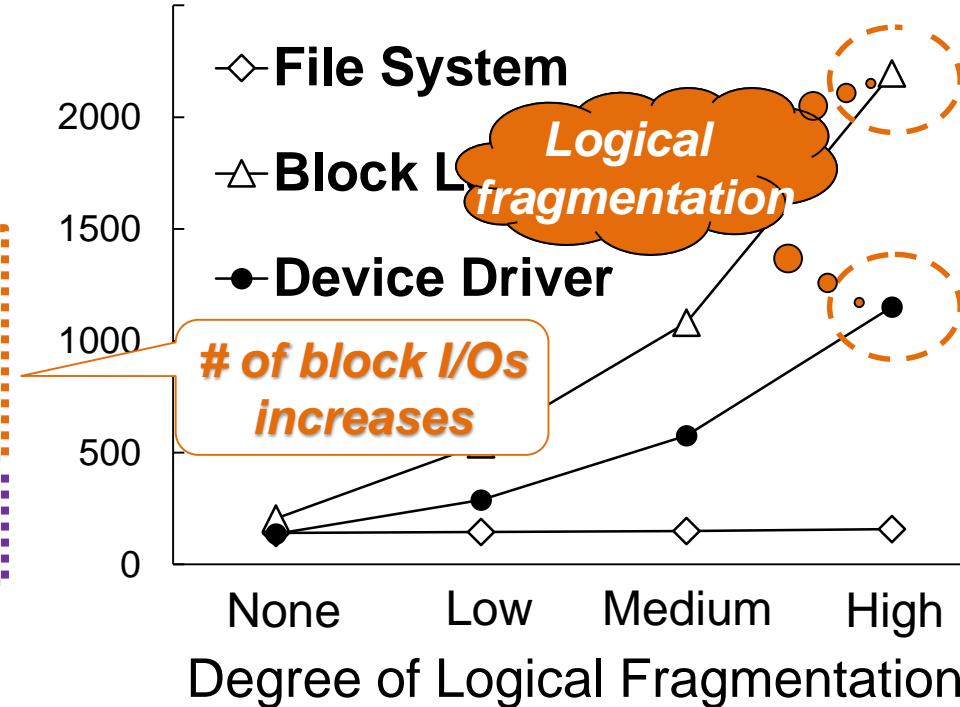
Observation 2: High Overhead of Logical Fragmentation



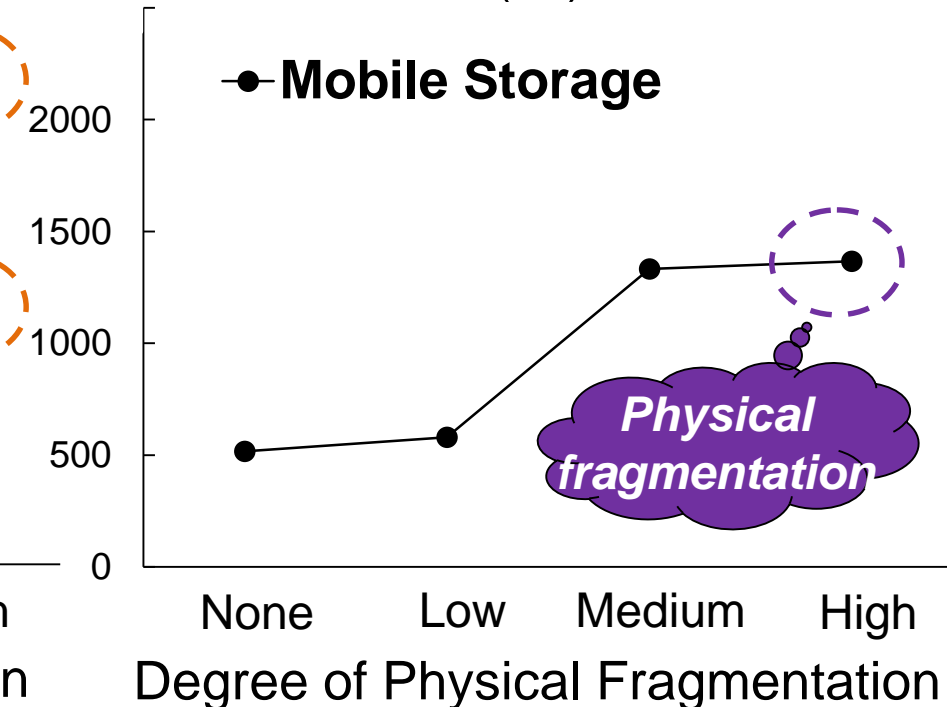
Q: How much the impact of logical/physical fragmentation on performance?



I/O Execution Time (us)

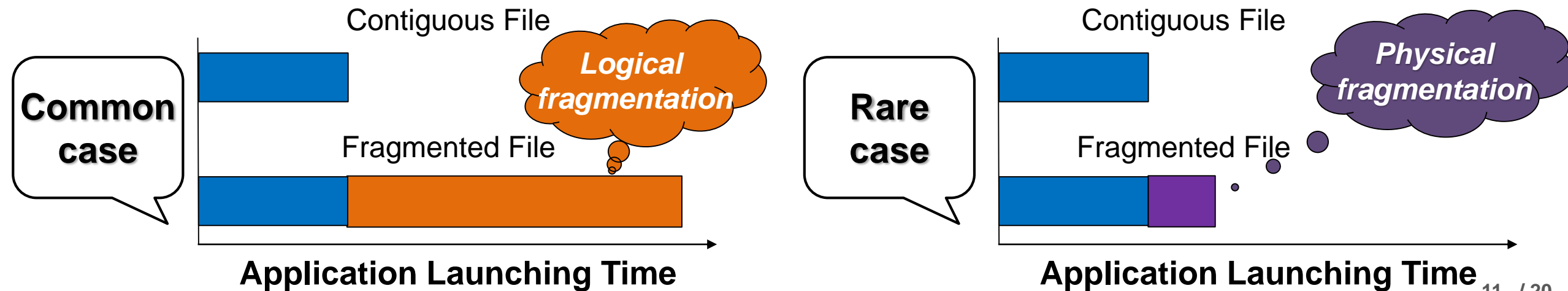
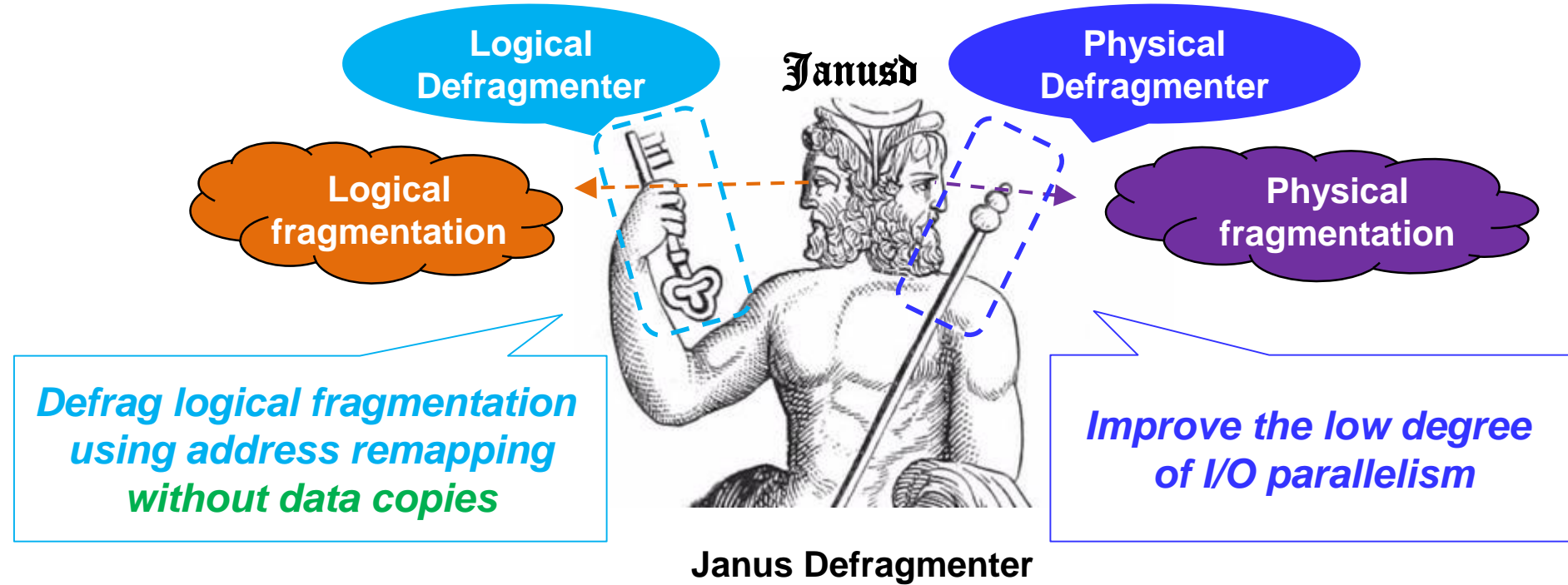


I/O Execution Time (us)



Logical fragmentation overhead overwhelms physical fragmentation overhead

Solution for Decoupled Fragmentation



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- Impact of File Fragmentation/Defragmentation
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 - ◆ Decoupled Fragmentation
 - ◆ Dominant Impact of Logical Fragmentation
- **Janusd: a Decoupled Defragmenter**
- **Experimental Results**
- **Conclusions**

Overview of Decoupled Defragmenter

Decoupled Defragmenter (Janusd)

e4defrag

*Logical
Defragmenter*

Detect
logical
fragmentation

Remap
LBAs of logical
fragments

Deliver
modified LBAs
*Use new custom
interface*

**Firmware
(FTL)**

*Physical
Defragmenter*

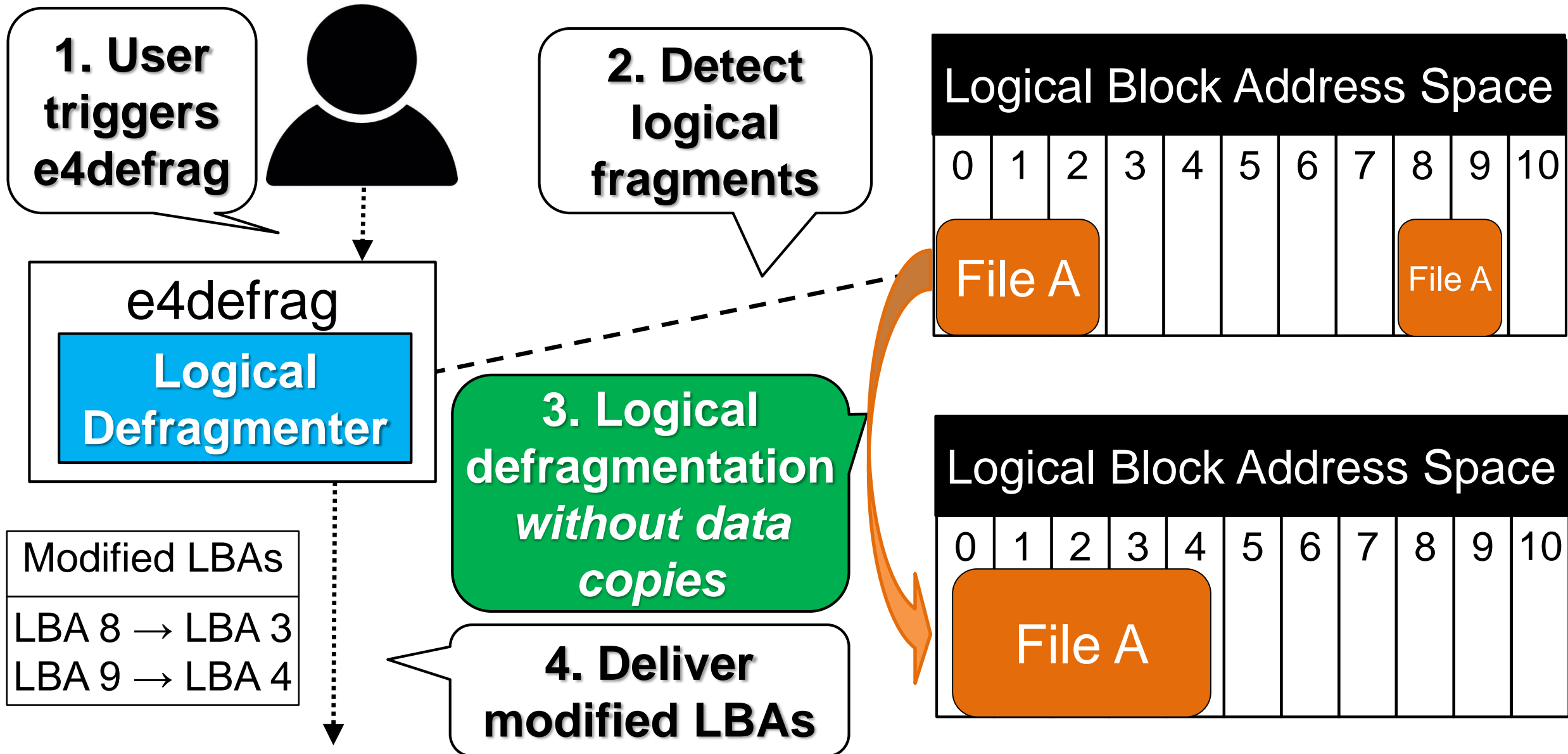
Detect
physical
fragmentation

Disperse
physical
fragments

Update FTL's
mapping table
*Maintain log for
reverse mapping*

*Improves I/O performance of mobile storage
while minimizing lifetime degradation*

Logical Defragmenter (JanusdL)



Logical Defragmenter (JanusdL)

Modified LBAs

LBA 8 → LBA 3

LBA 9 → LBA 4

**Logical
Defragmenter**

Mobile Storage

L2P Mapping Table

NAND Flash Memory

Defrag log

**6. Maintain
remapping
history**

L2P Mapping Table

0 **NAND PAGE A**

5. Remapping

4

8 **NAND PAGE D**

9 **NAND PAGE E**

L2P Mapping Table

0 **NAND PAGE A**

1 **NAND PAGE B**

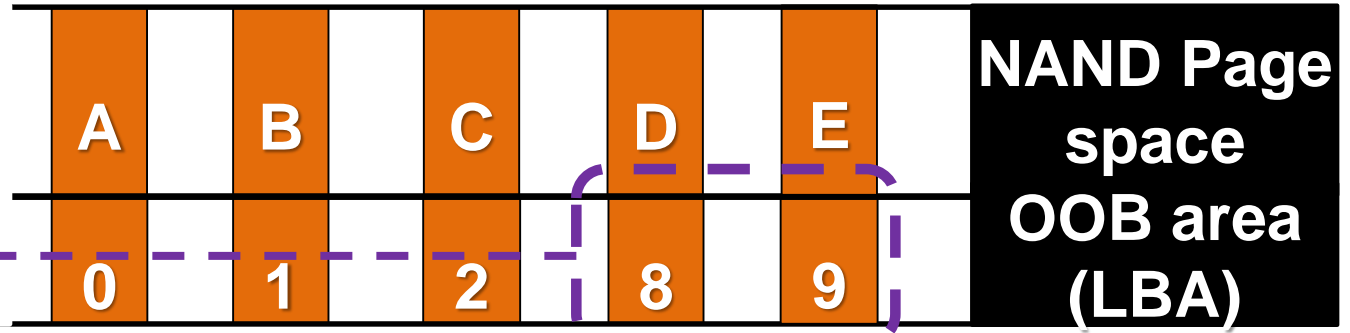
2 **NAND PAGE C**

3 **NAND PAGE D**

4 **NAND PAGE E**

8

9

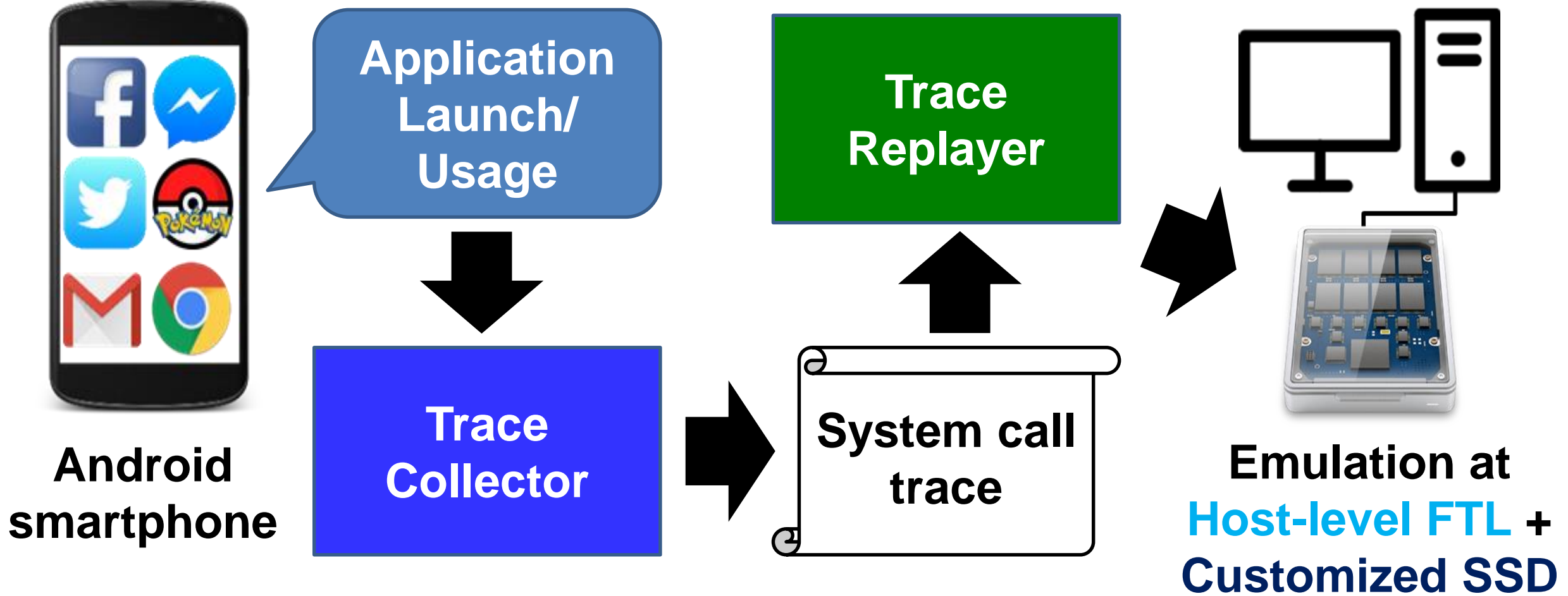


Evaluation Scenarios

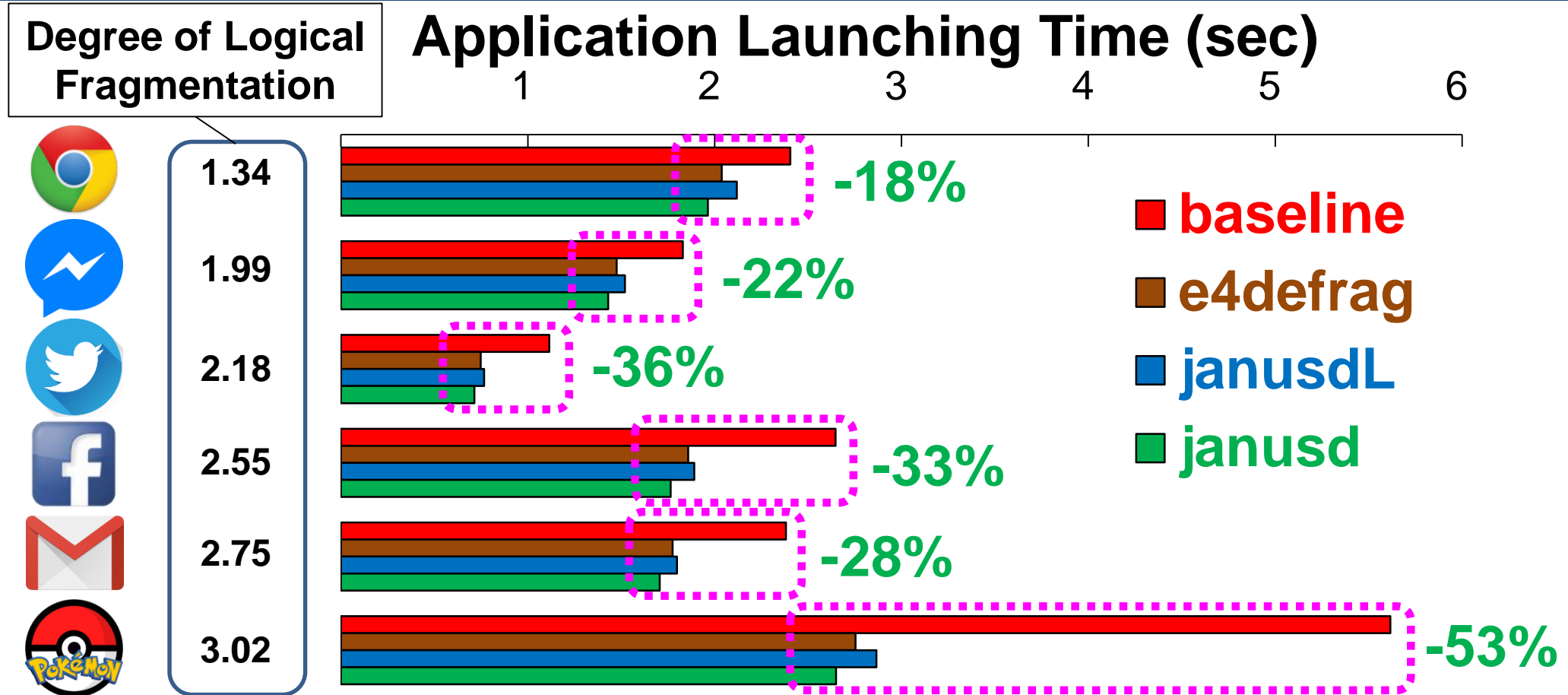
- **We collected six different application usage traces**
 - **Application launching scenarios**
 - **Interactive application usage scenario (10 minutes)**

Scenario	Scenario Description
Chrome	Launching app → Viewing webpages
Messenger	Launching app → Viewing chat records
Gmail	Launching app → Viewing emails
Facebook	Launching app → Viewing online news
Twitter	Launching app → Viewing online news
Game	Launching Pokemon Go → Playing game

Experimental Settings

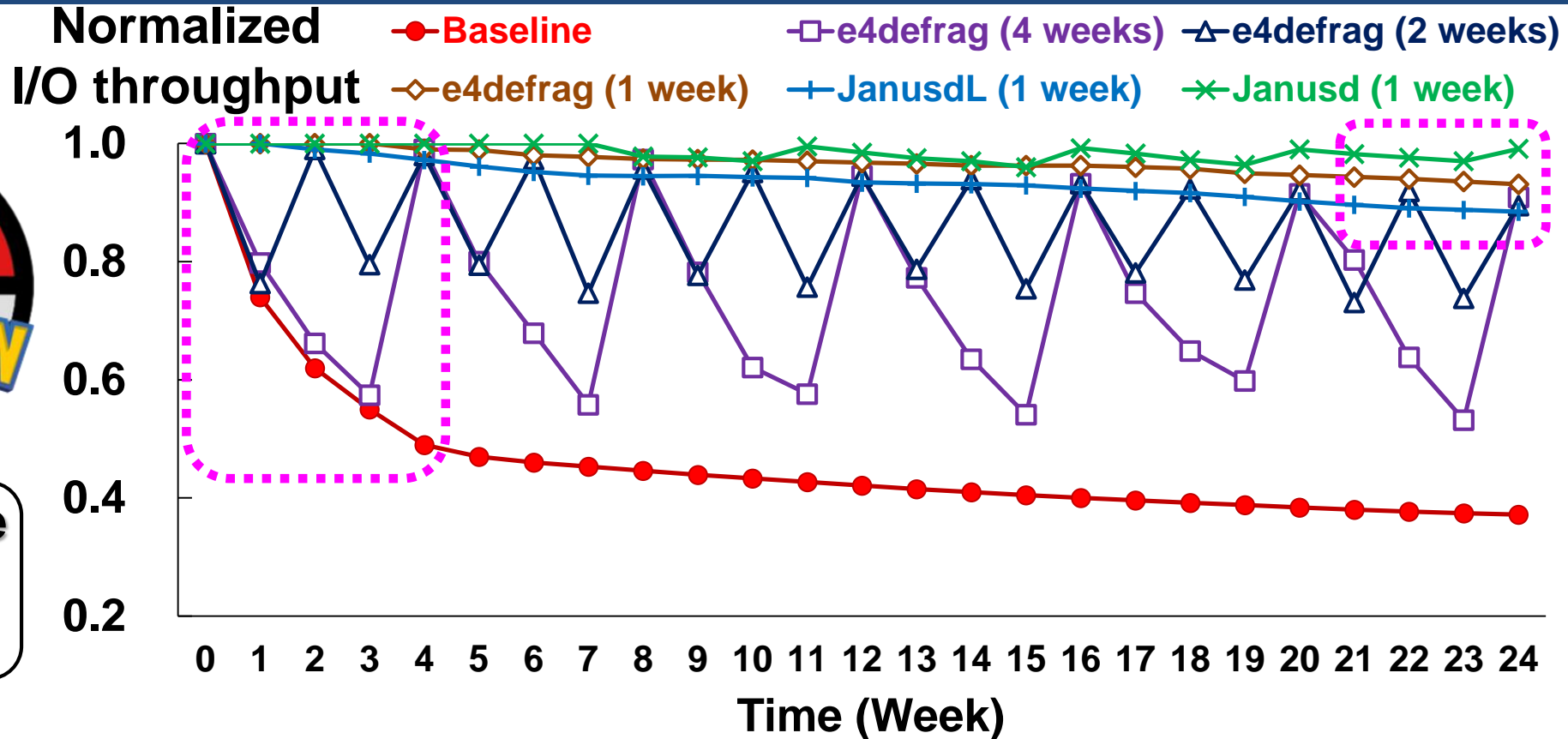


Result 1: Application Launching Time



1. **The more file fragmentation, the greater the performance improvement**
2. **Janusd** achieves better performance than **conventional defragmenter (e4defrag)**

Result 2: I/O Throughput



- 
1. **Performance degradation occurs** even when we defrag smartphone every 2 weeks
 2. **Conventional defragmenter** has limitations in solving physical fragmentation

Conclusion

- We have presented a **decoupled defragmenter** for improving the file system performance
 - JanusdL defrags logical fragmentation **without data copies** by **remapping** LBAs with FTL's mapping table
 - JanusdP defrags physical fragmentation by improving **I/O parallelism** of files
 - Improved application launching times by **32%** on average
 - Reduced the amount of data copies by **99.99%** on average
- Future expends
 - Free space defragmentation tool
 - Defrag-on-write() which triggers JanusdL right before write()

감사합니다 Natick
Grazie Danke Ευχαριστίες Dalu
Thank You Köszönöm
Tack
Спасибо Dank Gracias
谢谢 **Merci** Seé
ありがとう